## Programmable Controllers

## Preface

Thank you for purchasing programmable controllers, this manual is the basic instructions of the programming manual, and the programming of the note matters, please fully understand the content of the correct use of this product.

## Explanation

- Only operators with certain electrical knowledge can perform other operations such as wiring the product. If there is any unclear usage, please consult our company's technicians.
- When using this product, please confirm whether it meets the requirements and safety. If this product malfunctions and may cause machine failure or loss, please set up backup and safety functions by yourself.
- The contents described in the manual are subject to specification changes without notice.


## safety matters

- When using under the following conditions and environments, please consult our technical staff and confirm the specifications. At the same time, you must leave room for rated functions and other use and take safety insurance measures into consideration. Control the security measures to a minimum.
- When used outdoors, where there is potential chemical pollution, electrical radiation, and conditions and environments that are not recorded in product samples or instructions.
- Used in nuclear energy control, railways, aviation, vehicle equipment, combustion equipment, medical equipment, safety machinery, administrative agencies and special industries, etc.
- Systems, machinery, devices, etc. that are expected to have a great impact on people and property.
- Used for high-reliability equipment such as gas, water pipes, electricity supply systems and 24 -hour non-stop operation systems.


## Responsibility statement

- Corresponding to the content of this manual, although carefully edited and checked, if you have any questions or find errors, please contact our company.
- The examples listed in the manual and other technical materials are for user understanding and reference only, and actions are not guaranteed.
- Due to changes in specifications and products, the content described in the document is
- for reference only, and the actual product shall prevail. Our company reserves the right of final interpretation.


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### 1.1 PLC host configuration

According to different hardware structure functions, PLC is divided into standard, customized, bus, compact PLC and expansion modules.
The main unit of the programmable controller is the basic unit; in order to expand its input and output points, expansion modules are provided; in addition, special expansion modules for special control can also be connected to meet a variety of applications.

### 1.2 Expansion module composition

PLC provides digital modules, analog modules, temperature modules, weighing modules, functional modules, etc. The expansion module can only be connected to any host of our company. In order to ensure the correct installation and operation of this product, please read the relevant manual carefully before use. A host can match up to 16 extensions.

## 2 Device device function

### 2.1 Introduction of soft components

A device with a certain function inside a PLC, and a soft element is a device with a certain function inside a PLC. These devices are composed of electronic circuits, registers, and memory units. In the ladder diagram, such as buttons, switches, relays (Relay), timers (Timer) and counters (Counter) and so on.
The basic internal devices of PLC are as follows:

| Device name | Description |
| :---: | :---: |
| Input relay <br> (X) | The input relay is the basic unit of internal memory storage corresponding to PLC and external input points (terminals used to connect with external input switches and receive external input signals). It is driven by an input signal sent from the outside, making it 0 or 1 . The state of the input relay cannot be changed by the method of programming, that is, the basic unit corresponding to the input relay cannot be rewritten. The input relay corresponding to no input signal can only be left empty and cannot be used for other purposes. <br> Device representation: $X 0, X 1, \ldots X 7, X 10, X 11 \ldots$, the device symbol is represented by $X$, and the sequence is numbered in octal. <br> The address number of the I/O expansion module: It is the same as the number of the host, and it is numbered in octal based on the last point of the host. |
| Output relay <br> (Y) | The output relay is the basic unit of internal memory storage corresponding to PLC and external output points (used to connect with external loads). It can be driven by input relay contacts, other internal device contacts and its own contacts. It uses a normally open contact to connect to an external load, and other contacts, like input contacts, can be used multiple times without limitation. There is no output relay corresponding to the output, it is empty, if necessary, it can be used as an internal relay. Device representation: $Y 0, Y 1, \ldots Y 7, Y 10, Y 11 \ldots$, the device symbol is represented by $Y$, and the sequence is numbered in octal. <br> The address number of the I/O expansion module: It is the same as the number of the host, and it is numbered in octal based on the last point of the host. |


| Auxiliary relay <br> (M) | - Auxiliary relay is a kind of auxiliary relay inside the PLC. Its function is the same as the auxiliary (middle) relay in the electrical control circuit. Each auxiliary relay also corresponds to a basic unit of the memory. It can be input relay contacts, output relay contacts and other internal The device's contact is driven, and its own contact can also be used for unlimited times. The auxiliary relay cannot directly drive the external output, and it needs to pass the output point to output. <br> - Device representation: MO, M1,...M7, M8..., the device symbol is represented by $M$, and the order is numbered in decimal. |
| :---: | :---: |
| Status relay <br> (S) | The status relay is a stepping action control program input method, and the control program can be written by using the instruction STL to control the transfer of the status relay S. If the step program is not used at all in the program, the status relay $S$ can also be used as an auxiliary relay $M$, or as an alarm point for external fault diagnosis. <br> - Device representation: $S 0, S 1, \ldots . S 1023$, the device symbol is represented by S , and the order is numbered in decimal. |
| Timer <br> ( T ) | - The timer is used to complete the timing control. The timer contains coils, contacts and timing value registers. When the coil is energized and the predetermined time is reached, its contacts will act. The timer's timing value is given by the set value. Each timer has a specified clock cycle (timing unit: $1 \mathrm{~ms} / 10 \mathrm{~ms} / 100 \mathrm{~ms}$ ). <br> - Device representation: $\mathrm{T}, \mathrm{T} 1, \ldots . \mathrm{T} 255$, the device symbol is represented by T , and the order is numbered in decimal. Different number ranges correspond to different clock cycles. |
| Counter <br> (C) | - The counter is used to realize counting operation. To use the counter, the set value of counting (that is, the number of pulses to be counted) should be given in advance. The counter contains coils, contacts and counting memory. When the coil turns from Off $\rightarrow$ On, it is regarded as the counter has a pulse input, and its count value is increased by one. There are 16 -bit and 32 -bit and high-speed counters for users to choose. <br> Device representation: $\mathrm{C} 0, \mathrm{Cl}, \ldots \mathrm{C} 255$, the device symbol is represented by C , and the order is numbered in decimal. |
| Data register <br> (D) | When PLC performs various sequence control and timing value and count value related control, it often needs to do data processing and numerical calculation, and the data register is specially used to store data or various parameters. Each data register has a 16 -bit binary value, that is, a word is stored, and two data registers with adjacent numbers are used to process double words. <br> - Device representation: DO, D1,...D1 1999, the device symbol is represented by C, and the sequence is numbered in decimal. |
| Index register $(E, F)$ | - E, F, and general data registers are 16-bit data registers, which can be written and read freely, and can be used for word devices, bit devices and constants for indirect addressing. <br> Device representation: E0~E7, F0~F7, the device symbol is represented by E, F, and the order is numbered in decimal. |
| Constant $(\mathrm{K}, \mathrm{H})$ | - K represents a decimal integer value, and H represents a hexadecimal value. They are used as the set value and current value of timers and counters, or the operands of application instructions. |

### 2.2 Numerical value, constant ( $\mathrm{K}, \mathrm{H}$ )

| constant | K | Decimal | K-32,768~K32,767 (16-bit operation) <br> K-2,147,483,648 $\sim$ K2,147,483,647 (32-bit operation) |
| :---: | :---: | :---: | :--- |
|  | H | Hexadecimal | H0 $\sim$ HFFFF (16-bit operation) <br> H0 $\sim$ HFFFFFFFF (32-bit operation) |

There are 5 types of numerical values that can be used for PLC numerical values to perform calculation tasks. The tasks and functions of various numerical values are described below.

1. Binary

The numerical calculation or storage in the PLC adopts binary system. The binary value and related terms are as follows:

| Types of | Description |
| :--- | :--- |
| Bit | Bit is the most basic unit of binary value, and its state is either 1 or 0 |
| Nibble | It is composed of 4 consecutive digits (such as b3 $\sim$ b0) which can be used to represent a decimal number 0~9 or <br> hexadecimal 0~F |
| Byte | It is composed of two consecutive nibbles (that is, 8 bits, b7~b0), which can represent 00~FF in hexadecimal |
| word | It is composed of two consecutive bytes (that is, 16 bits, b15 $\sim$ b0), which can represent the hexadecimal 4-digit value <br> $0000 \sim$ FFFF |
| Double word | It is composed of two consecutive words (that is, 32 bits, b31 $\sim$ b0), which can represent 8-bit hexadecimal values <br> $0000000 \sim$ FFFFFFFF |

## 2. Octal

The PLC's external input and output terminal numbers adopt octal coding
Example: External input: X0~X7, X10~X17...(device number)
External output: YO~Y7, Y10~Y17...(device number)

## 3. Decimal

As the setting value of timer T, counter C, etc., for example: TMR CO K50. (K constant)
S, M, T, C, D, E, F, P, I and other device numbers, for example: M10, T30. (Device number)
Used as an operand in application instructions, for example: MOV K123 DO. (K constant)

## 4. $B C D$

A decimal data is represented by half a byte or 4 digits, so consecutive 16 digits can represent 4 -digit decimal numerical data.

## 5. Hexadecimal

Used as an operand in application instructions, for example: MOV HIA2B DO. (H constant)
Constant K :
In the PLC system, the decimal value is usually represented by the word "K" in front of the value.
Example: K 100 , expressed as a decimal system, and its value is 100 .

- When $K$ is used with bit devices $X, Y, M$, and $S$, it can be combined into data in the form of nibble, byte, word or double word.
Example: K2Y10, K4M100. Here K1 represents a combination of 4 bits, and K2~K4 represent combinations of 8,12 and 16 bits respectively.
Constant H :
The hexadecimal value in the PLC is usually represented by the "H" character in front of the value, for example: H 100 , which is expressed in hexadecimal and the value is 100 .

The numerical comparison table is as follows:


### 2.3 Input and output relay (X, Y)

Input and output relays are all numbered in octal
Host number: The number of input and output terminals is fixed from $X 0$ and $Y 0$, and the number of numbers varies with the number of points of the host.

I/O expansion: The number of input and output terminals is calculated according to the connection sequence of the host.

|  | range |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| name | 14 point | 16 point | 24 point | 32 point | 40 point | 48 point | 60 point | 68 point | Expansion 1/0 (Note 1) |
| Input (X) | $\begin{gathered} \hline \text { X0~X7 } \\ (8 \\ \text { point) } \end{gathered}$ | $\begin{gathered} \text { X0~X7 } \\ (8 \\ \text { point) } \end{gathered}$ | $\begin{aligned} & \text { X0~X13 } \\ & (8 \text { point) } \end{aligned}$ | $\begin{gathered} \hline \text { X0~X17 } \\ (16 \\ \text { point) } \end{gathered}$ | $\begin{gathered} \hline \text { X0~X27 } \\ (24 \\ \text { point) } \end{gathered}$ | $\begin{gathered} \hline \text { X0~X27 } \\ (24 \\ \text { point }) \end{gathered}$ | $\begin{gathered} \hline \mathrm{XO} \mathrm{\sim X43} \\ \quad(36 \\ \text { point) } \end{gathered}$ | $\begin{gathered} \hline \text { X0~X43 } \\ \quad(36 \\ \text { point) } \end{gathered}$ | X※~X377 |
| Output (Y) | $\begin{gathered} \hline \text { YO~Y5 } \\ (6 \\ \text { point) } \end{gathered}$ | $\begin{gathered} \hline \text { YO~Y7 } \\ (8 \\ \text { point) } \end{gathered}$ | $\begin{aligned} & \text { Y0~Y } 13 \\ & \text { (8 point) } \end{aligned}$ | $\begin{gathered} \hline \text { YO~Y17 } \\ (16 \\ \text { point }) \end{gathered}$ | $\begin{gathered} \hline \text { YO~Y17 } \\ (16 \\ \text { point) } \end{gathered}$ | $\begin{gathered} \hline \text { YO~Y27 } \\ (24 \\ \text { point }) \end{gathered}$ | $\begin{gathered} \hline \text { YO~Y27 } \\ (24 \\ \text { point }) \end{gathered}$ | $\begin{gathered} \hline \text { YO~Y37 } \\ (32 \\ \text { point }) \end{gathered}$ | Y $\sim$ Y377 |

Description:
Note 1: Expansion I/O input and output starting number starts with the last number of connecting host input/output points. The numbers of the extended I/O are arranged in sequential order. If the last point of the host is $\mathrm{X} \mathrm{n} \square$ (the number range in $\square$ is $0-7)$, the start number of the digital extended input is $X(n+1) 0$. The same is true for the extended output start number. The maximum input number can reach $X 377$, and the maximum output number can reach Y377.

Example: The last point of the host is X 27 , and the start number of the extended input is X 30 . The last point of the host is X43, and the start number of the extended input is X50.

1. Input relay: X0~X377

The number of the input relay (or input terminal) is coded in octal, the maximum number of points can reach 256 points, and the range is as follows: $\mathrm{X} 0 \sim \mathrm{X} 7, \mathrm{X} 10 \sim \mathrm{X17}, \ldots, \mathrm{X} 370 \sim \mathrm{X} 377$.

Function of input contact $X$ :
The input contact $X$ is connected with the input device, and the input signal is read into the PLC. There is no limit to the number of times the $A$ or $B$ contact of each input contact $X$ can be used in the program. The On/Off of the input contact $X$ will only change with the On/Off of the input device. You cannot use the programming software to force the On/Off of the input contact $X$.

## 2. Output relay: YO~Y377

The number of the output relay (or output terminal) is coded in octal, the maximum number of points can reach 256 points, and the range is as follows: YO~Y7, Y10~Y17,..., Y370~Y377.

Function of output contact $Y$ :
The output contact $Y$ sends out On/Off signals to drive the load connected to the output contact $Y$. There are two types of output contacts, one is a relay and the other is a transistor. There is no limit to the number of times that the A or B contact of each output contact $Y$ can be used in the program, but the number of the output coil $Y$ is only recommended in the program. It can be used once, otherwise, according to the PLC's program scanning principle, the power to determine the output state will fall on the last output $Y$ circuit in the program.

## - Input processing

1. The PLC will read the On/Off status of the external input signal into the input image area once before executing the program.
2. If the input signal changes on/off during program execution, the state in the input image area will not change, and the new On/Off state of the input signal will be read until the next scan starts.
3. There is a delay of about 10 ms from the time the external signal $O n \rightarrow$ Off or $O f f \rightarrow$ On changes to the time when the contact in the program is recognized as On/Off (but it may be affected by the program scan cycle).

- Program processing

After the PLC reads the On/Off status of each input signal in the input image area, it starts to execute each instruction in the program sequentially from address 0 , and the processing result, namely the On/Off of each output coil, is also successively stored in each device image area. Inside.

- Output processing

1. When the END instruction is executed, the On/Off status of $Y$ in the device image area is sent to the output image area for latch, and this image area is actually the coil of the output relay.
2. There is about 10 ms delay between the relay coil $O n \rightarrow$ Off or Off $\rightarrow$ On changing to the contact On/Off.
3. Using a transistor module, there will be a delay of about 10~20us from the On $\rightarrow$ Off or $O f f \rightarrow$ On change to the contact On/Off.

### 2.4 Auxiliary relay ( $M$ )

All auxiliary relays are numbered in decimal system, please refer to the corresponding table for the serial number of each series:

| Auxiliary relay (M) | General use | M0~M499, 500 points. Can use parameter settings to change to the power failure retention area | 4096 points in total |
| :---: | :---: | :---: | :---: |
|  | For power failure | M500~M999, M2000~M4095, 2,596 points. Can use parameter settings to change to non-latched area |  |
|  | Special use | M1000~M1999, 1,000 points. Part of it is maintained |  |

Function of auxiliary relay:
Auxiliary relay $M$ and output relay $Y$ have output coils and $A, B$ contacts, and there is no limit to the number of times they can be used in the program. Users can use auxiliary relay $M$ to combine control loops, but they cannot directly drive external loads. According to its nature, it can be divided into the following three types:

1. General auxiliary relay: If the general auxiliary relay encounters a power failure when the PLC is running, its status will all be reset to Off, and its status will remain Off when it is re-powered.
2. Auxiliary relay for power failure retention: If the auxiliary relay for power failure retention encounters a power failure when the PLC is running, its state will all be maintained, and its state will be the state before the power failure when the power is turned on again.
3. Special auxiliary relay: each special auxiliary relay has its specific function, please don't use the undefined special auxiliary relay. Special auxiliary relays cannot be used as ordinary relay $M$.

### 2.5 Status relay (S)

The status relays are all numbered in decimal system, please refer to the corresponding table for the serial number of each series:

| Status <br> relay <br> (S) | Initial | S0 ~ S9, 10 points. Can be modified to be latched by setting up parameters. | Total 1024 points |
| :---: | :---: | :---: | :---: |
|  | Zero return | S10~S19, 10 points, used with IST instruction. Can be modified to be latched by setting up parameters. |  |
|  | General purpose | S20 ~ S499, 480 points. Can be modified to be latched by setting up parameters. |  |
|  | Latched | S500 ~ S899, 400 points. Can be modified to be non-latched by setting up parameters. |  |
|  | Alarm | S900 ~ S1023, 124 points. Can be modified to be latched by setting up parameters. |  |

Function of status relay:
The state relay $S$ can be easily set up in the engineering automation control program. It is the most basic device of the step ladder diagram. STL, RET, etc. must be included in the step ladder diagram (or Sequential Function Chart, SFC) Use with instructions.

The device number of stepping relay $S$ is $S 0 \sim S 1023$ with 1,024 points. Each stepping relay $S$ and output relay $Y$ have output coils and $A, B$ contacts, and there is no limit to the number of times they can be used in the program, but they cannot directly drive external loads. When stepping relay $(S)$ is not used for stepping ladder diagram, it can be used as a general auxiliary relay. Its nature can be divided into the following four types:

| Initial step relay | SO~S9, a total of 10 points. <br> The step point used as the initial state in the Sequential Function Chart (SFC). |
| :---: | :--- |
| Zero return step relay: | S10~S19, 10 points. <br> When the ZL 60 IST instruction is used in the program, S10~S19 are planned for home return. <br> If the IST instruction is not used, it will be used as a general stepping relay. |
| General purpose step <br> relay | S20~S499, 480 points. <br> In the sequence function chart (SFC) as a general purpose step point, if there is a power <br> failure when the PLC is running, its status will be cleared. |
| Latched step relay | S20~S127, 108 points. <br> In the sequence function diagram (SFC), if the stepping relay for power failure retention <br> encounters a power failure when the PLC is running, its state will all be maintained, and its <br> state will be the state before the power failure when the power is retransmitted. |
| Alarm step relay | S900~S1023, 124 points. <br> The step relay for alarm and the alarm point drive command ZL 46 ANS are used as alarm <br> contacts to record relevant warning information and to eliminate external faults. |

### 2.6 Timer (T)

The timers are all numbered in decimal, please refer to the corresponding table for the serial number:

| Timer T | 100ms general purpose | TO ~ T199, 200 points. <br> When M1028 is OFF, T64 to T 126 is 100 ms When M1028 is ON, T64 to T126 is 10 ms | Total 256 points |
| :---: | :---: | :---: | :---: |
|  | 10 ms general purpose | T200~T245, T250~T255, 40 points. <br> When M1038 is OFF, For T200 to T245 and T250 to T255 is 10 ms When M1038 is ON, For T200 to T245 and T250 to T255 is 1ms |  |
|  | 1 ms accumulative | T246 ~ T249, 4 points. |  |

Timer function:
The timer uses $1 \mathrm{~ms}, 10 \mathrm{~ms}$, and 100 ms as a timing unit. The timing method adopts counting up. When the current value of the timer = the set value, the output coil is turned on. The set value is a decimal K value. Data register $D$ can also be used As a set value.

The actual setting time of the timer = timer unit * setting value.
According to its nature, it can be divided into the following three types:

1. General purpose timer:

The timer is generally used to time one time when the TMR instruction is executed. When the TMR instruction is executed, if the timing reaches, the output coil is turned on.


1. When $X 0=O n$, The PV in timer TO will count up by 100 ms . When the PV = SV K100, the output coil TO will be On.
2. When $X 0=$ Off or the power is off, the PV in timer TO will be cleared as 0 , and the output coil TO will be Off.

2, Accumulative type timer:
The timer executes once when the program reaches TMR instruction. When TMR instruction is executed, the output coil will be On when the timing reaches its target.


1. When $\mathrm{XO}=\mathrm{On}$, The PV in timer T246 will count up by 100 ms . When the PV $=$ SV K100, the output coil T 246 will be On.
2. When $\mathrm{XO}=$ Off or the power is off, timer T246 will temporarily stop the timing and the PV remain unchanged. When $X O$ is On again, the timing will resume and the $P V$ will count up and when the $P V=S V$ K100, the output coil T246 will be On.

### 2.7 Counter (C)

All counters are numbered in decimal system, please refer to the corresponding table for the serial number of each series:

| Counter C | 16-bit counting up, for general purpose | $\mathrm{C} 0 \sim \mathrm{C} 99,100$ points. |  |
| :--- | :--- | :--- | :--- |
|  | 16-bit counting up, for latched | $\mathrm{C} 100 \sim \mathrm{C} 199,100$ points. |  |
|  | 32-bit counting <br> up/down, for general purpose | $\mathrm{C} 200 \sim \mathrm{C} 234,20$ points. |  |
| 32-bit counting up/down <br> high-speed counter C | software 1-phase 1 input | Hardware 1-phase 2 inputs | C235~C238, 4 points |
|  | Hardware 2-phase 2 inputs | $\mathrm{C} 241 \sim \mathrm{C} 243,3$ points |  |

1. Features of counter:

| project | 16 bits counters | 32 bits counters |  |
| :---: | :---: | :---: | :---: |
| Type | General purpose | General purpose | High speed |
| Counting direction | Counting up | Counting up, counting down |  |
| Set value | 0~32,767 | -2,147,483,648 ~+2,147,483,647 |  |
| SV designation | Constant K or data register D | Constant K or data register D (designating 2 values) |  |
| Present value | Counting will stop when the SV is reached. | Counter will continue when the SV is reached. |  |
| Output contact | On and being retained when the counting reaches SV. | On and keeps being On when counting up reaches SV. Reset to Off when counting down reaches SV. |  |
| Reset | PV will be return to 0 when RST instruction is executed and the contact will be reset to Off. |  |  |
| Contact action | Acts when the scanning is completed. | Acts when the scanning is completed. | Acts immediately when the counting reaches its target, has nothing to do with the scan period. |

2, Functions of counters

When the pulse input signals of the counter go from Off to On and the present value in the counter equals the set value, the output coil will be On. The set value should be a K value in decimal and the data register D can also be a set value.
16-bit counters C0 ~ C199:

- The setup range of 16 -bit counter: $\mathrm{KO} \sim \mathrm{K} 32,767$. KO is the same as K 1 . The output contact will be On immediately when the first counting starts.
- PV in the general purpose counter will be cleared when the power of the PLC is switched off. If the counter is a latched type, the counter will retain the PV and contact status before the power is off and resume the counting after the power is on again.
- If you use MOV instruction, send a value bigger than the SV to the present value register of CO , next time when X1 goes from Off to On, the contact of counter CO will be On and its PV will equal SV.
- The SV in the counter can be constant K (set up directly) or the values in register D (set up indirectly, excluding special data registers D1000~ D1999).
- If you set up a constant K as the SV , it should be a positive value. Data register D as SV can be positive or negative. When the PV reaches up to 32,767 , the next PV will turn to $-32,768$.
Example:

LD X0

a) When $\mathrm{XO}=\mathrm{On}$, RST instruction will be executed, PV in C0 will be " 0 " and the output contact will be reset to Off.
b) When X1 goes from Off to On, the PV in the counter will count up (plus 1)
c) When the counting of C0 reaches SV K5, the contact of C 0 will be On and PV of $\mathrm{CO}=\mathrm{SV}$ $=K 5$. The X1 trigger signal comes afterwards will not be accepted by C0 and the PV of C0 will stay at K5.

32-bit general purpose addition/subtraction counters C200~C234:

- The setup range of 32-bit counter: K-2,147,483,648 ~ K2,147,483,647.
- Addition or subtraction of the counters is designated by On/Off status of special auxiliary relays M1200 ~ M1234. For example, when M1200 = Off, C200 will be an addition counter; when M1200=On, C200 will be a subtraction counter.
- The SV can be constant K or data register D (excluding special data registers D1000~D1999). Data register $D$ as SV can be a positive or negative value and an SV will occupy two consecutive data registers.
- PV in the general purpose counter will be cleared when the power of the PLC is switched off. If the counter is a latched type, the counter will retain the PV and contact status before the power is off and resume the counting after the power is on again.
- When the PV reaches up to $2,147,483,647$, the next PV will turn to $-2,147,483,648$. When the PV reaches down to $-2,147,483,648$, the next PV will turn to $2,147,483,647$.


1. X10 drives M1200 to determine whether C200 is an addition or subtraction counter.
2. When X11 goes from Off to On, RST instruction will be executed and the PV in C 200 will be cleared to "0" and the contact will be Off.
3. When X 12 goes from Off to On, the PV in the counter will count up (plus 1) or count down (minus 1).

4. When the PV in C200 changes from $\mathrm{K}-6$ to $K-5$, the contact of C 200 will go from Off to On. When the PV in C200 changes from K-5 to $K-6$, the contact of C 200 will go from On to Off.
5. If you use MOV instruction, HPP to send a value bigger than the $S V$ to the present value register of CO , next time when XI goes from Off to On, the contact of counter C0 will be On and its PV will equal SV.

32-bit high-speed addition/subtraction counters C235~C255:

1. The setup range of 32 -bit counter: $\mathrm{K}-2,147,483,648 \sim K 2,147,483,647$
2. Addition or subtraction of C235 ~ C244 is designated by On/Off status of special auxiliary relays M1235 ~ M1244. For example, when M1235 $=$ Off, C235 will be an addition counter; when M1235 $=$ On, C235 will be a subtraction counter.
3. Addition or subtraction of C246 ~ C255 is designated by On/Off status of special auxiliary relays M1246 ~ M1255. For example, when M1246 = Off, C246 will be an addition counter; when M1246=On, C246 will be a subtraction counter.
4. The SV can be constant K or data register D (excluding special data registers D1000~D1999). Data register D as SV can be a positive or negative value and an SV will occupy two consecutive data registers.
5. If using DMOV instruction, HPP to send a value which is large than the setting to any high-speed counter, next time when the input point $X$ of the counter goes from Off to On, this contact will remain unchanged and it will perform addition and subtraction with the present value.
6. When the PV reaches up to $2,147,483,647$, the next PV will turn to $-2,147,483,648$. When the $P V$ reaches down to $-2,147,483,648$, the next PV will turn to $2,147,483,647$.

PLC models support high-speed counters. C235~C240 are program interrupted one-phase high-speed counters, with counting frequencies up to 10 KHz . C241~C254 are hardware high speed counters (Hardware High Speed Counter hereinafter referred to as HHSC). There are four HHSCs (HHSCO~3). The pulse input frequency of HHSCO~3 and HHSCl can reach 200 kHz . among them:

| Numbering | Abbreviation |
| :--- | :--- |
| C251 | HHSC0 |
| C252 | HHSC1 |
| C253 | HHSC2 |
| C254 | HHSC3 |

> Each HHSC can only be assigned to one number at a time. Use the DCNT command as the assignment.
> Each HHSC has three counting modes:
(1) 1 phase 1 input, also known as pulse/direction (Pulse/Direction) mode
(2) 1 phase 2 input, also called forward/reverse (FWD/REV) mode
(3) 2 phase 2 input, also known as AB-phase (AB-phase) mode
(4) Please refer to the table below for distinguishing by serial number.

|  | software high-speed counter |  |  |  |  |  | Hardware high-speed counter |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 phase 1 input |  |  |  |  |  | 1 phase 1 input |  |  |  | 2 phase 2 input |  |  |  |
|  | C235 | C236 | C237 | C238 | C239 | C240 | C241 | C242 | C243 | C244 | C251 | C252 | C253 | C254 |
| X0 | U/D |  |  |  |  |  | U/D |  |  |  | A |  |  |  |
| X1 |  | U/D |  |  |  |  |  |  |  |  | B |  |  |  |
| X2 |  |  | U/D |  |  |  |  | U/D |  |  |  | A |  |  |
| X3 |  |  |  | U/D |  |  |  |  |  |  |  | B |  |  |
| X4 |  |  |  |  | U/D |  |  |  | U/D |  |  |  | A |  |
| X5 |  |  |  |  |  | U/D |  |  |  |  |  |  | B |  |
| X6 |  |  |  |  |  |  |  |  |  | U/D |  |  |  | A |
| X7 |  |  |  |  |  |  |  |  |  |  |  |  |  | B |

Description:

| U | Count up |
| :--- | :--- |
| D | Count down |
| A | Phase A input |
| B | Phase B input |

## 3. Counting modes

The counting modes of the hardware high-speed counters in CPU can be set in D1225 ~ D1228:

| Counting modes |  | Wave pattern |  |
| :---: | :---: | :---: | :---: |
| Type | Set value in special D | Counting up(+1) | Counting down(-1) |
| 1-phase <br> 1 input | 1 <br> (Normal frequency) | U/D <br> U/D FLAG $\qquad$ | $\qquad$ |
|  | 2 <br> (Double frequency) | U/D <br> U/D FLAG $\qquad$ | $\qquad$ |
| 1-phase <br> 2 inputs | 1 <br> (Normal frequency) | $0 \longdiv { \sim }$ |  |
|  | 2 <br> (Double frequency) |  |  |


| Counting modes |  | Wave pattern |
| :---: | :---: | :---: |
| Type | Set value in special D | Counting up（＋1）Counting down（－1） |
| 2－phase <br> 2 inputs | 1 <br> （Normal frequency） |  |
|  | 2 <br> （Double frequency） |  |
|  | 3 <br> （Triple frequency） |  |
|  | 4 <br> （4 times frequency） |  |

4，High－speed counter related flag signals and special registers for related settings：

| Flag | Function |
| :---: | :---: |
| M1235～M1244 | C235～C244 High speed counter counting direction specified． <br> When M12ロロ＝Off，C2口ロ：Count on。 <br> When M12ロロ＝On，C2口■：Count off。 |
| D1225 | The counting mode of the 1st group counters（C251） |
| D1226 | The counting mode of the 2nd group counters（C252） |
| D1227 | The counting mode of the 3rd group counters（C253） |
| D1228 | The counting mode of the 4th group counters（C254） |
| D1225～D1228 | PLC hardware high speed counter HHSC0～HHSC3 counting mode setting，not the following setting values are preset for the double frequency counting mode． <br> 2：for the double frequency counting mode，（factory value）． <br> 3：it is the triple frequency counting mode． <br> 4：it is the quadruple frequency counting mode．（desired value） |

2-phase $A B$ input high-speed counter:


Contact Y0, C251


1. When X 10 is On, RST instruction will be executed. The PV in C251 will be cleared to "0" and the output contact will be reset to be Off.
2. In C251, when X11 is On and C251 receives the A-phase signals from X0 and B-phase signals from XI, the PV in the counter will count up (plus 1) or count down (minus 1). You can select different counting modes.
3. When the counting of C251 reaches SV K 5 , the contact of C 251 will be On. If there are still input signals coming in, the counting will continue.
4, The counting modes (normal frequency, double frequency, triple frequency or 4 times frequency) of C251 (HHSC0) can be set up by D1225.

### 2.8 Numbering and Functions of Registers [D], [E], [F]

A data register is for storing a 16 -bit datum of values between $-32,768$ to $+32,767$. The highest bit is " + " or "-" sign. Two 16 -bit registers can be combined into a 32 -bit register ( $D+1$; $D$ of smaller No. is for lower 16 bits). The highest $b$ it is " + " or "-" sign and it can store a 32-bit datum of values between $2,147,483,648$ to $+2,147,483,647$.

| Data register D | General purpose | JS/T 24 points and below points host | D0 ~ D499, 500 points.Fixed as general purpose. Cannot be modified to be latched by setting up parameters. |
| :---: | :---: | :---: | :---: |
|  |  | JS 32 points and above points host JM host | D0 ~ D199, 200 points. Can be modified to be latched by setting up parameters. |
|  |  | JH2 series host | $\text { D0 ~ D199, D12000~D30000, } 28200$ points. |
|  | Latched | JS/T 24 points and below points host | D500 ~ D499, 500 points |
|  |  | JS 32 points and above points host JM, JH2 host | $\begin{aligned} & \text { D200~D999, D2000~D1 1999, } \\ & \text { 10,800 points. } \end{aligned}$ |
|  | Special purpose | D1000~D1999, 1,000 pints. Some are latched. |  |
|  | Index register E, F | E0~E7, F0~F7, 16 points. |  |

1. Registers can be divided into the following four types according to their nature:

- General purpose register: When the PLC is powered off, the value data in the register will be cleared to 0.
- Latched register: When the power of PLC is switched off, the data in the register will not be cleared but will retain at the value before the power is off. You can use RST or ZRST instruction to clear the data in the latched register.
- Special purpose register: Every register of this kind has its special definition and purpose, mainly for storing the system status, error messages and monitored status.
- Index register E, F: The index register is a 16-bit register, E0~E7, FO~F7 total 16 points.

2, Index Register [E], [F]


Higher 16 bits Lower 16 bits


> Index registers E, F are 16-bit data registers and can be written and read.
> If you need to use a 32-bit register, you have to designate
> E. In this case, F will be covered by E and cannot be used anymore; otherwise, the content in E (32-bit) will be incorrect. We suggest you use DMOVP K0 E instruction, the content in E (including F) will be cleared to "0" when the power of PLC is switched on.
> The combination of E, F when you use a 32-bit index register:
> (FO, E0), (F1, E1), (F2, E2), $\cdots$ (F7, E7)

- The index register is the same as normal operands, can be used for moving or comparison on word devices (KnX, KnY, KnM, KnS, T, C, D) and bit devices (X, Y, M, S). supports constant (K, H) index register. has 16 points of index registers EO ~ E7, FO ~F7.
- When you use the instruction mode to generate constant $(\mathrm{K}, \mathrm{H})$ index register function, please use symbol "@". For example, "MOV K10@EO DOFO" .


### 2.9 Pointer [N], Pointer [P], Interruption Pointer [l]

| Pointer | N | Master control loop |  | NO~N7, 8 points | Control point of master control loop |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | P | For CJ, CALL instructions |  | P0~P255, 256 points | Position pointer of CJ, CALL |
|  | 1 | Interruption | External interruption | ```IO0\square(X0), I10\square(X1), I20\square(X2), I30\square(X3) I40\square (X4), 150\square(X5), 160\square(X6), 170\square(X7), 190\square (X10), 191\square(X11), 192\square(X12) , 193\square(X13), 194\square(X14), 195\square(X15), 196\square(X16) , 197\square(X17), 16 point(\square=1, rising- edge trigger, }\square=0\mathrm{ , falling-edge trigger)``` |  |
|  |  |  | Timed interruption | 16 $\square \square$, 17 $\square \square$, 2 points ( $\square \square=02 \sim 99$, time base $=1 \mathrm{~ms}$ ) $18 \square \square$, 1 points ( $\square \square=05 \sim 99$, time base $=0.1 \mathrm{~ms}$ ) | of interruption subroutine |
|  |  |  | High-speed counter interruption | 1010, 1020, 1030, 1040, 1050, 1060, 6 points |  |
|  |  |  | Pulse interruption | 1110, 1120, 1130, 1140, 4 points |  |
|  |  |  | Communication interruption | 1150, 1160, 1170, 3 points |  |

## Description:

1. Input point $X$ as a high-speed counter cannot be used as an external interruption signal. For example, if C251 occupies X0, X1, X2 and X3, the external input interruption No. $100 \square(X 0), I 10 \square(X 1), I 20 \square(X 2)$, and $130 \square(X 3)$ cannot be used.
2. If an interrupt subroutine is executed, the next interrupt subroutine will not be executed until the execution of the interrupt is complete.
3. The time it takes for an interrupt subroutine in a PLC to be executed affects the efficiency of the PLC. It is suggested that the size of an interrupt subroutine not be large.
4, Pointer $N$ : Used with MC and MCR instructions. MC is the master control start instruction. When MC instruction is executer, the instructions between MC and MCR will still be executed normally.
5, Pointer P: Used with application commands ZL 00 CJ, ZL 01 CALL, ZL 02 SRET.

- CJ Conditional Jump:


1. When $\mathrm{XO}=\mathrm{On}$, the program will jump from address 0 to N (designated label P1) and keep on the execution. The addresses in the middle will be ignored.
2. When $\mathrm{XO}=\mathrm{Off}$, the program will execute from address 0 and keep on executing. At this time, CJ instruction will not be executed.

- CALL Call Subroutine, SRET Subroutine Return:

instruction list

| 000000 | LD | X0 |
| :--- | :--- | :--- |
| 000001 | CALL | P2 |
| 000004 | LD | X1 |
| 000005 | OUT | Y1 |
| 000006 | FEND |  |
| 000007 | P2 |  |
| 000008 | LD | M0 |
| 000009 | SET | Y0 |
| 000010 | LD | M1 |
| 000011 | SET | Y2 |
| 000012 | SRET |  |
| 000013 | END |  |

1. When $\mathrm{XO}=\mathrm{On}$, CALL instruction will be executed and the program will jump to P 2 and executed the designated subroutine. When SRET instruction is executed, the program will return to address 24 and keep on the execution.
2, There is no need to edit the FEND and SRET codes in the ladder diagram. After the compilation is passed, the instruction list will be automatically generated.

Interruption Pointer I: Used with application commands ZL 04 EI, ZL 05DI, and ZL 03 IRET, the purpose can be divided into the following six types. The interrupt insertion action must be combined with commands such as EI interrupt insertion enable, DI interrupt insertion prohibition, and IRET interrupt insertion return.

1. External interruption: Due to the special hardware design inside the CPU, the input signals coming in at input terminals X0 ~ X5 when rising-edge or falling-edge triggers will not be affected by the scan cycle. The currently executed program will be interrupted immediately and the execution will jump to the designated interruption subroutine pointer $100 \square(X 0)$, $110 \square$ (X1), I20■(X2), I30 $\square(X 3), I 40 \square(X 4), I 50 \square(X 5)$. Till the execution reaches IRET instruction, the program will return to the original position and keep on its execution.
Example: X2 rising edge interrupt.


instruction list

| 11 | [1201] 12 | Instruction List $\mathbf{x}$ |
| :--- | :--- | :--- |
| 000000 | EI |  |
| 000001 | LD | MQ |
| 000002 | OUT | Y0 |
| 000003 | FEND |  |
| 000004 | I201 |  |
| 000005 | LD | M1000 |
| 000006 | INC | D0 |
| 000009 | IRET |  |
| 000010 | END |  |

2. Timed interruption: PLC automatically interrupts the currently executed program every a fixed period of time and jumps to the execution of a designated interruption subroutine.
Example: 1602 timer interrupt.



| 11 | [l602] 12 | Instruction List $\boldsymbol{x}$ |  |
| :--- | :--- | :--- | :--- |
| 000000 | EI |  |  |
| 000001 | LDP | M1013 |  |
| 000004 | M0V | D0 | D10 |
| 000009 | MOV | D2 | D12 |
| 000014 | MOV | D4 | D14 |
| 000019 | MOV | D6 | D16 |
| 000024 | ZRST | D0 | D7 |
| 000029 | FEND |  |  |
| 000030 | I602 |  |  |
| 000031 | LD | M1000 |  |
| 000032 | DINC | D0 |  |
| 000037 | IRET |  |  |
| 000038 | END |  |  |

3, Interruption when the counting reaches the target: The high-speed counter comparison instruction ZL 53 DHSCS can designates that when the comparison reaches the target, the currently executed program will be interrupted and jump to the designated interruption subroutine executing the interruption pointers $1010,1020,1030,1040,1050$ and 1060.
4, Pulse interruption: The pulse output instruction ZL 57 PLSY can be set up that the interruption signal is sent out synchronously when the first pulse is sent out by enabling flags M1342 and M1343. The corresponding interruptions are 1130 and 1140 . You can also set up that the interruption signal is sent out after the last pulse is sent out by enabling flags M1340 and M1341. The corresponding interruptions are I110 and I120.

## 2. 10 Functions of Special Auxiliary Relays and Special Registers

The types and functions of special auxiliary relay (special $M$ ) and special data register (special D) are as follows. In the following tables, there is a "*" mark in the upper right corner of the number. You can refer to the function description in the next section. If the attribute column is marked as "R", it means that it can only be read. If it is marked as "R/W" means that it can be read and written. If it is marked as "-", it means no change. Marked as "\#", it means that the system will be set according to the PLC status, and the user can read the setting value and compare the description of the manual to further understand the system information.

| Special <br> M | Function Description |
| :--- | :--- | M1000 | Normally open contact (a contact). This contact is ON when running and it is ON when the status |
| :--- |
| is set to RUN. |$\quad$| Normally OFF contact (b contact). This contact is OFF in running and it is OFF when the status is |
| :--- |
| set to RUN. |


| M1029 | Y0, Y1 pulse sending completion flag |
| :--- | :--- |
| M1030 | Y2, Y3 pulse sending completion flag |
| M1031 | Clear all latched memory |
| M1032 | Clear all latched memory |
| M1033 | Memory latched at STOP |
| M1036 | Y4, Y5 pulse sending completion flag |
| M1037 | Y6, Y7 pulse sending completion flag |
| M1038 | Switching T200 T245 ,T250~ T255 timer resolution (10ms/1ms). ON = 1ms. |
| M1039 | Constant scan mode |
| M1040 | Step transition starts |
| M1041 | Step transition starts |
| M1044 | Zero point condition |
| M1045 | All outputs clear inhibit |
| M1046 | STL state setting (On) |
| M1047 | STL monitor enable |
| M1048 | Flag for alarm point state |
| M1049 | Monitor flag for alarm point |
| M1050 | I001 masked |
| M1051 | I101 masked |
| M1052 | I201 masked |
| M1053 | I301 masked |
| M1054 | I401 masked |
| M1055 | I501 masked |
| M1056 | I6 $\quad$ masked |
| M1057 | I7 |
| M1059 | I010~ I060 masked |
| M1060 | System error message 1 |
| M1061 | System error message 2 |
| M1062 | System error message 3 |
| M1063 | System error message 4 |
| M1064 | Operator error |
| M1065 | Syntax error |
| M1066 | Program error |
| M1067 | Program execution error |
| M1068 | Execution error locked (D1068) |
| M1070 | PWM command Y0 output frequency unit switching, 100us when On, 1ms when Off |
| M1071 | PWM command Y0 output frequency unit switching, 100us when On, 1ms when Off |
| M1076 | Battery voltage is too low or malfunction |
| M1077 | Battery voltage is too low or malfunction |
| M1080 | COM2 Monitor request |
| M1081 | FLT command change direction flag |
| M1082 | Flag changed for RTC |
| M1102 | Y10, Y11 pulse sending completion flag |
| M1103 | Y12, Y13 pulse sending completion flag |


| M1104 | Y14, Y15 pulse sending completion flag |
| :--- | :--- |
| M1105 | Y16, Y17 pulse sending completion flag |
| M1106 | Y20, Y21 pulse sending completion flag |
| M1107 | Y22, Y23 pulse sending completion flag |
| M1108 | Y24, Y25 pulse sending completion flag |
| M1109 | Y26, Y27 pulse sending completion flag |
| M1110 | Y30, Y31 pulse sending completion flag |
| M1111 | Y32, Y33 pulse sending completion flag |
| M1112 | Y34, Y35 pulse sending completion flag |
| M1113 | Y36, Y37 pulse sending completion flag |
| M1114 | Y40, Y41 pulse sending completion flag |
| M1115 | Y42, Y43 pulse sending completion flag |
| M1116 | Y44, Y45 pulse sending completion flag |
| M1117 | Y46, Y47 pulse sending completion flag |
| M1118 | Y50, Y51 pulse sending completion flag |
| M1119 | Y52, Y53 pulse sending completion flag |
| M1120 | COM2(RS-485) Communication protocol holding (communication protocol will be the original <br> setting even if D1120 is changed) |
| M1121 | RS-485 Communication data transmission ready |
| M1122 | Sending request |
| M1123 | Receiving completed |
| M1124 | Receiving wait |
| M1125 | Communication reset |
| M1127 | MODRD/RDST/MODRW instructions. Data receiving is completed, RS instruction doesn include. |
| M1128 | Sending / Receiving |
| M1129 | Receiving time out |
| M1131 | MODRD/RDST/MODRW, M1131=On when data convert to HEX |
| M1132 | When it is On, it denotes that there is no relative communication instruction in PLC program. |
| M1138 | COM1 (RS-232) communication protocol holding. D1036 modification invalid after setting. |
| M1143 | When SLAVE mode, ASCII/RTU selection for COM1 (RS-232). OFF for ASCII mode and ON for <br> (used with MODRD/ MODWR/MODRW instructions) (it is Off when in ASCII mode and it is On <br> when in RTU mode) <br> RTU mode. <br> M1144 Ouput start switch of accel/decel pulse output function of adjustable slope |
| M1145 | Acceleration flag of accel/decel pulse output function of adjustable slope |
| M1146 | Target attained frequency flag of accel/decel pulse output function of adjustable slope |
| M1147 | Deceleration flag of accel/decel pulse output function of adjustable slope |
| M1148 | Complete function flag of accel/decel pulse output function of adjustable slope |
| M1149 | Stop counting temporality flag of accel/decel pulse output function of adjustable slope |
| M1140 | MODRD/MODWR/MODRW data received error |
| M1141 | MODRD/MODWR/MODRW command error |
| M1142 | VFD-A command data received error |
|  | 1.COM2(RS-485) ASCII/RTU mode selections when PLC is SLAVE (it is Off when in ASCII mode |


| M1150 | Declare DHSZ command used for multi-group settings comparison mode |
| :--- | :--- |
| M1151 | Finish executing multi-group settings comparison mode |
| M1152 | Declare DHSZ command used to be frequency control mode |
| M1153 | Finish executing frequency control mode |
| M1154 | Start designated deceleration function flag of accel/decel pulse output function of adjustable <br> slope |
| M1161 | 8/16 bits mode (it is On when in 8 bits mode) |
| M1162 | Using flag for the integral of decimal system and the floating point of binary systems. ON for the <br> floating point of binary. |
| M1167 | HKY input is 16 bits mode |
| M1168 | SMOV working mode indication |
| M1169 | PWD mode selection |
| M1170 | Start executing single step |
| M1171 | Execute single step |
| M1172 | 2-phase pulse output switch (on is start) |
| M1173 | On is continuous output switch |
| M1174 | Output pulse number attained flag |
| M1178 | VR0 potentiometer starts |
| M1179 | VR1 potentiometer starts |
| M1184 | Startup MODEM |
| M1185 | Start to initiate MODEM |
| M1186 | Fail to initiate MODEM |
| M1187 | Finish initiating MODEM |
| M1188 | Display if current MODEM is on-line or not |
| M1200 | C200 counting mode (on: count down |
| M1201 | C202 counting mode (on: count down) |
| M1202 | C202 counting mode (on: count down) |
| M1203 | C203 counting mode (on: count down) |
| M1204 | C204 counting mode (on: count down) |
| M1205 | Y54, Y55 pulse sending completion flag |
| M1206 | Y52, Y53 pulse sending completion flag |
| M1207 | PLSV instruction acceleration and deceleration time enable |
| M1226 | C226 counting mode (on: count down) |
| M1227 | C227 counting mode (on: count down) |
| M1228 | C228 counting mode (on: count down) |
| M1229 | C229 counting mode (on: count down) |
| M1230 | C230 counting mode (on: count down) |
| M1231 | C231 counting mode (on: count down) |
| M1232 | C232 counting mode (on: count down) |
| M1233 | C233 counting mode (on: count down) |
| M1234 | C234 counting mode (on: count down) |
| M1235 | C235 counting mode (on: count down) |
| M1236 | C236 counting mode (on: count down) |
| M1237 | C237 counting mode (on: count down) |


| M1238 | C238 counting mode (on: count down) |
| :--- | :--- |
| M1239 | C239 counter mode setting (on: count down) |
| M1240 | C240 counter mode setting (on: count down) |
| M1241 | C241 counter mode setting (on: count down) |
| M1242 | C242 counter mode setting (on: count down) |
| M1243 | C243 counter mode setting (on: count down) |
| M1244 | C244 counter mode setting (on: count down) |
| M1245 | C245 counter mode setting (on: count down) |
| M1246 | C246 counter monitor (on: count down) |
| M1247 | C247 counter monitor (on: count down) |
| M1248 | C248 counter monitor (on: count down) |
| M1249 | C249 counter monitor (on: count down) |
| M1250 | C250 counter monitor (on: count down) |
| M1251 | C251 counter monitor (on: count down) |
| M1252 | C252 counter monitor (on: count down) |
| M1253 | C253 counter monitor (on: count down) |
| M1254 | C254 counter monitor (on: count down) |
| M1256 | System used |
| M1258 | Swap Y0 and Y1 pulse output signal |
| M1259 | Swap Y2 and Y3 pulse output signal |
| M1260 | Let X5 be the reset input signal of all high-speed counter |
| M1261 | DHSCR command High-speed comparison flag |
| M1264 | HHSCO Reset function enable |
| M1265 | HHSCO Start function enable |
| M1266 | HHSC1 Reset function enable |
| M1267 | HHSC1 Start function enable |
| M1268 | HHSC2 Reset function enable |
| M1269 | HHSC2 Start function enable |
| M1270 | HHSC3 Reset function enable |
| M1271 | HHSC3 Start function enable |
| M1272 | HHSC0 Reset control |
| M1273 | HHSC0 Start control |
| M1274 | HHSC1 Reset control |
| M1275 | HHSC1 Start control |
| M1276 | HHSC2 Reset control |
| M1277 | HHSC2 Start control |
| M1278 | HHSC3 Reset control |
| M1279 | HHSC3 Start control |
| M1280 | l00 masked |
| M1281 | I10 masked |
| M1282 | I20 masked |
| M1283 | I30 masked |
| M1284 | I40 masked |
| M1285 | I50 masked |
|  |  |


| M1286 | I6 masked |
| :--- | :--- |
| M1287 | I7 masked |
| M1288 | I8 masked |
| M1289 | IO10 masked |
| M1290 | IO20 masked |
| M1291 | IO30 masked |
| M1292 | I040 masked |
| M1293 | I050 masked |
| M1294 | IO60 masked |
| M1295 | I110 masked |
| M1296 | I120 masked |
| M1297 | I130 masked |
| M1298 | I140 masked |
| M1299 | I150 masked |
| M1300 | I160 masked |
| M1301 | I170 masked |
| M1302 | I180 masked |
| M1303 | Swap high and low byte |
| M1308 | Y0, Y1 emergency stop without deceleration |
| M1309 | Y2, Y3 emergency stop without deceleration |
| M1310 | Y4, Y5 emergency stop without deceleration |
| M1311 | Y6, Y7 emergency stop without deceleration |
| M1312 | Y10, Y11 emergency stop without deceleration |
| M1313 | Y12, Y13 emergency stop without deceleration |
| M1314 | Y14, Y15 emergency stop without deceleration |
| M1315 | Y16, Y17 emergency stop without deceleration |
| M1316 | Y20, Y21 emergency stop without deceleration |
| M1317 | Y22, Y23 emergency stop without deceleration |
| M1318 | Y24, Y25 emergency stop without deceleration |
| M1319 | Y26, Y27 emergency stop without deceleration |
| M1320 | Y30, Y31 emergency stop without deceleration |
| M1321 | Y32, Y33emergency stop without deceleration |
| M1322 | Y34, Y35 emergency stop without deceleration |
| M1323 | Y36, Y37 emergency stop without deceleration |
| M1324 | Y40, Y41 emergency stop without deceleration |
| M1325 | Y42, Y43 emergency stop without deceleration |
| M1326 | Y44, Y45 emergency stop without deceleration |
| M1327 | Y46, Y47 emergency stop without deceleration |
| M1328 | Y50, Y51 emergency stop without deceleration |
| M1329 | Y52, Y53 emergency stop without deceleration |
| M1330 | Y54, Y55 emergency stop without deceleration |
| M1331 | C239 Start/Reset function enable |
| M1332 | C239 Start/Reset function enable |
| M1333 | C240 Start/Reset function enable |


| M1340 | After the CH0 (Y0, Y1) pulse is sent out, an interrupt I110 is generated |
| :--- | :--- |
| M1341 | After the CH0 (Y0, Y1) pulse is sent out, an interrupt I110 is generated |
| M1342 | CH0 (Y0, Y1) pulse is sent out at the same time, interrupt I130 is generated |
| M1343 | CH0 (Y0, Y1) pulse is sent out at the same time, interrupt I130 is generated |
| M1344 | Y0, Y1 pulse sending flag |
| M1345 | Y2, Y3 pulse sending flag |
| M1346 | Y4, Y5 pulse sending flag |
| M1347 | Y6, Y7 pulse sending flag |
| M1348 | Y10, Y11 pulse sending flag |
| M1349 | Y12, Y13 pulse sending flag |
| M1350 | Y14, Y15 pulse sending flag |
| M1351 | Y16, Y17 pulse sending flag |
| M1352 | Y20, Y21 pulse sending flag |
| M1353 | Y22, Y23 pulse sending flag |
| M1354 | Y24, Y25 pulse sending flag |
| M1355 | Y26, Y27 pulse sending flag |
| M1356 | Y30, Y31 pulse sending flag |
| M1357 | Y32, Y33 pulse sending flag |
| M1358 | Y34, Y35 pulse sending flag |
| M1359 | Y36, Y37 pulse sending flag |
| M1360 | Y40, Y41 pulse sending flag |
| M1361 | Y42, Y43 pulse sending flag |
| M1362 | Y44, Y45 pulse sending flag |
| M1363 | Y46, Y47 pulse sending flag |
| M1364 | Y50, Y51 pulse sending flag |
| M1365 | Y52, Y53 pulse sending flag |
| M1366 | Y54, Y55 pulse sending flag |
| M1367 | Y56, Y57 pulse sending flag |
| M1415 | Indicating reading from Salve ID\#8 is completed |


| Special D | Function Description |
| :---: | :---: |
| D1000 | Watchdog timer (WDT) value (Unit: 1ms) |
| D1002 | Program capacity |
| D1003 | Sum of program memory (sum of the PLC internal program memory. User can identify the content of PLC control program by this register) |
| D1004 | Check code for grammar |
| D1005 | System use |
| D1008 | STEP address when WDT timer is ON |
| D1009 | The occuring history of LV signals will be stored in D1009 |
| D1010 | Present scan time (Unit: 0.1 ms ) |
| D1011 | Minimum scan time (Unit: 0.1ms) |
| D1012 | Maximum scan time (Unit: 0.1ms) |
| D1015 | 0~32,767(unit: 0.1 ms ) addition type of high-speed connection timer |
| D1018 | Geocities PI (Low byte) |
| D1019 | Geocities Pl(High byte) |
| D1020 | X0 ~ X7 input filter (Unit: ms) |
| D1021 | ES/EH/EH2/SV: X10 ~ X17 input filter (Unit: ms)SC: X10 ~ X17 input filter (time base: scan cycle), range: 0 ~ 1,000 (Unit: times) |
| D1024 | System use flag |
| D1025 | Communication error code |
| D1028 | Index register E0 |
| D1029 | Index register F0 |
| D1034 | Frequency measurement card working mode |
| D1035 | Set the number of X input point of RUN/STOP |
| D1036 | COM1 (RS-485) Communication protocol |
| D1037 | HKY key repeat time (ms) |
| D1038 | When PLC MPU is slave, the setting of data response delay time. The setting range is $0 \sim 10,000$, and the time unit is 0.1 ms . |
| D1039 | Constant scan time (ms) |
| D1040 | On state number 1 of STEP point S |
| D1041 | On state number 2 of STEP point S |
| D1042 | On state number 3 of STEP point S |
| D1043 | On state number 4 of STEP point S |
| D1044 | On state number 5 of STEP point S |
| D1045 | On state number 6 of STEP point S |
| D1046 | On state number 7 of STEP point S |
| D1047 | On state number 8 of STEP point S |
| D1049 | On number of alarm point |
| D1050 | Modbus communication data conversion. PLC will automatically convert the ASCII data saved in D1070~D1085 to HEX. |
| D1051 | Modbus communication data conversion. PLC will automatically convert the ASCII data saved in D1070~D1085 to HEX. |
| D1052 | Modbus communication data conversion. PLC will automatically convert the ASCII data saved in D1070~D1085 to HEX. |
| D1053 | Modbus communication data conversion. PLC will automatically convert the ASCII data saved in D1070~D1085 to HEX. |
| D1054 | Modbus communication data conversion. PLC will automatically convert the ASCII data saved in D1070~D1085 to HEX. |


| D1055 | Modbus communication data conversion. PLC will automatically convert the ASCII data saved in D1070~D1085 to HEX. |
| :---: | :---: |
| D1056 | Present value of SX/EX MPU analog input channel 0 (CH0) and EH MPU AD card channel 0 (CHO) |
| D1057 | Present value of SX/EX MPU analog input channel 1 (CH1) and EH MPU AD card channel 1 (CH1) |
| D1058 | Present value of EX MPU analog input channel 2 (CH2) |
| D1059 | Present value of EX MPU analog input channel 3 (CH3) |
| D1061 | System error message: Record of the number of errors in the power failure retention area |
| D1064 | System used flag |
| D1065 | System used flag |
| D1066 | Algorithm error code |
| D1067 | Algorithm error code |
| D1068 | Lock the algorithm error address |
| D1069 | M1065~M1067 error address |
| D1070 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1071 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1072 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1073 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1074 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1075 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1076 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1077 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1078 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1079 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1080 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |


| D1081 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| :---: | :---: |
| D1082 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1083 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1084 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1085 | When the PLC built-in RS-485 communication command receives feedback signals from receiver. The signals will be saved in the registers D1070~D1085. User can use the contents saved in the registers to check the feedback data. |
| D1089 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1090 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1091 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1092 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1093 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1094 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1095 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1096 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1097 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1098 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1099 | When the PLC built-in RS-485 communication command is executed, the transmitting signals will be stored in the registers D1089~D1099. User can use the contents saved in registers to check the feedback data. |
| D1100 | Corresponding state after LV signal action |
| D1101 | File register start address |


| D1102 | File register copy number |
| :---: | :---: |
| D1103 | Set as the starting D number of the file register (must be greater than 2000) |
| D1104 | Acceleration and deceleration pulse output use control register (D) start number |
| D1119 | System use |
| D1120 | COM2(RS-485) Communication protocol |
| D1121 | PLC communication address (the address that save PLC communication address, it is latched) |
| D1122 | Residual words of transmitting data |
| D1123 | Residual words of receiving data |
| D1124 | Start character definition (STX) |
| D1125 | First ending character definition (EXT1) |
| D1126 | Second ending character definition (EXT2) |
| D1129 | RS-485 time-out setting (ms) |
| D1130 | MODBUS return error code record |
| D1133 | System program version |
| D1135 | System hardware version |
| D1137 | Address of operator error occurs |
| D1139 | Connection number of BCD module expansion unit (the maximum is two units) |
| D1141 | System used |
| D1144 | Adjustable slope acceleration and deceleration pulse output YO control register (D) start number |
| D1145 | Number of left-side high-speed special extension modules; maximum 8 modules extendable (only supports (for SV) |
| D1146 | Connection number of DISP module expansion unit |
| D1147 | Memory card type (Memory state) b0=0:NO CARD (H0000) b0=1:FLASH CARD exists $\mathrm{b} 8=0$ :The switch of FLASH CARD is Off (H0001) b8=1:The switch of FLASH CARD is On (H0101) |
| D1148 | System use flag signal |
| D1149 | Function card type: 0:NO CARD 1:RS-232(DVP-F232),DU01(DVPDU01) 2:RS-422(DVPF422) 3:COM3(DVP-F232S,DVP-F485S) 4:Potentiometer switch(DVP-F6VR) 5:DIP switch(DVP-F8ID) 6:Transitor output card(DVP-F2OT) 7:Digital input(DVP-F4IP) 8:2AD card (DVP-F2AD) 9:2DA card(DVP-F2DA) |
| D1150 | DHSZ instruction multiple set value comparison mode table count register |
| D1151 | DHSZ instruction frequency control mode table count register |
| D1152 | DHSZ D value changed High word |
| D1153 | DHSZ D value changed low word |
| D1154 | Adjustable slope acceleration and deceleration pulse Y0 output function deceleration interval time ( $10 \sim 32,767 \mathrm{~ms}$ ) recommended value |
| D1155 | Adjustable slope acceleration and deceleration pulse YO output function deceleration interval time (-1~32,700 HZ) recommended value |
| D1156 | Special D specified by RTMU instruction (No. K0~K9) |
| D1157 | Special D that indicated by RTMU command (K0~K9) |
| D1158 | Special D that indicated by RTMU command (KO~K9) |
| D1159 | Special D that indicated by RTMU command (KO~K9) |
| D1160 | Special D that indicated by RTMU command (K0~K9) |
| D1161 | Special D that indicated by RTMU command (K0~K9) |
| D1162 | Special D that indicated by RTMU command (K0~K9) |
| D1163 | Special D that indicated by RTMU command (K0~K9) |
| D1164 | Special D that indicated by RTMU command (K0~K9) |


| D1165 | Special D that indicated by RTMU command (K0~K9) |
| :--- | :--- |
| D1166 | Rising-edge or falling-edge switch of X10 (SCV1.4 or over models) |
| D1168 | RS instruction, interrupt request when receiving specified data character (I150) |
| D1169 | RS instruction, interrupt request when receiving specified data legnth (I160) |
| D1170 | Single step (Single step) PC value during execution |
| D1172 | Two-phase pulse output frequency (12Hz~20KHz) |
| D1173 | Two-phase pulse output mode selection (K1 and K2) |
| D1174 | Target number of two-phase output pulse (low word) |
| D1175 | Target number of two-phase output pulse (high word) |
| D1176 | Current output number of two-phase pulse (low word) |
| D1177 | Current output number of two-phase pulse (high word) |
| D1178 | VR0 value |
| D1179 | VR1 value |
| D1180 | Interrupt I401 to capture high-speed count value(low word) |
| D1181 | Interrupt I401 to capture high-speed count value(high word) |
| D1182 | Pointer register E1 |
| D1183 | Pointer register F1 |
| D1184 | Pointer register E2 |
| D1185 | Pointer register F2 |
| D1186 | Pointer register E3 |
| D1187 | Pointer register F3 |
| D1188 | Pointer register E4 |
| D1189 | Pointer register F4 |
| D1190 | Pointer register E5 |
| D1191 | Pointer register F5 |
| D1192 | Pointer register E6 |
| D1193 | Pointer register F6 |
| D1194 | Pointer register E7 |
| D1195 | Pointer register F7 |
| D1196 | Number system stting (provided for DVP-SX series only) |
| D1197 | System use flag signal |
| D1198 | Interrupt I501 to capture high-speed count value(low word) |
| D1199 | Interrupt I501 to capture high-speed count value(high word) |
| D1200 | M0~M999 auxiliary relay start address when power failure |
| D1201 | M0~M999 auxiliary relay power failure retention end address |
| D1202 | M2000~M4095 auxiliary relay power failure start address |
| D1203 | M2000~M4095 auxiliary relay power failure retention end address |
| D1204 | T0~T199, the start address of the 100ms timer to keep the power off |
| D1205 | T0~T199, 100ms timer power failure retention end address |
| D1206 | T200~T239, the start address of the 10ms timer to keep the power off |
| D1207 | T200~T239, 10ms timer power failure retention end address |
| D1208 | C0~C199, 16-bit counter power failure retention start address |
| D1209 | C0~C199, 16-bit counter power failure retention end address |
| D1210 | C200~C234, 32-bit counter start address when power failure |
| D1211 | TCP/IP Port |
| D1212 | IP0.1 |
| D1213 | IP2.3 |
| D1214 | S0~S899 step point power failure retention start address |
| D1215 | S0~S899 step point power failure retention end address |


| D1216 | D0~D999 register power failure retention start address |
| :--- | :--- |
| D1217 | D0~D999 register power failure retention end address |
| D1218 | D2000~D9999 register start address for power failure retention |
| D1219 | D2000~D9999 register power failure retention end address |
| D1220 | mask0.1 |
| D1221 | mask2.3 |
| D1222 | gateway0.1 |
| D1223 | gateway2.3 |
| D1225 | The first group of counter (HHSC0) counting mode setting |
| D1226 | The second group counter (HHSC1) counting mode setting |
| D1227 | The third group counter (HHSC2) counting mode setting |
| D1228 | The fourth group of counter (HHSC3) counting mode setting |
| D1256 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1257 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1258 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1259 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1260 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1261 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1262 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1273 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1274 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| D1271 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved <br> in D1256~D1295. User can check according to the content of these registers. |
| in D1256~D1295. User can check according to the content of these registers. |  |


| D1275 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| :---: | :---: |
| D1276 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1277 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1278 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1279 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1280 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1281 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1282 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1283 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1284 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1285 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1286 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1287 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1288 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1289 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1290 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1291 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1292 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1293 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1294 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1295 | MODRW command of RS-485 is built-in. The characters that sent during executing is saved in D1256~D1295. User can check according to the content of these registers. |
| D1296 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1297 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1298 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1299 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |


| D1300 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| :---: | :---: |
| D1301 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1302 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1303 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1304 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1305 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1306 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1307 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1308 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1309 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1310 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1311 | MODRW command of RS-485 is built-in. PLC system will convert ASCII in the content of the register that user indicates to HEX and save it in D1296 ?D1311. |
| D1313 | Real time clock (RTC) second 00~59 |
| D1314 | Real time clock (RTC) minute 00~59 |
| D1315 | Real time clock (RTC) hour 00~23 |
| D1316 | Real time clock (RTC) day 01~31 |
| D1317 | Real time clock (RTC) month 01~12 |
| D1318 | Real time clock (RTC) week 1~7 |
| D1319 | Real time clock (RTC) year 00-99 |
| D1340 | The 1st step start frequency and the last step end frequency of CH 0 pulse |
| D1341 | Maximum output frequency (Low word) (it is fixed to 200KHz) |
| D1342 | Maximum output frequency (High word) (it is fixed to 200 KHz ) |
| D1343 | Acceleration /Deceleration time of CH0 pulse |
| D1352 | SC: the 1st step starting frequency and last step ending frequency of Y11.EH: the 2nd step starting frequency and last step ending frequency of $\mathrm{CH} 1(\mathrm{Y} 2, \mathrm{Y} 3)$ output |
| D1353 | SC: Y11 acceleration/deceleration time setting.EH: acceleration/deceleration time setting of 2nd group CH1 (Y2, Y3) output. |
| D1355 | Starting reference for Master to read from Salve ID\#1 |
| D1356 | Starting reference for Master to read from Salve ID\#2 |
| D1357 | Starting reference for Master to read from Salve ID\#3 |
| D1358 | Starting reference for Master to read from Salve ID\#4 |
| D1359 | Starting reference for Master to read from Salve ID\#5 |
| D1360 | Starting reference for Master to read from Salve ID\#6 |
| D1361 | Starting reference for Master to read from Salve ID\#7 |
| D1362 | Starting reference for Master to read from Salve ID\#8 |
| D1363 | Starting reference for Master to read from Salve ID\#9 |
| D1364 | Starting reference for Master to read from Salve ID\#10 |


| D1365 | Starting reference for Master to read from Salve ID\#11 |
| :--- | :--- |
| D1366 | Starting reference for Master to read from Salve ID\#12 |
| D1367 | Starting reference for Master to read from Salve ID\#13 |
| D1368 | Starting reference for Master to read from Salve ID\#14 |
| D1369 | Starting reference for Master to read from Salve ID\#15 |
| D1370 | Starting reference for Master to read from Salve ID\#16 |
| D1371 | When M1070 is on, it determines the time unit of CH0 PWM pulse output |
| D1372 | When M1070 is on, it determines the time unit of CH1 PWM pulse output |
| D1379 | Y4, Y5 start frequency |
| D1380 | Y6, Y7 start frequency |
| D1381 | Y4, Y5 acceleration and deceleration time |
| D1382 | Y6, Y7 acceleration and deceleration time |
| D1383 | Y10, Y11 acceleration and deceleration time |
| D1384 | Y12, Y13 acceleration and deceleration time |
| D1385 | Y14, Y15 acceleration and deceleration time |
| D1386 | Y16, Y17 acceleration and deceleration time |
| D1387 | Y20, Y21 acceleration and deceleration time |
| D1388 | Y22, Y23 acceleration and deceleration time |
| D1389 | Y24, Y25 acceleration and deceleration time |
| D1390 | Y26, Y27 acceleration and deceleration time |
| D1391 | Y30, Y31 acceleration and deceleration time |
| D1392 | Y32, Y33 acceleration and deceleration time |
| D1393 | Y34, Y35 acceleration and deceleration time |
| D1394 | Y36, Y37 acceleration and deceleration time |
| D1395 | Y40, Y41 acceleration and deceleration time |
| D1396 | Y42, Y43 acceleration and deceleration time |
| D1397 | Y44, Y45 acceleration and deceleration time |
| D1398 | Y46, Y47 acceleration and deceleration time |
| D1399 | Y50, Y51 acceleration and deceleration time |
| D1400 | Y10, Y11 start frequency |
| D1401 | Y12, Y13 start frequency |
| D1402 | Y14, Y15 start frequency |
| D1403 | Y16, Y17 start frequency |
| D1404 | Y20, Y21 start frequency |
| D1405 | Y22, Y23 start frequency |
| D1406 | Y24, Y25 start frequency |
| D1407 | Y26, Y27 start frequency |
| D1408 | Y30, Y31 start frequency |
| D1409 | Y32, Y33 start frequency |
| D1410 | Y34, Y35 start frequency |
| D1411 | Y34, Y35 start frequency |
| D1412 | Y40, Y41 start frequency |
| D1413 | Y42, Y43 start frequency |
| D1414 | Y44, Y45 start frequency |
| D1415 | Y50, Y51 start frequency |
| D1416 | Y50, Y51 start frequency |
| D1417 | Y52, Y53 start frequency |
| D1418 | Y54, Y55 start frequency |
| D1419 | Y56, Y57 start frequency |
|  |  |


| D1420 | Y52, Y53 acceleration and deceleration time |
| :--- | :--- |
| D1421 | Y54, Y55 acceleration and deceleration time |
| D1422 | Y56, Y57 acceleration and deceleration time |
| D1423 | Write the start communication address setting of slave ID\# 9 |
|  | Write the start communication address setting of slave station ID\# 10 D1425:: Write the start |
| D1424 | communication address setting of slave station ID\# 11 |
| D1426 | Y0, Y1 maximum speed (LOW WORD) |
| D1427 | Y0, Y1 maximum speed (HIGH WORD) |
| D1428 | Y2, Y3 maximum speed (LOW WORD) |
| D1429 | Y2, Y3 maximum speed (HIGH WORD) |
| D1430 | Y4, Y5 maximum speed (LOW WORD) |
| D1431 | Y4, Y5 maximum speed (HIGH WORD) |
| D1432 | Y6, Y7 maximum speed (LOW WORD) |
| D1433 | Y6, Y7 maximum speed (HIGH WORD) |
| D1434 | Y10, Y11 maximum speed (LOW WORD) |
| D1435 | Y10, Y11 maximum speed (HIGH WORD) |
| D1436 | Y12, Y13 maximum speed (LOW WORD) |
| D1437 | Y12, Y13 maximum speed (HIGH WORD) |
| D1438 | Y14, Y15 maximum speed (LOW WORD) |
| D1439 | Y14, Y15 maximum speed (HIGH WORD) |
| D1440 | Y16, Y17 maximum speed (LOW WORD) |
| D1441 | Y16, Y17 maximum speed (HIGH WORD) |
| D1442 | Y20, Y21 maximum speed (LOW WORD) |
| D1443 | Y20, Y21 maximum speed (HIGH WORD) |
| D1444 | Y22, Y23 maximum speed (LOW WORD) |
| D1445 | Y22, Y23 maximum speed (HIGH WORD) |
| D1446 | Y24, Y25 maximum speed (LOW WORD) |
| D1447 | Y24, Y25 maximum speed (HIGH WORD) |
| D1448 | Y26, Y27 maximum speed (LOW WORD) |
| D1449 | Y26, Y27 maximum speed (HIGH WORD) |
| D1450 | Setting the data write length for slave ID\#1 |
| D1451 | Set the data write length for slave ID\#3 |
| D1452 | Set the data write length for slave ID\#3 |
| D1453 | Set the data write length for slave ID\#4 |
| D1454 | Setting the data write length for slave ID\#5 |
| D1455 | Setting the data write length for slave ID\#6 |
| D1456 | Setting the data write length for slave ID\#7 |
| D1457 | Set the data write length for slave ID\#8 |
| D1458 | Setting the data write length for slave ID\#9 |
| D1459 | Set the data write length for slave ID\#10 |
| D1460 | Setting the data write length for slave ID\#11 |
| D1461 | Set the data write length for slave ID\#12 |
| D1462 | Set the data write length for slave ID\#13 |
| D1463 | Setting the data write length for slave ID\#14 |
| D1464 | Set the data write length for slave ID\#15 |
| D1465 | Setting the data write length for slave ID\#16 |
| D1472 | Y44, Y45 sent pulses (low word) |
| D1473 | Y44, Y45 sent pulse number (high word) |
| D1474 | Y44, Y45 target position (low word) |


| D1475 | Y44, Y45 target position (high word) |
| :---: | :---: |
| D1478 | Y44, Y45 acceleration (floating point number low word) |
| D1479 | Y44, Y45 acceleration (floating point number high word) |
| D1480 | Y44, Y45 current speed (floating point low word) |
| D1481 | Y44, Y45 current speed (floating point high word) |
| D1482 | Y44, Y45 target speed (floating point low word) |
| D1483 | Y44, Y45 target speed (floating point high word) |
| D1488 | Y46, Y47 sent pulse number (low word) |
| D1489 | Y46, Y47 sent pulse number (high word) |
| D1490 | Y46, Y47 target position (low word) |
| D1491 | Y46, Y47 target position (high word) |
| D1494 | Y46, Y47 acceleration (floating point number low word) |
| D1495 | Y46, Y47 acceleration (floating point number low word) |
| D1496 | Y46, Y47 current speed (floating point low word) |
| D1497 | Y46, Y47 current speed (floating point high word) |
| D1498 | Y46, Y47 target speed (floating point low word) |
| D1499 | Y46, Y47 target speed (floating point high word) |
| D1504 | Y50, Y51 sent pulse number (low word) |
| D1505 | Y50, Y51 sent pulse number (high word) |
| D1506 | Y50, Y51 target position (low word) |
| D1507 | Y50, Y51 target position (high word) |
| D1510 | Y50, Y51 acceleration (floating point number low word) |
| D1511 | Y50, Y51 acceleration (floating point number high word) |
| D1512 | Y50, Y51 current speed (floating point low word) |
| D1513 | Y50, Y51 current speed (floating point high word) |
| D1514 | Y50, Y51 target speed (floating point low word) |
| D1515 | Y50, Y51 target speed (floating point high word) |
| D1520 | Y52, Y53 sent pulse number (low word) |
| D1521 | Y52, Y53 sent pulse number (high word) |
| D1522 | Y52, Y53 target position (low word) |
| D1523 | Y52, Y53 target position (high word) |
| D1526 | Y52, Y53 acceleration (floating point number low word) |
| D1527 | Y52, Y53 acceleration (floating point number high word) |
| D1528 | Y52, Y53 current speed (floating point low word) |
| D1529 | Y52, Y53 current speed (floating point high word) |
| D1530 | Y52, Y53 target speed (floating point low word) |
| D1531 | Y52, Y53 target speed (floating point high word) |
| D1536 | Y54, Y55 sent pulse number (low word) |
| D1537 | Y54, Y55 sent pulse number (high word) |
| D1538 | Y54, Y55 target position (low word) |
| D1539 | Y54, Y55 target position (high word) |
| D1542 | Y54, Y55 acceleration (floating point number low word) |
| D1543 | Y54, Y55 acceleration (floating point high word) |
| D1544 | Y54, Y55 current speed (floating point low word) |
| D1545 | Y54, Y55 current speed (floating point high word) |
| D1546 | Y54, Y55 target speed (floating point low word) |
| D1547 | Y54, Y55 target speed (floating point high word) |
| D1552 | Y56, Y57 sent pulse number (low word) |
| D1553 | Y56, Y57 sent pulse number (high word) |


| D1554 | Y56, Y57 target position (low word) |
| :---: | :---: |
| D1555 | Y56, Y57 target position (high word) |
| D1558 | Y56, Y57 acceleration (floating point number low word) |
| D1559 | Y56, Y57 acceleration (floating point number high word) |
| D1560 | Y56, Y57 current speed (floating point low word) |
| D1561 | Y56, Y57 current speed (floating point high word) |
| D1562 | Y56, Y57 target speed (floating point low word) |
| D1563 | Y56, Y57 target speed (floating point high word) |
| D1568 | Y0, Y1 absolute position after ZRN instruction |
| D1569 | Y2, Y3 Absolute position after ZRN instruction |
| D1570 | Y4, Y5 absolute position after ZRN instruction |
| D1571 | Y6, Y7 absolute position after ZRN instruction |
| D1572 | Y10, Y11 Absolute position after executing ZRN instruction |
| D1573 | Y12, Y13 Absolute position after ZRN instruction |
| D1574 | Y14, Y15 Absolute position after ZRN instruction |
| D1575 | Y16, Y17 Absolute position after ZRN instruction is executed |
| D1576 | Y20, Y21 Absolute position after executing ZRN instruction |
| D1577 | Y22, Y23 Absolute position after executing ZRN instruction |
| D1578 | Y24, Y25 Absolute position after ZRN instruction is executed |
| D1579 | Y26, Y27 Absolute position after executing ZRN instruction |
| D1580 | Y30, Y31 Absolute position after executing ZRN instruction |
| D1581 | Y32, Y33 Absolute position after ZRN instruction |
| D1582 | Y34, Y35 absolute position after ZRN instruction |
| D1583 | Y36, Y37 Absolute position after executing ZRN instruction |
| D1584 | Y40, Y41 Absolute position after executing ZRN instruction |
| D1585 | Y42, Y43 Absolute position after executing ZRN instruction |
| D1586 | Y44, Y45 absolute position after ZRN instruction |
| D1587 | Y46, Y47 Absolute position after ZRN instruction |
| D1588 | Y50, Y51 Absolute position after executing ZRN instruction |
| D1589 | Y52, Y53 absolute position after ZRN instruction |
| D1590 | Y54, Y55 Absolute position after ZRN instruction |
| D1591 | Y56, Y57 absolute position after ZRN instruction |
| D1648 | Y0, Y1 sent pulse number (low word) |
| D1649 | Y0, Y1 sent pulse number (high word) |
| D1650 | Y0, Y1 target position (low word) |
| D1651 | Y0, Y1 target position (high word) |
| D1654 | Y0, Y1 acceleration (floating point number low word) |
| D1655 | Y0, Y1 acceleration (floating point number high word) |
| D1656 | Y0, Y1 current speed (floating point low word) |
| D1657 | Y0, Y1 current speed (floating point high word) |
| D1658 | Y0, Y1 target speed (floating point low word) |
| D1659 | Y0, Y1 target speed (floating point high word) |
| D1664 | Y2, Y3 sent pulse number (low word) |
| D1665 | Y2, Y3 sent pulse number (high word) |
| D1666 | Y2, Y3 target position (low word) |
| D1667 | Y2, Y3 target position (high word) |
| D1670 | Y2, Y3 acceleration (floating point number low word) |
| D1671 | Y2, Y3 acceleration (floating point high word) |
| D1672 | Y2, Y3 current speed (floating point low word) |


| D1673 | Y2, Y3 current speed (floating point high word) |
| :---: | :---: |
| D1674 | Y2, Y3 target speed (floating point low word) |
| D1675 | Y2, Y3 target speed (floating point high word) |
| D1680 | Y4, Y5 sent pulse number (low word) |
| D1681 | Y4, Y5 sent pulse number (high word) |
| D1682 | Y4, Y5 target position (low word) |
| D1683 | Y4, Y5 target position (high word) |
| D1686 | Y4, Y5 acceleration (floating point number low word) |
| D1687 | Y4, Y5 acceleration (floating point high word) |
| D1688 | Y4, Y5 current speed (floating point low word) |
| D1689 | Y4, Y5 current speed (floating point high word) |
| D1690 | Y4, Y5 target speed (floating point low word) |
| D1691 | Y4, Y5 target speed (floating point high word) |
| D1696 | Y6, Y7 sent pulse number (low word) |
| D1697 | Y6, Y7 sent pulse number (high word) |
| D1698 | Y6, Y7 target position (low word) |
| D1699 | Y6, Y7 target position (high word) |
| D1702 | Y6, Y7 acceleration (floating point number low word) |
| D1703 | Y6, Y7 acceleration (floating point number high word) |
| D1704 | Y6, Y7 current speed (floating point low word) |
| D1705 | Y6, Y7 current speed (floating point high word) |
| D1706 | Y6, Y7 target speed (floating point low word) |
| D1707 | Y6, Y7 target speed (points high word) |
| D1712 | Y10, Y11 sent pulse number (low word) |
| D1713 | Y10, Y11 sent pulse number (high word) |
| D1714 | Y10, Y11 target position (low word) |
| D1715 | Y10, Y11 target position (high word) |
| D1718 | Y10, Y11 acceleration (floating point number low word) |
| D1719 | Y10, Y11 acceleration (floating point number high word) |
| D1720 | Y10, Y11 current speed (floating point low word) |
| D1721 | Y10, Y11 current speed (floating point high word) |
| D1722 | Y10, Y11 target speed (floating point low word) |
| D1723 | Y10, Y11 target speed (floating point high word) |
| D1728 | Y12, Y13 sent pulse number (low word) |
| D1729 | Y12, Y13 sent pulse number (high word) |
| D1730 | Y12, Y13 target position (low word) |
| D1731 | Y12, Y13 target position (high word) |
| D1734 | Y12, Y13 acceleration (floating point number low word) |
| D1735 | Y12, Y13 acceleration (floating point number high word) |
| D1736 | Y12, Y13 current speed (floating point low word) |
| D1737 | Y12, Y13 current speed (floating point high word) |
| D1738 | Y12, Y13 target speed (floating point low word) |
| D1739 | Y12, Y13 target speed (floating point high word) |
| D1744 | Y14, Y15 sent pulse number (low word) |
| D1745 | Y14, Y15 sent pulse number (high word) |
| D1746 | Y14, Y15 target position (low word) |
| D1747 | Y14, Y15 target position (high word) |
| D1750 | Y14, Y15 acceleration (floating point number low word) |
| D1751 | Y14, Y15 acceleration (floating point number high word) |


| D1752 | Y14, Y15 current speed (floating point low word) |
| :---: | :---: |
| D1753 | Y14, Y15 current speed (floating point high word) |
| D1754 | Y14, Y15 target speed (floating point low word) |
| D1755 | Y14, Y15 target speed (floating point high word) |
| D1760 | Y16, Y17 sent pulse number (low word) |
| D1761 | Y16, Y17 sent pulse number (high word) |
| D1762 | Y16, Y17 target position (low word) |
| D1763 | Y16, Y17 target position (high word) |
| D1766 | Y16, Y17 acceleration (floating point number low word) |
| D1767 | Y16, Y17 acceleration (floating point number high word) |
| D1768 | Y16, Y17 current speed (floating point low word) |
| D1769 | Y16, Y17 current speed (floating point high word) |
| D1770 | Y16, Y17 target speed (floating point low word) |
| D1771 | Y16, Y17 target speed (floating point high word) |
| D1776 | Y20, Y21 sent pulse number (low word) |
| D1777 | Y20, Y21 sent pulse number (high word) |
| D1778 | Y20, Y21 target position (low word) |
| D1779 | Y20, Y21 target position (high word) |
| D1782 | Y20, Y21 acceleration (floating point number low word) |
| D1783 | Y20, Y21 acceleration (floating point number high word) |
| D1784 | Y20, Y21 current speed (floating point low word) |
| D1785 | Y20, Y21 current speed (floating point high word) |
| D1786 | Y20, Y21 target speed (floating point low word) |
| D1787 | Y20, Y21 target speed (floating point high word) |
| D1792 | Y22, Y23 sent pulse number (low word) |
| D1793 | Y22, Y23 sent pulse number (high word) |
| D1794 | Y22, Y23 target position (low word) |
| D1795 | Y22, Y23 target position (high word) |
| D1798 | Y22, Y23 acceleration (floating point number low word) |
| D1799 | Y22, Y23 acceleration (floating point high word) |
| D1800 | Y22, Y23 current speed (floating point low word) |
| D1801 | Y22, Y23 current speed (floating point high word) |
| D1802 | Y22, Y23 target speed (floating point low word) |
| D1803 | Y22, Y23 target speed (floating point high word) |
| D1808 | Y24, Y25 sent pulse number (low word) |
| D1809 | Y24, Y25 sent pulse number (high word) |
| D1810 | Y24, Y25 target position (low word) |
| D1811 | Y24, Y25 target position (high word) |
| D1814 | Y24, Y25 acceleration (floating point low word) |
| D1815 | Y24, Y25 acceleration (floating point number high word) |
| D1816 | Y24, Y25 current speed (floating point low word) |
| D1817 | Y24, Y25 current speed (floating point high word) |
| D1818 | Y24, Y25 target speed (floating point low word) |
| D1819 | Y24, Y25 target speed (floating point high word) |
| D1824 | Y26, Y27 sent pulse number (low word) |
| D1825 | Y26, Y27 sent pulse number (high word) |
| D1826 | Y26, Y27 target position (low word) |
| D1827 | Y26, Y27 target position (high word) |
| D1830 | Y26, Y27 acceleration (floating point number low word) |


| D1831 | Y26, Y27 acceleration (floating point number high word) |
| :---: | :---: |
| D1832 | Y26, Y27 current speed (floating point low word) |
| D1833 | Y26, Y27 current speed (floating point high word) |
| D1834 | Y26, Y27 target speed (floating point low word) |
| D1835 | Y26, Y27 target speed (floating point high word) |
| D1840 | Y30, Y31 sent pulse number (low word) |
| D1841 | Y30, Y31 sent pulse number (high word) |
| D1842 | Y30, Y31 target position (low word) |
| D1843 | Y30, Y31 target position (high word) |
| D1846 | Y30, Y31 acceleration (floating point number low word) |
| D1847 | Y30, Y31 acceleration (floating point number high word) |
| D1848 | Y30, Y31 current speed (floating point low word) |
| D1849 | Y30, Y31 current speed (floating point high word) |
| D1850 | Y30, Y31 target speed (floating point low word) |
| D1851 | Y30, Y31 target speed (floating point high word) |
| D1856 | Y32, Y33 sent pulse number (low word) |
| D1857 | Y32, Y33 sent pulse number (high word) |
| D1858 | Y32, Y33 target position (low word) |
| D1859 | Y32, Y33 target position (high word) |
| D1862 | Y32, Y33 acceleration (floating point low word) |
| D1863 | Y32, Y33 acceleration setting (floating point number high word) |
| D1864 | Y32, Y33 current speed (floating point low word) |
| D1865 | Y32, Y33 current speed (floating point high word) |
| D1866 | Y32, Y33 target speed (floating point low word) |
| D1867 | Y32, Y33 target speed (floating point high word) |
| D1872 | Y34, Y35 sent pulse number (low word) |
| D1873 | Y34, Y35 sent pulse number (high word) |
| D1874 | Y34, Y35 target position (low word) |
| D1875 | Y34, Y35 target position (high word) |
| D1878 | Y34, Y35 acceleration (floating point number low word) |
| D1879 | Y34, Y35 acceleration (floating point number high word) |
| D1880 | Y34, Y35 current speed (floating point low word) |
| D1881 | Y34, Y335 current speed (floating point high word) |
| D1882 | Y34, Y35 target speed (floating point low word) |
| D1883 | Y34, Y35 target speed (floating point high word) |
| D1888 | Y36, Y37 sent pulse number (low word) |
| D1889 | Y36, Y37 sent pulse number (high word) |
| D1890 | Y36, Y37 target position (low word) |
| D1891 | Y36, Y37 target position (high word) |
| D1894 | Y36, Y37 acceleration (floating point number low word) |
| D1895 | Y36, Y37 acceleration (floating point number high word) |
| D1896 | Y36, Y37 current speed (floating point low word) |
| D1897 | Y36, Y37 current speed (floating point high word) |
| D1898 | Y36, Y37 target speed (floating point low word) |
| D1899 | Y36, Y37 target speed (floating point high word) |
| D1904 | Y40, Y41 sent pulse number (low word) |
| D1905 | Y40, Y41 sent pulse number (high word) |
| D1906 | Y40, Y41 target position (low word) |
| D1907 | Y40, Y41 target position (high word) |


| D1910 | Y40, Y41 acceleration (floating point number low word) |
| :--- | :--- |
| D1911 | Y40, Y41 acceleration (floating point number high word) |
| D1912 | Y40, Y41 current speed (floating point low word) |
| D1913 | Y40, Y41 current speed (floating point high word) |
| D1914 | Y40, Y41 target speed (floating point low word) |
| D1915 | Y40, Y41 target speed (floating point high word) |
| D1920 | Y42, Y43 sent pulse number (low word) |
| D1921 | Y42, Y43 sent pulse number (high word) |
| D1922 | Y42, Y43 target position (low word) |
| D1923 | Y42, Y43 target position (high word) |
| D1926 | Y42, Y43 acceleration (floating point number low word) |
| D1927 | Y42, Y43 acceleration (floating point number high word) |
| D1928 | Y42, Y43 current speed (floating point low word) |
| D1929 | Y42, Y43 current speed (floating point high word) |
| D1930 | Y42, Y43 target speed (floating point low word) |
| D1931 | Y42, Y43 target speed (floating point high word) |
| D1936 | Y0, Y1 deceleration time |
| D1937 | Y2, Y3 deceleration time |
| D1938 | Y4, Y5 deceleration time |
| D1939 | Y6, Y7 deceleration time |
| D1940 | Y10, Y11 deceleration time |
| D1941 | Y12, Y13 deceleration time |
| D1942 | Y14, Y15 deceleration time |
| D1943 | Y16, Y17 deceleration time |
| D1944 | Y20, Y21 deceleration time |
| D1945 | Y22, Y23 deceleration time |
| D1946 | Y24, Y25 deceleration time |
| D1947 | Y26, Y27 deceleration time |
| D1948 | Y30, Y31 deceleration time |
| D1949 | Y32, Y33 deceleration time |
| D1950 | Y34, Y35 deceleration time |
| D1951 | Y36, Y37 deceleration time |
| D1952 | Y40, Y41 deceleration time |
| D1953 | Y42, Y43 deceleration time |
| D1954 | Y44, Y45 deceleration time |
| D1955 | Y46, Y47 deceleration time |
| D1956 | Y50, Y51 deceleration time |
| D1957 | Y52, Y53 deceleration time |
| D1958 | Y54, Y55 deceleration time |
| D1959 | Y56, Y57 deceleration time |
| D1966 | Number of modules |
| D1968 | Expansion module 1 read address |
| D1969 | Expansion module 2 read address |
| D1970 | Expansion module 3 read address |
| D1971 | Expansion module 4 read address |
| D1972 | Expansion module 5 read address |
| D1973 | Expansion module 6 read address |
| D1974 | Expansion module 7 read address |
| D1975 | Expansion module 8 read address |
|  |  |


| D1976 | Expansion module 9 read address |
| :--- | :--- |
| D1977 | Expansion module 10 read address |
| D1978 | Expansion module 11 read address |
| D1979 | Expansion module 12 read address |
| D1980 | Expansion module 13 read address |
| D1981 | Expansion module 14 read address |
| D1982 | Expansion module 15 read address |
| D1983 | Expansion module 16 read address |
| D1984 | Expansion module 1 write address |
| D1985 | Expansion module 2 write address |
| D1986 | Expansion module 3 write address |
| D1987 | Expansion module 4 write address |
| D1988 | Expansion module 5 write address |
| D1989 | Expansion module 6 write address |
| D1990 | Expansion module 7 write address |
| D1991 | Expansion module 8 write address |
| D1992 | Expansion module 9 write address |
| D1993 | Expansion module 10 write address |
| D1994 | Expansion module 11 write address |
| D1995 | Expansion module 12 write address |
| D1996 | Expansion module 13 write address |
| D1997 | Expansion module 14 write address |
| D1998 | Expansion module 15 write address |
| D1999 | Expansion module 16 write address |

### 2.11 Functions of Special Auxiliary Relays and Special Registers

## - PLC Operation Flag (M1000-M1003)



1. M1001: M1001 (B contact) is constantly "On" during operation and detection. When PLC is in RUN status, M1001 remains "On" .
2, M1002: M1002 is to enable single positive pulse at the moment when RUN is activated (Normally OFF). The pulse width $=1$ scan time. Use this contact for all kinds of initial settings.
2. M1003: M1003 is to enable single negative pulse at the moment when RUN is activated (Normally ON).

3. Monitor timer is used for monitoring PLC scan time. When the scan time exceeds the set time in the monitor timer, the red ERROR LED indicator remains beaconing and all outputs will be "Off" .
4. The initial set value of the time in the monitor timer is 200 ms . If the program is long or the operation is too complicated, MOV instruction can be used for changing the set value. See the example below for SV = 300 ms .
5. The maximum set value in the monitor timer is $32,767 \mathrm{~ms}$. Please be noted that if the SV is too big, the timing of detecting operational errors will be delayed. Therefore, it is suggested that you remain the scan time of shorter than 200ms.
6. Complicated instruction operations or too many extension modules being connected to the CPU will result in the scan time being too long. Check D1010 ~ D1012 to see if the scan time exceeds the SV in D1000. In this case, besides modifying the SV in D1000, you can also add WDT instruction (ZL 07) into the PLC program. When the CPU execution progresses to WDT instruction, the internal monitor timer will be cleared as "0" and the scan time will not exceed the set value in the monitor timer.。

- Program Capacity D1002

Program capacity of PLC model: 30K steps


1. When errors occur in syntax check, ERROR LED indicator will flash and special relay M1004 = On.

2, Timings for PLC syntax check:
a) When the power goes from "Off" to "On".
3. The syntax check may start due to illegal use of instruction operands (devices) or incorrect program syntax loop. The error can be detected by the error code in D1004 and error table. The address where the error exists will be stored in D1137. (The address value in D1 137 will be invalid if the error is a general loop error.)

## Scan Time-out

Timer
M1008, D1008

1. $\mathrm{M} 1008=$ On: Scan time-out occurs during the execution of the program, and PLC ERROR LED indicator remains beaconing.
2, Monitor the content (STEP address when WDT timer is "On").。

- Scan Time Monitor D1010~D1012

The present value, minimum value and maximum value of scan time are stored in D1010~D1012.
1, D1010: Present scan time value.
2, D1011: Minimum scan time value.
3. D1012: Maximum scan time value.

- Internal Clock Pulse M1011~M1014

The PLC mainframe has the following 4 kinds of clock pulses. As long as the PLC is powered on, these 4 kinds of clock pulses will act automatically.


M1014 (60 sec)


- Real Time Clock

M1016, M1017
M1076 D1313~D1319

1. Special M and special D relevant to RTC

| No. | Name | Function |
| :---: | :---: | :--- |
| M1016 | Year (in A.D.) in RTC | Off: display the last 2 digits of year in A.D. <br> On: display the last 2 digits of year in A.D. plus 2,000 |
| M1017 | $\pm 30$ seconds <br> correction | From "Off" to "On", the correction is enabled. $0 \sim 29$ second: minute <br> intact; second reset to 0 30~59 second: minute + 1; second reset to 0 |
| M1076 | RTC malfunction | Set value exceeds the range; dead battery |
| M1082 | Flag change on RTC | On: Modification on RTC |
| D1313 | Second | $0 \sim 59$ |
| D1314 | Minute | $0 \sim 59$ |
| D1315 | Hour | $0 \sim 23$ |
| D1316 | Day | $1 \sim 31$ |
| D1317 | Month | $1 \sim 12$ |
| D1318 | Week | $1 \sim 7$ |
| D1319 | Year | $0 \sim 99$ (last 2 digits of Year in A.D.) |

2. If the set value in RTC is incorrect, the time will be recovered as "Saturday, 00:00 Jan. 1, 2000" when PLC is powered and restarted.
3. D1313 ~ D 1319 will immediately update the RTC only when in TRD instruction or monitoring mode.
4. D1018 and D1019 are combined as 32-bit data register for storing the floating point value of $\Pi(\mathrm{PI})$.

2, Floating point value $=\mathrm{H} 40490$ FDB

Adjustment on Input Terminal Response Time D1020, D1021

1. D1020 can be used for setting up the response time of receiving pulses at X0 $\sim X 7$. (Setup range: $0 \sim$ 60; Unit: ms).
2. D1021 can be used for setting up the response time of receiving pulses at X10~X17.(Setup range: $0 \sim$ 60; Unit: ms).
3. When the power of PLC goes from "Off" to "On", the content of D1020 and D1021 turn to 10 automatically.


4, If the following programs are executed during the program, the response time of $X 0 \sim X 7$ will be set to Oms. The fastest response time of input terminals is $50 \mu \mathrm{~s}$ due to that all terminals are connected with RC filter loop.

normally ON contact
5. There is no need to make adjustment on response time when using high-speed counters and interruptions during the program.

- Execution completed flag: After the pulse output is completed, $M$ will be set to ON

| Applicable instructions: ZL 155 DABSR, ZL 156 ZRN, ZL 158 DRVI, ZL 159 DRVA, PLSY, PLSR |  |  |  |
| :--- | :--- | :--- | :--- |
| Output <br> device | Pulse output complete flag | Output device | Pulse output complete <br> flag |
| Y0, Y1 | M1029 | Y30, Y31 | M1110 |
| Y2, Y3 | M1030 | Y32, Y33 | M1111 |
| Y4, Y5 | M1036 | Y34, Y35 | M1112 |
| Y6, Y7 | M1037 | Y36, Y37 | M11113 |
| Y10, Y11 | M1102 | Y40, Y41 | M1114 |
| Y12, Y13 | M1103 | Y42, Y43 | M1115 |
| Y14, Y15 | M1104 | Y44, Y45 | M1116 |


| Y16, Y17 | M1105 | Y46, Y47 | M1117 |
| :--- | :--- | :--- | :--- |
| Y20, Y21 | M1106 | Y50, Y51 | M1118 |
| Y22, Y23 | M1107 | Y52, Y53 | M1119 |
| Y24, Y25 | M1108 | Y54, Y55 | M1205 |
| Y26, Y27 | M1109 | Y56, Y57 | M1206 |

1, After the pulse output is completed, the corresponding pulse completion flag $M$ point will be set to On, and when the pulse sending command is Off, the corresponding pulse completion flag $M$ point will turn Off. When the instruction is restarted next time, the corresponding pulse completion flag bit M becomes Off again, and then becomes On again after completion.
2, ZL 63 INCD: When the comparison of the specified number of groups is completed, M1029 will be On for one scan cycle.
3, ZL 67 RAMP, ZL 69 SORT:

- When the command is executed, M1029 = On, and M1029 must be cleared by the user.
- When this command is Off, M1029 becomes Off.


### 2.12 Communication Addresses of Devices in PLC

| Device | Range |  | Type | PLC Com. Address (hex) | Modbus Com. Address (dec) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| S | 000~255 |  | Bit | 0000~00FF | 000001~000256 |
| S | 246~511 |  | Bit | 0100~01FF | 000257~000512 |
| S | 512~767 |  | Bit | 0200~02FF | 000513~000768 |
| S | 768~1023 |  | Bit | 0300~03FF | 000769~001024 |
| X | 000~377 (Octal) |  | Bit | 0400~04FF | 101025~101280 |
| Y | 000~377 (Octal) |  | Bit | 0500~05FF | 001281~001536 |
| T | 000~255 |  | Bit | 0600~06FF | 001537~001792 |
|  |  |  | Word | 0600~06FF | 401537~401792 |
| M | 000~255 |  | Bit | 0800~08FF | 002049~002304 |
| M | 256~511 |  | Bit | 0900~09FF | 002305~002560 |
| M | 512~767 |  | Bit | 0A00~0AFF | 002561~002816 |
| M | 768~1023 |  | Bit | OBOO~OBFF | 002817~003072 |
| M | 1024~1279 |  | Bit | 0C00~0CFF | 003073~003328 |
| M | 1280~1535 |  | Bit | OD00~0DFF | 003329~003584 |
| M | 1536~1791 |  | Bit | B000~BOFF | 045057~045312 |
| M | 1792~2047 |  | Bit | B100~B1FF | 045313~045568 |
| M | 2048~2303 |  | Bit | B200~B2FF | 045569~045824 |
| M | 2304~2559 |  | Bit | B300~B3FF | 045825~046080 |
| M | 2560~2815 |  | Bit | B400~B4FF | 046081~046336 |
| M | 2816~3071 |  | Bit | B500~B5FF | 046337~046592 |
| M | 3072~3327 |  | Bit | B600~B6FF | 046593~046848 |
| M | 3328~3583 |  | Bit | B700~B7FF | 046849~047104 |
| M | 3584~3839 |  | Bit | B800~B8FF | 047105~047360 |
| M | 3840~4095 |  | Bit | B900~B9FF | 047361~047616 |
| C | 0~199 | 16-bit | Bit | 0E00~0EC7 | 003585~003784 |
|  |  |  | Word | 0E00~0EC7 | 403585~403784 |
|  | 200~255 | 32-bit | Bit | OEC8~0EFF | 003785~003840 |
|  |  |  | Word | 0700~076F | 403785~403840 |
| D | 000~256 |  | Word | 1000~10FF | 404097~404352 |
| D | 256~511 |  | Word | 1100~11FF | 404353~404608 |
| D | 512~767 |  | Word | 1200~12FF | 404609~404864 |
| D | 768~1023 |  | Word | 1300~13FF | 404865~405120 |
| D | 1024~1279 |  | Word | 1400~14FF | 405121~405376 |
| D | 1280~1535 |  | Word | 1500~15FF | 405377~405632 |
| D | 1536~1791 |  | Word | 1600~16FF | 405633~405888 |
| D | 1792~2047 |  | Word | 1700~17FF | 405889~406144 |
| D | 2048~2303 |  | Word | 1800~18FF | 406145~406400 |
| D | 2304~2559 |  | Word | 1900~19FF | 406401~406656 |
| D | 2560~2815 |  | Word | 1A00~1AFF | 406657~406912 |


| Device | Range | Type | PLC Com. Address (hex) | Modbus Com. Address (dec) |
| :---: | :---: | :---: | :---: | :---: |
| D | 2816~3071 | Word | 1B00~1BFF | 406913~407168 |
| D | 3072~3327 | Word | 1C00~1CFF | 407169~407424 |
| D | 3328~3583 | Word | 1D00~1DFF | 407425~407680 |
| D | 3584~3839 | Word | 1E00~1EFF | 407681~407936 |
| D | 3840~4095 | Word | 1F00~1FFF | 407937~408192 |
| D | 4096~4351 | Word | 9000~90FF | 436865~437120 |
| D | 4352~4607 | Word | 9100~91FF | 437121~437376 |
| D | 4608~4863 | Word | 9200~92FF | 437377~437632 |
| D | 4864~5119 | Word | 9300~93FF | 437633~437888 |
| D | 5120~5375 | Word | 9400~94FF | 437889~438144 |
| D | 5376~5631 | Word | 9500~95FF | 438145~438400 |
| D | 5632~5887 | Word | 9600~96FF | 438401~438656 |
| D | 5888~6143 | Word | 9700~97FF | 438657~438912 |
| D | 6144~6399 | Word | 9800~98FF | 438913~439168 |
| D | 6400~6655 | Word | 9900~99FF | 439169~439424 |
| D | 6656~6911 | Word | 9A00~9AFF | 439425~439680 |
| D | 6912~7167 | Word | 9B00~9BFF | 439681~439936 |
| D | 7168~7423 | Word | 9C00~9CFF | 439937~440192 |
| D | 7424~7679 | Word | 9D00~9DFF | 440193~440448 |
| D | 7680~7935 | Word | 9E00~9EFF | 440449~440704 |
| D | 7936~8191 | Word | 9F00~9FFF | 440705~440960 |
| D | 8192~8447 | Word | A000~AOFF | 440961~441216 |
| D | 8448~8703 | Word | A 100~A1FF | 441217~441472 |
| D | 8704~8959 | Word | A200~A2FF | 441473~441728 |
| D | 8960~9215 | Word | A300~A3FF | 441729~441984 |
| D | 9216~9471 | Word | A400~A4FF | 441985~442240 |
| D | 9472~9727 | Word | A500~A5FF | 442241~442496 |
| D | 9728~9983 | Word | A600~A6FF | 442497~442752 |
| D | 9984~10239 | Word | A700~A7FF | 442753~443008 |
| D | 10234~10495 | Word | A800~A8FF | 443009~443246 |
| D | 10496~10751 | Word | A900~A9FF | 443247~443502 |
| D | 10752~11007 | Word | AA00~AAFF | 443503~443758 |
| D | 11008~11263 | Word | ABOO~ABFF | 443759~444014 |
| D | 11264~11519 | Word | ACOO~ACFF | 444015~444270 |
| D | 11520~11775 | Word | AD00~ADFF | 444271~444526 |
| D | 11776~11999 | Word | AEOO~AEDF | 444527~444750 |

## 3 Basic Instructions

### 3.1 Basic Instructions and Step Ladder Instructions

1, Basic Instructions

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| LD | Loading in A contact | X, Y, M, S, T, C |
| LDI | Loading in B contact | X, Y, M, S, T, C |
| AND | Series connection- A contact | X, Y, M, S, T, C |
| ANI | Series connection- B contact | X, Y, M, S, T, C |
| OR | Parallel connection- A contact | X, Y, M, S, T, C |
| ORI | Parallel connection- B contact | X, Y, M, S, T, C |
| ANB | Series connection- loop blocks | N/A |
| ORB | Parallel connection- loop blocks | N/A |
| MPS | Store the current result of the <br> internal PLC operations | N/A |
| MRD | Reads the current result of the <br> internal PLC operations | N/A |
| MPP | Pops (recalls and removes) the <br> currently stored result | N/A |

2, Output instructions

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| OUT | Output coil | Y, S, M |
| SET | Latched (On) | Y, S, M |
| RST | Clear the contacts or the registers | Y, M, S, T, C, D, E, F |

3. Timers, Counters

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| TMR | 16 -bit timer | T-K or T-D |
| CNT | 16-bit counter | C-K or C-D (16 bits) |
| DCNT | 32-bit counter | Z, C-K or C-D (32 bits)) |

4, Main control instructions

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| $M C$ | Master control start | N0~N7 |
| $M C R$ | Master control reset | N0~N7 |

1. nstructions for detecting the contacts of rising-/falling-edge

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| LDP | Rising-edge detection operation | S, X, Y, M, T, C |
| LDF | Falling-edge detection operation | S, X, Y, M, T, C |
| ANDP | Rising-edge series connection | S, X, Y, M, T, C |
| ANDF | Falling-edge series connection | S, X, Y, M, T, C |
| ORP | Rising-edge parallel connection | S, X, Y, M, T, C |


| ORF | Falling－edge parallel connection | S，X，Y，M，T，C |
| :--- | :--- | :--- |

2，Rising－／falling－edge output instructions

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| PLS | Rising－edge output | Y，M |
| PLF | Falling－edge output | Y，M |

3．End instruction

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| END | Program ends | N／A |

4，Other instructions

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| NOP | No operation | N／A |
| INV | Inverting operation | N／A |
| P | Pointer | P0～P255 |
| I | Interruption program marker | I口ロロ |

## 5，Step ladder instructions

| Instruction Code | Function | Operands |
| :--- | :--- | :--- |
| STL | Step transition ladder start <br> instruction | S |
| RET | Step transition ladder return <br> instruction | N／A |

## 3.2 [LD] , [LDI] , [AND]

| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| LD | Loading in A contact |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 |  | D0~D11999

1. Instruction description: LD instruction is used to connect the contacts to the bus. Support $X, Y, M, S$ components can be modified, for example: LD XOE1

## 2, Program Example:



| Mnemonic | Function |  |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| LDI | Loading in B contact |  |  |  |  |  |  |  |
|  | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |  |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |  |

1. Instruction description: LD instruction is used to connect the contacts to the bus. Support X, Y, M, S components can be modified, for example: LDI XOEI
2, Program Example:


| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| AND | Series connection- A contact |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |

1, Instruction description: Use the AND instruction to connect a contact in series. The number of serial contacts is not limited, and this instruction can be used multiple times. Support X, Y, M, S components can be modified, for example: AND XOEI

## 2, Program Example:



## 3.3 [ANI] , [OR] , [ORI]

| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ANI | Series connection- B contact |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |

1. Instruction description: One contact can be connected in series with AN instruction. The number of serial contacts is not limited, and this instruction can be used multiple times. Support X, Y, M, S components can be modified, for example: ANI XOE I

## 2, Program Example



| Mnemonic | Function |  |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| OR | Parallel connection- A contact |  |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |  |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |  |

1. Instruction description: OR is used as a parallel connection command for a contact. Support $\mathrm{X}, \mathrm{Y}$, M, S components can be modified, for example: OR XIE1

## 2, Program Example:



| Mnemonic | Function |  |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ORI | Parallel connection- B contact |  |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |  |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |  |

1. Instruction description: ORI is used as a parallel connection command for a contact. Support X, Y, M, S components can be modified, for example: ORI XIEI
2, Program Example:


## 3.4 [ANB] , [ORB]

| Mnemonic | Function |
| :--- | :---: |
| ANB | Series connection-loop blocks |
| Operand | $\mathrm{N} / \mathrm{A}$ |

1, Instruction description: To perform the "AND" operation of the preserved logic results and content in the accumulative register.

## 2, Program Example:

| Ladder diagram: | Instruction code: | Operation: |
| :--- | :--- | :--- | :--- |
| Lloading in contact A of X 0 |  |  |


| Mnemonic | Function |
| :--- | :---: |
| ORB | Parallel connection- loop blocks |
| Operand | N/A |

1. Instruction description: To perform the "OR" operation of the preserved logic results and content in the accumulative register.

## 2, Program Example:



| Instruction code: | Operation: |  |
| :--- | :--- | :--- |
| LD | X0 | Loading in contact $A$ of $X 0$ |
| ANI | X1 | Connecting to $\operatorname{contact~} B$ of $X 1$ in series |
| LDI | X2 | Loading in contact $B$ of $X 2$ |
| AND | X3 | Connecting to contact $A$ of $X 3$ in series |
| ORB |  | Connecting circuit block in parallel |
| OUT | Y1 | Driving $Y 1$ coil |

## 3.5 [MPS] , [MRD], [MPP]

| Mnemonic | Function |
| :--- | :---: |
| MPS | Store the current result of the internal PLC operations |
| Operand | N/A |

1. Explanations: To save the content in the accumulative register into the operational result (the pointer of operational result will plus 1).

| Mnemonic | Function |
| :--- | :---: |
| MRD | Reads the current result of the internal PLC operations |
| Operand | $\mathrm{N} / \mathrm{A}$ |

1. Explanations: To read the operational result and store it into the accumulative register (the pointer of operational result stays intact).

| Mnemonic | Function |
| :--- | :---: |
| MPP | Pops (recalls and removes) the currently stored result |
| Operand | N/A |

1, Explanations: To retrieve the previous preserved logical operation result and store it into the accumulative register (the pointer of operational result will minus 1).
2, Program Example:


## 3.6 [OUT] , [SET], [RST]

| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| OUT | Output coil |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\times$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ | $\times$ | $\times$ |

1. Explanations: Output the result of the logic operation before the OUT instruction to the specified component. Support Y, M, S components can be modified, for example: OUT YIE2

2, Actions of coil contact:

| Operational result | OUT instruction |  |  |
| :--- | :---: | :---: | :---: |
|  | Coil | Contact |  |
|  |  | A contact (normally <br> open) | B contact (normally <br> closed) |
| FALSE | Off | Off | On |
| TRUE | On | On | Off |


| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| SET | Latched (On) |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\times$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ | $\times$ | $\times$ |

1. Explanations: When the SET instruction is driven, its designated device will be "On" and keep being On both when SET instruction is still being driven or not driven. Use RST instruction to set "Off" the device. Support Y, M, S components can be modified, for example: SET Y1E2

2, Program Example:

| Ladder diagram: |  |  |  | Instruction code: Op |  |  | Operation: |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | LD | X0 Lo |  | Loading in contact A of X0 |  |  |
|  | $\xrightarrow{\text { xo }}$ |  |  | ANI | YO Co |  | Connecting to contact B of Y 0 in series |  |  |
|  |  |  | Y1 | SET | Y1 |  | Y 1 latched ( On ) |  |  |
| Mnemonic | Function |  |  |  |  |  |  |  |  |
| RST | Clear the contacts or the registers |  |  |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | SO~S |  | TO~T255 | C0~C255 | D0~D9999 | E0~E7/F0~F7 |
|  | $\times$ | $\checkmark$ | $\checkmark$ |  |  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ |

1, Explanations: When the RST instruction is driven, the actions of the designated devices are:

| Device | Status |
| :--- | :--- |
| S, Y, M | Coil and contact will be set to "Off" |
| T, C | Present values of the timer or counter will be set to "O" , and the coil and contact will be set to <br> "Off" |
| D, E, F | The content will be set to "O" . |

## 2, Program Example:

| Ladder diagram: |  | Instruction code: |  | Operation: |
| :--- | :--- | :--- | :--- | :--- | :--- |
| X0 RST Y5  LD <br> RST Y5 Loading in contact A of X0   |  |  |  |  |

## 3.7 [TMR] , [ATMR], [CNT] , [DCNT]

| Mnemonic | Function |  |  |
| :--- | :--- | :--- | :---: |
| TMR | 16-bit timer |  |  |
| Operand | T-K | T0~T255, K0~K32,767 |  |
|  | T-D | T0~T255, D0~D11999 |  |

1, Explanations: When TMR instruction is executed, the designated coil of the timer will be On and the timer will start to time. When the set value in the timer is reached (present $\geqslant$ set value).

## 2, Program Example

| Ladder diagram: |  |  | Instruction code: |  | Operation: |
| :---: | :---: | :---: | :---: | :---: | :---: |
| X0 |  |  | LD | X0 | Loading in contact A of XO T 5 timer |
| TMR | T5 | K1000 | TMR | T5 K1000 | Set value in timer T 5 as $\mathrm{K} 1,000$ |


| Mnemonic | Function |  |
| :--- | :---: | :---: |
| ATMR | 16-bit contact type timer counter |  |
| Operand | T-K | T0~T255, K0~K32,767 |
|  | T-D | T0~T255, D0~D11999 |

1, Explanations: The instruction ATMR corresponds to the combination of AND and TMR. If the contact preceding ATMR is ON , the timer specified will begin to count. When the count value is greater than or equal to the setting value, the AND contact is ON. If the contact preceding ATMR is not ON, ATMR will automatically clear the count value.

## 2, Program Example:

| Ladder diagram: |  | Instruction code: | Operation: |
| :--- | :--- | :--- | :--- | :--- | :--- |
| Loading in contact A of XO |  |  |  |


| Mnemonic | Function |  |
| :--- | :---: | :---: |
| CNT | C-K | 16-bit counter |
| Operand | C-D | C0~C199, K0~K32,767 |
|  |  | C0~C199, D0~D11999 |

1, Explanations: When the CNT instruction goes from Off to On, the designated counter coil will be driven, and the present value, in the counter will plus 1 . When the counting reaches the set value (present value $=$ set value), the contact will be:

| NO (Normally Open) contact | Open collector |
| :--- | :--- |
| NC(Normally Close) contact | Close collector |

If there are other counting pulse inputs after the counting reaches its target, the contact and present value will stay intact. Use RST instruction to restart or reset the counting.

## 2, Program Example:



| Mnemonic | Function |  |
| :--- | :---: | :---: |
| DCNT | 32-bit counter |  |
|  | C-K | C200~C255, K-2,147,483,648~K2,147,483,647 |
|  | C-D | C200~C255, D0~D11999 |

## 1, Explanations:

- DCNT is the instruction for enabling the 32-bit high-speed counters C200~C255.
- For general purpose addition/subtraction counters C200~C234, when DCNT goes from Off to On, the present value in the counter will pulse 1 (counting up) or minus 1 (counting down) according to the modes set in special M1200~ M1235.
- For high-speed addition/subtraction counters C235 ~ C255, when the high-speed counting pulse input goes from Off to On, the counting will start its execution.
- When DCNT is Off, the counting will stop, but the existing present value in the counter will not be cleared. To clear the present value and the contact, you have to use the instruction RST C2XX. Use externally designated input points to clear the present values and contacts of high-speed addition/subtraction counters C235~C255.


## 2, Program Example:

Ladder diagram:


| Instruction code: | Operation: |  |
| :--- | :--- | :--- |
| LD | M0 | Loading in contact A of M0 |
| DCNT | C254 K1000 | Set value of counter C254 as K1,000 |

## 3.8 [MC/MCR] , [LDP], [LDF]

| Mnemonic | Function |
| :--- | :---: |
| MC/MCR | Master control Start/Reset |
| Operand | NO~N7 |

## 1. Explanations:

- MC is the main-control start instruction. When MC instruction is executed, the execution of instructions between MC and MCR will not be interrupted. When MC instruction is Off, the actions of the instructions between MC and MCR are:

| Instruction type | Explanation |
| :--- | :--- |
| General purpose timer | Present value $=0$ <br> Coil is Off, No action for the contact |
| Accumulative timer | Coil is Off, present value and contact stay intact |
| Subroutine timer | Present value $=0$ <br> Coil is Off, No action for the contact |
| Counter | Coil is Off, present value and contact stay intact |
| Coils driven by OUT instruction | All Off |
| Devices driven by SET and RST <br> instructions | Stay intact |
|  | All disabled. <br> The FOR-NEXT nested loop will still execute back and forth for N times. <br> Instructions between FOR-NEXT will act as the instructions between MC <br> and MCR. |

- MCR is the main-control end instruction that is placed in the end of the main-control program. There should not be any contact instructions prior to MCR instruction.
- MC-MCR main-control program instructions support the nested program structure (max. 8 layers) and please use the instruction in the order N0 ~ N7:


## 2, Program Example



| Instruction code: | Operation: |  |
| :--- | :--- | :--- |
| LD | X0 | Loading in A contact of X0 |
| MC | N0 | Enabling N0 common series connection contact |
| LD | X1 | Loading in A contact of X1 |
| OUT | Y0 | Driving Y0 coil |
| $:$ |  |  |
| LD | X2 | Loading in A contact of X2 |
| MC | N1 | Enabling N1 common series connection contact |
| LD | X3 | Loading in A contact of X3 |
| OUT | Y1 | Driving Y1 coil |
| $:$ |  |  |
| MCR | N1 | Disabling N1 common series connection contact |
| $:$ |  |  |
| MCR | N0 | Disabling N0 common series connection contact |
| $:$ |  |  |
| LD | X10 | Loading in A contact of X10 |
| MC | N0 | Enabling N0 common series connection contact |
| LD | X11 | Loading in A contact of X11 |
| OUT | Y10 | Driving Y10 coil |
| $:$ |  |  |
| MCR | N0 | Disabling N0 common series connection contact |


| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| LDP | Rising-edge detection operation |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |

## 1, Explanations:

The method of using LDP is the same as using LD, but the actions of the two instructions differ. LDP saves the current content and store the detected status of rising-edge to the accumulative register.

## 2, Program Example

| Ladder diagram: | Instruction code: | Operation: |  |
| :--- | :--- | :--- | :--- |
| LDP | X0 | Starting X0 rising-edge detection |  |
|  | AND | X1 | Series connecting A contact of X1 |
|  | OUT | Y1 | Driving Y1 coil |


| Mnemonic | Function |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| LDF | Falling-edge detection operation |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | MO~M4095 | SO~S 1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |

## 1, Explanations:

- The method of using LDF is the same as using LD, but the actions of the two instructions differ. LDF saves the current content and store the detected status of falling-edge to the accumulative register.


## 2, Program Example

| Ladder diagram: | Instruction code: |  | Operation: |
| :--- | :--- | :--- | :--- |
| LDF | X 0 | Starting X 0 falling-edge detection |  |
|  | AND | X 1 | Series connecting A contact of X 1 |
|  | OUT | Y 1 | Driving Y 1 coil |

## 3.9 [ANDP] , [ANDF] , [ORP] , [ORF]

| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ANDP | Rising-edge series connection |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |

## 1, Explanations:

- ANDP instruction is used in the series connection of the contacts' rising-edge detection.


## 2, Program Example



| Instruction code: |  | Operation: |
| :--- | :--- | :--- |
| LD | X0 | Loading in A contact of X0 |
| ANDP | X 1 | X1 rising-edge detection in series connection |
| OUT | Y 1 | Driving Y1 coil |


| Mnemonic | Function |  |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ANDF | Falling-edge series connection |  |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |  |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |  |

## 1, Explanations:

- ANDF instruction is used in the series connection of the contacts' falling-edge detection.


## 2, Program Example:

Ladder diagram:

Instruction code:

| LD | X0 | Loading in A contact of X0 |
| :--- | :--- | :--- |
| ANDF | X 1 | X 1 falling-edge detection in series connection |
| OUT | Y 1 | Drive Y 1 coil |


| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ORP | Rising-edge parallel connection |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |

## 1, Explanations:

- The ORP instructions are used in the parallel connection of the contact' s rising-edge detection.


## 2. Program Example

Ladder diagram:


Instruction code: Operation

| LD | X0 | Loading in A contact of X0 |
| :--- | :--- | :--- |
| ORP | X1 | X1 rising-edge detection in parallel connection |
| OUT | Y1 | Driving Y1 coil |


| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ORF | Falling-edge parallel connection |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\times$ |

## 1, Explanations:

The ORF instructions are used in the parallel connection of the contact's falling-edge detection.

## 2, Program Example

Ladder diagram:


Instruction code: Operation

| LD | X0 | Loading in A contact of X 0 |
| :--- | :--- | :--- |
| ORF | X 1 | X1 falling-edge detection in parallel connection |
| OUT | Y 1 | Driving Y 1 coil |

### 3.10 [PLS] , [PLF]

| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PLS | Rising-edge output |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\times$ | $\checkmark$ | $\checkmark$ | $\times$ | $\times$ | $\times$ | $\times$ |

## 1, Explanations:

- When X0 goes from Off to On (rising-edge trigger), PLS instruction will be executed and $S$ will send out pulses for once of 1 scan time.


## 2, Program Example

Ladder diagram: Instruction code: Operation:


| Mnemonic | Function |  |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PLF | Falling-edge output |  |  |  |  |  |  |
| Operand | X0~X377 | Y0~Y377 | M0~M4095 | S0~S1023 | T0~T255 | C0~C255 | D0~D11999 |
|  | $\times$ | $\checkmark$ | $\checkmark$ | $\times$ | $\times$ | $\times$ | $\times$ |

## 1, Explanations:

- When XO goes from On to Off (falling-edge trigger), PLF instruction will be executed and S will send out pulses for once of 1 scan time.


## 2, Program Example

Ladder diagram:


Instruction code:

| LD | X0 | Loading in A contact of X0 |
| :--- | :--- | :--- |
| PLF | M0 | M0 falling-edge output |
| LD | MO | Loading in contact A of MO |
| SET | YO | YO latched (On) |

Timing Diagram:


### 3.11 [END] , [NOP] , [INV]

| Mnemonic | Function |
| :--- | :---: |
| END | Program End |
| Operand | N/A |

## 1, Explanations:

END instruction has to be placed in the end of a ladder diagram or instruction program. PLC will start to scan from address 0 to END instruction and return to address 0 to restart the scan.

| Mnemonic | Function |
| :--- | :---: |
| NOP | No operation |
| Operand | N/A |

## 1, Explanations:

NOP instruction does not conduct any operations in the program; therefore, after the execution of NOP, the existing logical operation result will be kept. If you want to delete a certain instruction without altering the length of the program, you can use NOP instruction.

2, Program Example:

Ladder diagram:


| Instruction code: |  | Operation: |
| :--- | :--- | :--- |
| LD | X0 | Loading in B contact of X0 |
| NOP |  | No operation |
| OUT | Y1 | Driving Y1 coil |


| Mnemonic | Function |
| :--- | :---: |
| INV | Inverting Operation |
| Operand | N/A |

## 1, Explanations:

The logical operation result before INV instruction will be inverted and stored in the accumulative register.

2, Program Example:


| Instruction code: | Operation: |  |
| :--- | :--- | :--- |
| LD $\quad$ X0 | Loading in A contact of X0 |  |
| INV |  | Inverting the operation result |
| OUT Y1 | Driving Y1 coil |  |

## 3．1 Pointer［P］，Interruption Pointer［1］

| Mnemonic | Function |
| :--- | :---: |
| $\mathbf{P}$ | Pointer |
| Operand | P0～P255 |

## 1．Explanations：

Pointer P is used inZL 00 CJ and ZLOI CALL instructions．The use of $P$ does not need to start from No．0，and the No．of P cannot be repeated；otherwise，unexpected errors may occur．。

2，Program Example：


| Instruction code： |  | Operation： |
| :--- | :--- | :--- |
| LD | X0 | Loading in A contact of X0 |
| CJ | P 10 | From instruction CJ to P10 |
| $:$ |  |  |
| P10 |  | Pointer P10 |
| LD | X1 | Loading in A contact of X1 |
| OUT | Y1 | Driving Y1 coil |


| Mnemonic | Function |
| :--- | :---: |
| I | Interruption program marker（I） |
| Operand | $100 \square, 110 \square, 120 \square, 130 \square, 140 \square, 150 \square, 16 \square \square, 17 \square \square, 18 \square \square$ <br> $1010,1020,1030,1040,1050,1060,1110,1120,1130,1140,1150,1160,1170,1180$ |

## 1，Explanations：

A interruption program has to start with a interruption pointer（Iロロロ）and ends with 03 IRET．I instruction has to be used with 03 IRET， 04 El ，and 05 DI ．

2，Program Example：


| Instruction code: | Operation: |  |
| :--- | :--- | :--- |
| EI |  | Enabling interruption |
| LD | X1 | Loading A contact of X1 |
| OUT | Y1 | Driving Y1 coil |
| $\vdots$ |  |  |
| DI |  | Disabling interruption |
| $\vdots$ |  |  |
| FEND |  | Main program ends |
| I001 |  | Interruption pointer |
| LD | X2 | Loading in A contact of X2 |
| OUT | Y2 | Driving Y2 coil |
| $\vdots$ |  |  |
| IRET |  | Interruption return |

### 4.1 Step Ladder Instructions [STL], [RET]

| Mnemonic | Function | Operand |
| :--- | :---: | :---: |
| STL | Step Transition Ladder Start | SO~S1023 |

## 1, Explanations:

2, STL Sn constructs a step. When STL instruction appears in the program, the program will enter a step ladder diagram status controlled by steps. The initial status has to start from SO ~S9. RET instruction indicates the end of a step ladder diagram starting from SO ~ S9 and the bus returns to a normal ladder diagram instruction. SFC uses the step ladder diagram composed of STL/RET to complete the action of a circuit. The No. of S cannot be repeated.

| Mnemonic | Function | Operand |
| :--- | :---: | :---: |
| RET | Step Transition Ladder Return | N/A |

## 1, Explanations:

RET indicates the end of a step. There has to be a RET instruction in the end of a series of steps. One PLC program can be written in maximum 10 steps (SO ~ S9) and every step should end with a RET.

## 2, Program Example:

## Ladder diagram:



SFC:


### 4.2 Step ladder instruction action description

1, STL instruction:
STL instruction is used for designing the syntax of a sequential function chart (SFC), making the program designing similar to drawing a flow chart and allowing a more explicit and readable program. From the figure in the left hand side below, we can see very clearly the sequence to be designed, and we can convert the sequence into the step ladder diagram in the right hand side.
2, RET instruction:
RET instruction has to be written at the end of every step sequence, representing the end of a sequence. There can be more than one step sequence in a program. Therefore, we have to write in RET at the end of every step sequence. There is no limitation on the times of using RET which is used together with SO~S9.
3. If there is no RET instruction at the end of a step sequence, errors will be detected.


4, Actions of Step Ladder:
A step ladder is composed of many steps and every step controls an action in the sequence. The step ladder has to:
a) Drive the output coil
b) Designate the transition condition
c) Designate which step will take over the control from the current step

Example:


## Explanation:

When $\mathrm{S} 10=\mathrm{On}, \mathrm{YO}$ and Y 1 will be On . When $\mathrm{XO}=\mathrm{On}, \mathrm{S} 20$ will be On and Y 10 will be On. When $\mathrm{S} 10=\mathrm{Off}, \mathrm{YO}$ will be Off and Y 1 will be On.

5, Repeated Use of Output Coil:

- You can use output coils of the same No. in different steps.
- See the diagram in the right. There can be the same output device (Y0) among different statuses. YO will be On when S10 or S20 is On. Such as right diagram, there is the same output device YO in the different state. No matter S10 or S20 is On, Y0 will be On.
- YO will be Off when $S 10$ is transferring to $S 20$. After $S 20$ is On , YO will output again. Therefore in this case, YO
will be On when S 10 or S 20 is On.
- Normally in a ladder diagram, avoid repeated use of an output coil. The No. of output coil used by a step should also avoid being used when the step ladder diagram returns to a general ladder diagram.


6. Cautions for Driving Output Point:
a. See the figure below. After the step point and once LD or LDI instructions are written into the second line, the bus will not be able to connect directly to the output coil, and errors will occur in the compilation of the ladder diagram. You have to correct the diagram into the diagram in the right hand side for a correct compilation.

b. The instruction used for transferring the step (SET S $\square$ or OUT S $\square$ ) can only be executed after all the relevant outputs and actions in the current status are completed. See the figure below. The executed results by the PLC are the same, but if there are many conditions or actions in S 10 , it is recommended that you modify the diagram in the left hand side into the diagram in the right hand side. SET S20 is only executed after all relevant outputs and actions are completed, which is a more explicit sequence.

c. The RET instruction must be added after the step ladder program is completed, and RET must also be added after the STL, as shown in the figure below:


### 4.3 Step ladder programming

1. The first step in the SFC is called the "initial step", SO ~ S9. Use the initial step as the start of a sequence and end a complete sequence with RET instruction.
2, If $S T L$ instruction is not in use, step $S$ can be a general-purpose auxiliary relay.
2. When STL instruction is in use, the No. of step $S$ cannot be repeated.

4, Types of sequences:

- Single sequence: There is only one sequence without alternative divergence, alternative convergence, simultaneous divergence and simultaneous convergence in a program.
- Complicated single sequence: There is only one sequence with alternative divergence, alternative convergence, simultaneous divergence and simultaneous convergence in a program.
- Multiple sequences: There are more than one sequence in a program, maximum 10 sequences, SO ~ S9.

5. Separation of sequence: Multiple sequences are allowed to be written into the step ladder diagram.

- There are two sequences S0 and S1. The program writes in S0 ~ S30 first and S1 ~ S43 next.
- b) You can designate a step in the sequence to jump to any step in another sequence.
- When the condition below $S 21$ is true, the sequence will jump to step $S 42$ in sequence $S 1$, which is called "separating the step".


6. Restrictions on diverging sequence:

- You can use maximum 8 diverged steps in a divergence sequence.
- You can use maximum 16 loops in multiple divergence sequences or in simultaneous sequences combined into one sequence.
- You can designate a step in the sequence to jump to any step in another sequence.

7. Reset of the step and the inhibiting output:

- Use ZRST instruction to reset a step to be Off

8. Latched step:

The On/Off status of the latched step will be memorized when the power of the PLC is switched off. When the PLC is re-powered, the status before the power-off will be recovered and the execution will resume. Please be aware of the area for the latched steps.

10, Types of Sequences

- Single Sequence: The basic type of sequences

The first step in a step ladder diagram is called the initial step, which can be SO ~ S9. The steps following the initial step are general steps, which can be $\mathrm{S} 10 \sim$ S 1023 . If you are using IST instruction, $\mathrm{S} 10 \sim \mathrm{~S} 19$ will become the steps for zero return.

- Single sequence without divergence and convergence: After a sequence is completed, the control power on the steps will be given to the initial step.

Step ladder diagram


SFC:

11. Jumping Sequence

1. The control power over the step is transferred to a certain step on top.
2. The control power over the step is transferred to the step in another sequence.


## 12. Reset Sequence

When the condition at S 50 is true, S 50 will be reset and the sequence will be completed at this time.

13. Complicated Single Sequence: Including simultaneous divergence, alternative divergence, simultaneous
convergence and alternative convergence.

- Structure of simultaneous divergence:

When the condition at the current step is true, the step can be transferred to many steps. See the diagrams below. When X0 = On, S20 will be simultaneously transferred to S21, S22, S23 and S24.

Ladder diagram:


SFC:


- Structure of alternative divergence:

When the individual condition at the current status is true, the step will be transferred to another individual step. See the diagrams below. When $\mathrm{XO}=\mathrm{On}, \mathrm{S} 20$ will be transferred to S 30 ; when $\mathrm{XI}=\mathrm{On}, \mathrm{S} 20$ will be transferred to S31; when X2 = On, S20 will be transferred to S32.

Ladder diagram:


SFC:


- Structure of alternative convergence:

See the diagrams below. Depending on the condition of the input signal of which of S30, S40 and S50 becomes true first, the first one will be first transferred to S60.

## Ladder diagram:



SFC:


- Examples of alternative divergence \& alternative convergence:


SFC:


- Examples of simultaneous divergence \& simultaneous convergence:

Ladder diagram:


SFC:


## 5 Categories \& Use of Application Instructions

### 5.1 Composition of Application Instruction

1, An application instruction has two parts: the instruction and operands. The instruction part of the application instruction usually occupies 1 address (Step), The instruction part of an application instruction usually occupies 1 step, and one operand occupies 2 or 4 steps depending on the instruction is a 16-bit or 32-bit one.
Instruction: The function of the instruction
Operands: Devices for processing the operations of the instruction

2, Input of application instruction:
Some application instructions are only composed of the instruction part (mnemonic), e.g. El, DI, WDT $\cdots$. Most application instructions are composed of the instruction part and many operands. Different application instructions designate different operands. Take MOV instruction for example:


MOV instruction is to move the operand designated in $S$ to the operand designated in $D$.

| $S$ | Source operand: If there are more than 1 source operands, they will be represented as S1, <br> $S 2, \cdots$. |
| :---: | :--- |
| D | Destination operand: If there are more than 1 destination operands, they will be represented <br> as D1, D2, $\cdots$. |
| If the operand can only be constant K/H or a register, it will be represented as $\mathrm{m}, \mathrm{ml}, \mathrm{m} 2, \mathrm{n}, \mathrm{n} 1, \mathrm{n} 2, \cdots$. |  |

3. Length of operand (16-bit instruction or 32-bit instruction) Depending on the contents in the operand, the length of an operand can be 16-bit or 32-bit. Therefore, a 16 -bit instruction is for processing 16 -bit operands, and 32 -bit instruction is for processing 32-bit operands. The 32-bit instruction is indicated by adding a "D" before the 16-bit instruction.

16-bit MOV instruction: When $\mathrm{XO}=\mathrm{On}, \mathrm{K} 10$ will be sent to D 10 .


32-bit DMOV instruction: When $\mathrm{X} 1=\mathrm{On}$, the content in (D11, D10) will be sent to (D21, D20).


4, Continuous execution type / pulse execution type
In terms of command execution mode, it can be divided into two types: "continuous execution type" and "pulse execution type". Since the execution time required when the instruction is not executed is relatively short, the use of pulse execution instructions in the program as much as possible can reduce the scan cycle. The instruction marked with "P" after the instruction is the pulse execution type instruction. Some
instructions use pulse execution in most applications, such as INC, DEC and shift-related instructions, so the $\uparrow$ mark is added to the upper right of the mark of each instruction to indicate that the instruction usually uses pulse execution.

Pulse execution type:
When XO changes from Off $\rightarrow$ On, the MOVP instruction is executed once, and the scan instruction is no longer executed, so it is called pulse execution instruction.


## Continuous execution:

When $\mathrm{XI}=$ On each scan cycle, the MOV instruction is executed once, so it is called a continuous execution instruction.


When the two conditional contacts XO and $\mathrm{XI}=\mathrm{Off}$, the instruction will not be executed, and the content of the destination operand $D$ will not change.

5, Designation of operands
> Bit devices $\mathrm{X}, \mathrm{Y}, \mathrm{M}$, and S can be combined into word device, storing values and data for operaions in the form of $\mathrm{KnX}, \mathrm{KnY}, \mathrm{KnM}$ and KnS in an application instruction.
> Data register D, timer T, counter C and index register E, F are designated by general operands.
> A data register is usually in 16 bits, i.e. of the length of 1 register D. A designated 32-bit data register refers to 2 consecutive register Ds.
> If an operand of a 32-bit instruction designates DO , the 32-bit data register composed of ( $\mathrm{D} 1, \mathrm{DO}$ ) will be occupied. D1 is the higher 16 bits; D0 is the lower 16 bits. The same rule also apply to timer T, 16-bit timers and C0 ~ C199.
> When the 32-bit counters C200 ~ C254 are used as data registers, they can only be designataed by the operands of 32 -bit instructions.
6. Format of operand
> $\mathrm{X}, \mathrm{Y}, \mathrm{M}$, and S can only $\mathrm{On} /$ Off a single point and are defined as bit devices.
> 16-bit (or 32-bit) devices T, C, D, and registers E, F are defined as word devices.
> You can place $\mathrm{Kn}(\mathrm{n}=1$ refers to 4 bits. For 16 -bit instruction, $\mathrm{n}=\mathrm{K1} \sim \mathrm{~K} 4$; for 32 -bit instruction, $\mathrm{n}=\mathrm{K} 1 \sim$ K8) before bit devices $X, Y, M$ and $S$ to make it a word device for performing word-device operations. For example, K1M0 refers to 8 bits, M0 ~M7.
> When $\mathrm{XO}=\mathrm{On}$, the contents in MO ~M7 will be moved to bit0 $\sim 7$ in D10 and bit8 ~ 15 will be set to "0" .

7. Data processing of word devices combined from bit devices

| 16-bit instruction |  |
| :---: | :---: |
| Designated value: K-32,768~K32,767 |  |
| Values for designated K1 ~ K4: |  |
| K1 (4 bits)) | 0~15 |
| K2 (8 bits)) | 0~255 |
| K3 (12 bits)) | 0~4,095 |
| K4 (16 bits)) | -32,768~+32,767 |
| 32-bit instruction |  |
| Designated value: K-2,147,483,648~K2,147,483,647 |  |
| Values for designated K1 ~ K8 |  |
| K1 (4 bits)) | 0~15 |
| K2 (8 bits)) | 0~255 |
| K3 (12 bits)) | 0~4,095 |
| K4 (16 bits)) | 0~65,535 |
| K5 (20 bits) ) | 0~1,048,575 |
| K6 (24 bits) ) | 0~167,772,165 |
| K7 (28 bits) ) | 0~268,435,455 |
| K8 (32 bits)) | -2,147,483,648~+2,147,483,647 |

### 5.2 Handling of Numeric Values

1. Devices only with On/Off status are called bit devices, e.g. X, Y, M and S. Devices used exclusively for storing numeric values are called word devices, e.g. T, C, D, E and F. Bit device plus a specific bit device (place a digit before the bit device in Kn ) can be used in the operand of an application instruction in the form of numeric value.
2. $\mathrm{n}=\mathrm{K} 1 \sim \mathrm{~K} 4$ for a 16 -bit value; $\mathrm{n}=\mathrm{K} 1 \sim \mathrm{~K} 8$ for a 32-bit value. For example, K2M0 refers to an 8 -bit value composed of MO ~M7.

3. $\mathrm{K} 1 \mathrm{MO}, \mathrm{K} 2 \mathrm{MO}$, and K 3 MO are transmitted to 16 -bit registers and the vacant high bits will be filled in " 0 " . The same rule applied to when K1M0, K2M0, K3M0, K4M0, K5M0, K6MO, and K7MO are transmitted to 32-bit registers and the vacant high bits will be filled in " 0 " .
4. In the 16-bit (or 32-bit) operation, if the contents of the operand are designated as bit devices K1 ~K3 (or K4 ~K7), the vacant high bits will be regarded as " 0 ". Therefore, the operation is a positive-value one. The BCD value composed of $\mathrm{X} 4 \sim \mathrm{X} 13$ will be converted to BIN value and sent to DO.


5, You can choose any No. for bit devices, but please make the 1 s digit of X and Y " 0 " , e.g. $\mathrm{X0}, \mathrm{X} 10, \mathrm{X} 20, \cdots \mathrm{Y} 0$, $Y 10 \cdots$, and the 1 digit of $M$ and $S$ " 8 ' s multiple" (" 0 " is still the best choice), e.g. M0, M10, M2O $\cdots$.
6, Designating continuous device No.
Take data register D for example, continuous D refers to D0, D1, D2, D3, D4 $\cdots$.
For bit devices with specifically designated digit, continuous No. refers to:

| K1X0 | K1X4 | K1X10 | K1X14 $\ldots$ |
| :--- | :--- | :--- | :--- |
| K2Y0 | K2Y10 | K2Y20 | K2X30 $\ldots$ |
| K3M0 | K3M12 | K3M24 | K3M36. |
| K4S0 | K4S16 | K4S32 | K4S48 $\ldots$ |

Please follow the No. in the table and do not skip No. in case confusion may occur. In addition, if you use $K 4 Y O$ in the 32 -bit operation, the higher 16 bits will be regarded as " 0 ". For 32 -bit data, please use K8Y0.
The operations in PLC are conducted in BIN integers. When the integer performs division, e.g. $40 \div 3=13$ and the remainder is 1 . When the integer performs square root operations, the decimal point will be left out. Use decimal point operation instructions to obtain the decimal point.

## 7. Binary Floating Point

PLC represents floating points in 32 bits, following the IEEE754 standard:


The expressible size is:

$$
(-1)^{S} \times 2^{E-B} \times 1 . M, \text { in which } \mathrm{B}=127
$$

Therefore, the range for the 32-bit floating point is $\pm 2^{-126} \sim \pm 2^{+128}$, i.e. $\pm 1.1755 \times 10^{-38} \sim \pm 3.4028 \times 10^{+38}$ |
Example 1: Representing "23" in 32-bit floating point
Step 1: Convert "23" into a binary value: 23.0=10111
Step 2: Normalize the binary value: $10111=1.0111 \times 24$, in which 0111 is mantissa and 4 is exponent.
Step 3: Obtain the exponent:
$\because E-B=4 \rightarrow E-127=4$
$\therefore E=131=100000112$
Step 4: Combine the sign bit, exponent and mantissa into a floating point $01000001101110000000000000000000_{2}=41 \mathrm{~B}_{2} 0000_{16}$
Example 2: Representing "-23.0" in 32-bit floating point
The steps required are the same as those in Example 1. The only difference is you have to alter the sign bit into " 1 ". PLC uses registers of 2 continuous No. to combine into a 32-bit floating point. For example, we use registers (D1, D0) for storing a binary floating point as below:


8, Decimal Floating Point
> Since the binary floating point are not very user-friendly, we can convert it into a decimal floating point for use. Please be noted that the decimal point operation in PLC is still in binary floating point.
$>$ The decimal floating point is represented by 2 continuous registers. The register of smaller No. is for the constant。 while the register of bigger No. is for the exponent. Example: Storing a decimal floating point in registers (D1, D0)
Decimal floating point $=$ [constant D0] $\times 10$ [exponent D1 ]
Constant DO $= \pm 1,000 \sim \pm 9,999$
Exponent D1 $=-41 \sim+35$
The constant 100 does not exist in DO due to 100 is represented as $1,000 \times 10^{-1}$. The range of decimal floating point is $\pm 1175 \times 10^{-41} \sim \pm 3402 \times 10^{+35}$.
9. The decimal floating point can be used in the following instructions:

D EBCD: Converting binary floating point to decimal floating point
D EBIN: Converting decimal floating point to binary floating point

### 5.3 E, F Index Register Modification

- The index registers are 16-it registers. There are 16 points EO ~ E7 and FO ~ F7

1, $E$ and $F$ index registers are 16 -bit data registers. They can be read and written.
2, If you need a 32-bit register, you have to designate E. In this case, F will be covered up by E and cannot be used; otherwise, the contents in E may become incorrect. (We recommend you use MOVP instruction to reset the contents in D to 0 when the PLC is switched on.)
3. Combination of $E$ and $F$ when you designate a 32-bit index register: (E0, F0), (E1, F1), (E2, F2), ... (E7, F7) See the diagram in the left hand side. E, F index register modification refers to the content in the operand changes with the contents in E and F .
For example, E0 $=8$ and K20EO represents constant $\mathrm{K} 28(20+8)$. When the condition is true, constant K28 will be transmitted to register D24.


$$
\begin{aligned}
& \mathrm{E} 0=8 \quad \mathrm{~F} 0=14 \\
& 20+8=28 \quad 10+14=24 \\
& \text { Transmission } \mathrm{K} 28 \rightarrow \mathrm{D} 24
\end{aligned}
$$

3, Devices modifiable MPU: P, I, X, Y, M, S, K, H, KnX, KnY, KnM, KnS, T, C, D
If you need to modify device $P, I, X, Y, M, S, K n X, K n Y, K n M, K n S, T, C$ and $D$ by $E, F$, you have to select a 16-bit register, i.e. you can designate E or F. To modify constant K and H in a 32-bit instruction, you have to select a 32-bit register, i.e. you have to designate E.

## 6.1 (ZL 00-09) Program flow instructions



1. Explanations:

- $S$ : The destination pointer of conditional jump.
- When the user does not wish a particular part of PLC program in order to shorten the scan time and execute dual outputs, CJ instruction or CJP instruction can be adopted.。
- When the program designated by pointer $P$ is prior to CJ instruction, WDT timeout will occur and PLC will stop running. Please use it carefully.
- CJ instruction can designate the same pointer Prepeatedly. However, CJ and CALL cannot designate the same pointer $P$; otherwise an error will occur.
- Actions of all devices while conditional jumping is being executed.
a. $Y, M$ and $S$ remain their previous status before the conditional jump takes place.
b. Timer 10 ms and 100 ms that is executing stops.
c. Timer T192~T199 that execute the subroutine program will continue and the output contact executes normally.
d. The high-speed counter that is executing the counting continues counting and the output contact executes normally.
e. The ordinary counters stop executing.
f. If the "reset instruction" of the timer is executed before the conditional jump, the device will still be in the reset status while conditional jumping is being executed.
g. Ordinary application instructions are not executed.
h. The application instructions that are being executed, i.e. ZL 53 DHSCS, ZL 54 DHSCR, ZL 55 DHSZ, ZL 56 SPD, ZL 57 PLSY, ZL 58 PWM, ZL 59 PLSR, ZL 157 PLSV, ZL 158 DRVI, ZL 159 DRVA, continue being executed.

2, Program Example 1:

- When $\mathrm{XO}=\mathrm{On}$, the program automatically jumps from address 0 to N (the designated label P ) and keeps its execution. The addresses between 0 and N will not be executed.
- When XO = Off, as an ordinary program, the program keeps on executing from address 0 . CJ instruction will not be executed at this time.


3. Program Example 2:

- CJ instruction can be used in the following 5 conditions between $M C$ and $M C R$ instructions:
a) Without MC ~MCR.
b) From without $M C$ to within MC. Valid in the loop Pl as shown in the figure below.
c) In the same level $N$, inside of $M C \sim M C R$.
d) From within MC to without MCR.
e) Jumping from this MC $\sim M C R$ to another MC $\sim M C R$
(f) When MC instruction is executed, PLC will push the status of the switch contact into the self-defined stack in PLC. The stack will be controlled by the PLC, and the user cannot change it. When MCR instruction is executed, PLC will obtain the previous status of the switch contact from the top layer of the stack. Under the conditions as stated in b), d) and e), the times of pushing-in and obtaining stack may be different. In this case, the maximum stack available to be pushed in is 8 and the obtaining of stacks cannot resume once the stack becomes empty. Thus, when using CALL or CJ instructions, the user has to be aware of the pushing-in and obtaining of stacks.



## 4, Program Example 3:

The states of each device:

| Device | Contact state before CJ is executed | Contact state when CJ is being executed | Output coil state when CJ is being executed |
| :---: | :---: | :---: | :---: |
| Y, M, S | M1, M2, M3 Off | M1, M2, M3 Off $\rightarrow$ On | Y1**, M20, S1 Off |
|  | M1, M2, M3 On | M1, M2, M3 Onf $\rightarrow$ Off | Y1** M20, S1 On |
| 10, 100ms Timer | M4 Off | M4 Off $\rightarrow$ On | Timer TO is not enabled. |
|  | M4 On | M4 On $\rightarrow$ Off | Timer TO immediately stops and is latched. MO On $\rightarrow$ |

6 Application Instructions ZL 00-49

|  |  |  | Off, TO is reset as 0. |
| :---: | :---: | :---: | :---: |
| 10, 100ms Timer (accumulative) | M6 Off | M6 Off $\rightarrow$ On | Timer T240 is not enabled. |
|  | M6 On | M6 On $\rightarrow$ Off | Once the timer function is enabled and when met with CJ instruction, all accumulative timers will stop timing and stay latched. MO On $\rightarrow$ Off. T240 remains unchanged. |
| C0~C234 | M7 , M10 Off | M10 On/Off trigger | Counter does not count. |
|  | M7 Off, M10 On/Off trigger | M10 On/Off trigger | Counter C0 stops counting and stays latched. After M0 goes Off, C0 resumes its counting. |
| Application instruction | M11 Off | M11 Off $\rightarrow$ On | Application instructions are not executed. |
|  | M11 On | M11 On $\rightarrow$ Off | The skipped application instructions are not executed, butZL59, ZL157~159 keep being executed. |

*1: Y 1 is a dual output. When $M 0$ is Off, M 1 will control Y 1 . When M0 is On, M12 will control Y 1 .
*2:When the high-speed counters (C235 ~ C255) are driven and encounter the execution of CJ instruction, thecounting will resume, as well as the action of the output points.。


| $\begin{aligned} & \mathrm{ZL} \\ & 01 \end{aligned}$ |  | CALL |  |  |  |  | S |  | Call Subroutine |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit device |  |  |  | Word device |  |  |  |  |  |  |  |  |  |  | CALL, CALLP: 3 steps 16 bits |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| - Operand $S$ can designate $P$. <br> - P can be modified by index register E, F. |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- $\quad$ : pointer to call subroutine.
- Operand S can designate PO ~ P255.
- Edit the subroutine designated by the pointer after FEND instruction.
- The number of pointer $P$, when used by CALL, cannot be the same as the number designated by CJ instruction.
- If only CALL instruction is in use, it can call subroutines of the same pointer number with no limit on times.
- Subroutine can be nested for 5 levels including the initial CALL instruction. (If entering the sixth level, the subroutine won' † be executed.)


1, Explanations:

- The subroutine will return to main program by SRET after the termination of subroutine and execute the sequence program located at the next step to the CALL instruction.

2, Program Example 1:

- When $\mathrm{XO}=\mathrm{On}, \mathrm{CALL}$ instruction is executed and the program jumps to the subroutine designated by P2. When SRET instruction is executed, the program returns to address4 and continues its execution.
- There is no need to edit the FEND and SRET codes in the ladder diagram. After the compilation is passed, the instruction list will be automatically generated.



| $\begin{aligned} & \mathrm{ZL} \\ & 03 \end{aligned}$ |  |  | IRET |  |  |  |  |  | Interrupt Return |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit device |  |  |  | Word device |  |  |  |  |  |  |  |  |  |  |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F | IRET: 1 steps 16-bit |

1, Explanations:

- Interruption return refers to interrupt the subroutine.
- After the interruption is over, returning to the main program from IRET to execute the next instruction where the program was interrupted.



1, Explanations:

- El instruction allows interrupting subroutine in the program, e.g. external interruption, timed interruption, and high-speed counter interruption.
- In the program, using interruption subroutine between El and DI instruction is allowed. However, you can choose not to use $D$ instruction if there is no interruption-disabling section in the program.
- Pointer for interruption (I) must be placed after FEND instruction.
- Other interruptions are not allowed during the execution of interruption subroutine.
- When many interruptions occur, the priority is given to the firstly executed interruption. If several interruptions occur simultaneously, the priority is given to the interruption with the smaller pointer No.
- The interruption request occurring between DI and El instructions that cannot be executed immediately will be memorized and will be executed in the area allowed for interruption.
- When using the interruption pointer, DO NOT repeatedly use the high-speed counter driven by the same $X$ input contact.
- When immediate I/O is required during the interruption, write REF instruction in the program to update the status of $\mathrm{I} / \mathrm{O}$.


## 2, Program Example:

During the operation of PLC, when the program scans to the area between El and DI instructions and $\mathrm{X} 2=$ Off $\rightarrow$ On, interruption subroutine $A$ or $B$ will be executed. When the subroutine executes to IRET, the program will return to the main program and resumes its execution.


The corresponding command list is as follows:

| 11 | Instruction List $\times$ | Device Comment List | [P201] 12 |
| :--- | :--- | :--- | :--- |
| 000000 | EI |  |  |
| 000001 | LD | M0 |  |
| 000002 | OUT | Y0 |  |
| 000003 | FEND |  |  |
| 000004 | P201 |  |  |
| 000005 | LD | M1000 |  |
| 000006 | INC | D0 |  |
| 000009 | SRET |  |  |
| 000010 | END |  |  |

3, No. of interruption pointer I :
a. External interruptions: $(I 00 \square, X 0),(I 10 \square, X 1),(I 20 \square, X 2),(I 30 \square, X 3),(140 \square, X 4),(I 50 \square, X 5),(I 60 \square, X 6)$, $(170 \square, X 7),(190 \square, X 10),(191 \square, X 11),(192 \square, X 12),(193 \square, X 13),(194 \square, X 14),(195 \square, X 15),(196 \square, X 16),(197 \square$, X17) 16 points. ( $\square=0$ designates interruption in falling-edge, $\square=1$ designates interruption in risingedge).
b. Time interruptions: $16 \square \square, 17 \square \square$, 2 points. ( $\square \square=2 \sim 99 \mathrm{~ms}$, time base $=1 \mathrm{~ms}$ )

$$
\text { I8 } \square \square 1 \text { point. ( } \square \square=5 \sim 99 \mathrm{~ms} \text {, time base }=0.1 \mathrm{~ms} \text { ) }
$$

c. High-speed counter interruptions: IO10, IO20, I030, 1040, 1050, 10606 points. (used with ZL 53 DHSCS instruction to generate interruption signals)
d. When pulse output interruptions 1110,1120 (triggered when pulse output is finished), 1130,1140 (triggered when the first pulse output starts) are executed, the currently executed program is interrupted and jumps to the designated interruption subroutine.
e. Communication interruption: $1150,1151,1153,1160,1161$
f. The order for execution of interruption pointer I: external interruption, time interruption, high-speed counter interruption, pulse interruption, and communication interruption.

| $\begin{array}{l\|} \hline \mathrm{ZL} \\ 06 \end{array}$ | FEND |  |  |  |  |  |  |  | The End of The Main Program |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit device |  |  |  | Word device |  |  |  |  |  |  |  |  |  |  | FEND: 1 steps 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
|  | No contact to drive the instruction is required. |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- This instruction denotes the end of the main program. It has the same function as that of END instruction when being executed by PLC.
- CALL must be written after FEND instruction and add SRET instruction in the end of its subroutine. Interruption program has to be written after FEND instruction and IRET must be added in the end of the service program.
- If several FEND instructions are in use, place the subroutine and interruption service programs between the final FEND and END instruction
- After CALL instruction is executed, executing FEND before SRET will result in errors in the program.
- After FOR instruction is executed, executing FEND before NEXT will result in errors in the program.


1. Operands:
$S$ : The number of repeated nested loops

| $\begin{aligned} & \mathrm{ZL} \\ & 09 \end{aligned}$ |  | NEXT |  |  |  |  |  |  |  | End of a FOR-NEXT LOOP |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit device |  |  |  | Word device |  |  |  |  |  |  |  |  |  |  | NEXT: 1 steps 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
|  | No operand. No contact to drive the instruction is required |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- FOR instruction indicates FOR ~NEXT loops executing back and forth $N$ times before escZLng for the next execution.
- $\mathrm{N}=\mathrm{K} 1 \sim \mathrm{~K} 32,767$. N is regarded as K 1 when $\mathrm{N} \leqslant 1$.
- When FOR~NEXT loops are not executed, the user can use the CJ instruction to escape the loops.
- Error will occur when:

1) NEXT instruction is before FOR instruction.
2) FOR instruction exists but NEXT instruction does not exist.
3) There is NEXT instruction after FEND or END instruction.

2, The number of instructions between FOR ~NEXT differs.
Program Example 1:
After program $A$ has been executed for 3 times, it will resume its execution after NEXT instruction. Program $B$ will be executed for 4 times whenever program A is executed once. Therefore, program B will be executed $3 \times 4$ $=12$ times in total.

3. Program Example 2:

When X7 = Off, PLC will execute the program between FOR ~NEXT. When X7 = On, CJ instruction jumps to P6 and avoids executing the programs between FOR ~ NEXT.


4, Program Example 3:
When the programs between FOR ~ NEXT are not to be executed, the user can adopt CJ instruction for a jumping. When the most inner FOR ~ NEXT loop is in the status of XI = On, CJ instruction executes jumping to PO and skips the execution on PO.


## 6.2 (ZL 10-19) Transmission comparison

| ZL 10 |  | CMP |  |  |  |  | S 1 |  |  |  | D |  |  |  |  | Compare |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit device |  |  |  | Word device |  |  |  |  |  |  |  |  |  |  | CMP, CMPP: 7 steps 16-bit DCMP, DCMPP: 13 step |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| D |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| - If $S 1$ and $S 2$ are used in device $F$, only 16 -bit instruction is applicable. <br> - Operand D occupies 3 consecutive devices |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S1:Comparison Value 1 S2: Comparison Value 2 D: Comparison result.
- The contents in S1 and S2 are compared and the result will be stored in D.
- The two comparison values are compared algebraically and the two values are signed binary values. When b15 $=1$ in 16 -bit instruction or b31 $=1$ in 32-bit instruction, the comparison will regard the value as negative binary values.
- The designated device is YO , then $\mathrm{YO}, \mathrm{Y} 1$ and Y 2 are automatically occupied.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example:

When $\mathrm{X} 10=\mathrm{On}, \mathrm{CMP}$ instruction will be executed and one of $\mathrm{YO}, \mathrm{Y} 1$, and Y 2 will be On. When $\mathrm{X} 10=\mathrm{Off}$, CMP instruction will not be executed and $Y 0, Y 1$, and $Y 2$ remain their status before $\mathrm{X} 10=0 \mathrm{ff}$.

- If the user need to obtain a comparison result with $\geqslant \leqslant$, and $\neq$, make a series parallel connection between YO~Y2.

- To clear the comparison result, use RST or ZRST instruction.


| ZL 11 |  |  | ZCP |  |  |  | S |  | S2 |  |  |  | D |  |  | Zone Compare |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit device |  |  |  | Word device |  |  |  |  |  |  |  |  |  |  | ZCP, ZCPP: 9 steps 16-bit DZCP, DZCPP: 17 steps 32-bit |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |
| - If $\mathrm{S} 1, \mathrm{~S} 2$ and S are used in device F , only 16 -bit instruction is applicable <br> - The content in S 1 should be smaller than the content in S 2 <br> - Operand D occupies 3 consecutive devices |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations

- S1: Lower bound of zone comparison S2: Upper bound of zone comparison S : Comparison value
- $S$ is compared with its $\mathrm{S} 1, \mathrm{~S} 2$ and the result is stored in D .
- When $\mathrm{S} 1>\mathrm{S} 2$, the instruction performs comparison by using S 1 as the lower/upper bound.
- The two comparison values are compared algebraically and the two values are signed binary values. When b15 $=1$ in 16 -bit instruction or b31 $=1$ in 32-bit instruction, the comparison will regard the value as negative binary values.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example:

- Designate device MO, and operand D automatically occupies MO, M1 and M.
- When $X 0=O n, Z C P$ instruction will be executed and one of $M 0, M 1$, and $M 2$ will be On. When $X 0=O f f, Z C P$ instruction will not be executed and $M 0, M 1$, and $M 2$ remain their status before $\mathrm{XO}=$ Off.

- To clear the comparison result, use RST or ZRST instruction



1. Explanations:

- If $S$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- See the specifications of each model for their range of use.
- When this instruction is executed, the content of $S$ will be moved directly to $D$. When this instruction is not executed, the content of D remains unchanged.
- If the operation result refers to a 32-bit output, (i.e. application instruction MUL and so on), and the user needs to move the present value in the 32-bit high-speed counter, DMOV instruction has to be adopted.

2. Program Example:

- MOV instruction has to be adopted in the moving of 16 -bit data.
a) When $\mathrm{XO}=\mathrm{Off}$, the content in D10 will remain unchanged. If $\mathrm{XO}=\mathrm{On}$, the value K 10 will be moved to D10 data register.
b) When $\mathrm{XI}=$ Off, the content in D 10 will remain unchanged. If $\mathrm{XI}=\mathrm{On}$, the present value TO will be moved to D10 data register.
- DMOV instruction has to be adopted in the moving of 32-bit data.

When X2 = Off, the content in (D31, D30) and (D41, D40) will remain unchanged. If X2 $=$ On, the present value of (D21, D20) will be sent to (D31, D30) data register. Meanwhile, the present value of C235 will be moved to (D41, D40) data register.


| $\begin{aligned} & \mathrm{ZL} \\ & 13 \end{aligned}$ |  | SMOV |  |  |  |  |  |  |  | S | m1 | m2 |  | D |  | n |  |  | Shift Move |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  |  |  | SMOV: 11 steps 16-bit |
|  | X | Y | M | M | S | K | H |  | KnX | KnY | KnM | KnS | T | C |  | D | E | F |  |
| s |  |  |  |  |  |  |  |  | * | * | * | * | * | * |  | * | * | * |  |
| ml |  |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |
| m2 |  |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |
| D |  |  |  |  |  |  |  |  |  | * | * | * | * | * |  |  | * | * |  |
| n |  |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- s : Source of data ml : Start digit to be moved of the source data m 2 : Number of digits (nibbles) to be moved of the source data $D$ : Destination device $n$ : Start digit of the destination position for the moved digits
- This instruction is able to re-allocate or combine data. When the instruction is executed, m 2 digits of contents starting from digit ml (from high digit to low digit) of S will be sent to m 2 digits starting from digit n (from high digit to low digit) of D .
- Range: $\mathrm{ml}=1 \sim 4 ; \mathrm{m} 2=1 \sim \mathrm{ml} ; \mathrm{n}=\mathrm{m} 2 \sim 4$.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example 1:

- XO=On, specify the 4th digit of the decimal value of D10 (also known as the thousands digit) and start to transfer the contents of the 2 digits calculated from the low digit to the 3rd digit of the decimal value of D20 (that is, the hundreds digit) Digits) from the bottom 2 digits. The contents of 103 and 100 of D20 remain unchanged after this instruction is executed.
- When the BCD value exceeds the range of 0-9,999, the PLC judges it as an operation error and the instruction is not executed.

If D10=K1,234, D20=K5,678 before execution, after execution, D10 remains unchanged, D20=K5,128.

| xo |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| SMOV | D10 | K4 | K2 | D20 | K3 |



4, Program Example 2

1) Dip switches connected to the input terminal of non-sequential numbers can be synthesized using this command.
2) Transfer the right 2 digits of the DIP switch to the right 2 digits of D2, and the left 1 digit of the DIP switch to the right 1 digit of $D 1$.
3) Use the SMOV instruction to transfer the first digit of $D 1$ to the third digit of $D 2$ to combine the two sets of DIP switches into one group.


| $\begin{aligned} & \hline \text { ZL } \\ & 14 \end{aligned}$ |  | CML |  |  |  |  |  |  |  | D |  |  |  |  |  | Compliment |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | CML: 5 steps 16 -bit DCML: 9 steps 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |

1. Explanations:

- S:Source of data D: Destination device.
- This instruction can be used for phase-reversed output.
- Reverse the phase $(0 \rightarrow 1,1 \rightarrow 0)$ of all the contents in $S$ and send the contents to D . Given that the content is a constant $K$, $K$ will be automatically converted into a BIN value.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2. Program Example 1:

When X10 = On, b0 ~ b3 in D1 will be phase-reversed and send to YO ~ Y3.

3. Program Example 2:

The loop below can also adopt CML instruction (see right below).


| ZL 15 |  | BMOV |  |  |  |  | $S$ D |  |  |  | n |  |  |  |  | Block Move |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | BMOV: 7 step | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  | * | * | * |  |  |  |  |

1. Explanations:

- S:Start of source devices D:Start of destination devices n: Number of data to be moved
- Range of $\mathrm{n}: 1 \sim 512$
- The contents in n registers starting from the device designated by S will be moved to n registers starting from the device designated by D. If $n$ exceeds the actual number of available source devices, only the devices that fall within the valid range will be used.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example 1:
When X10 = On, the contents in registers D0 ~ D3 will be moved to the 4 registers D20 ~ D23.


3, Program Example 2:
Assume the bit devices KnX, KnY, KnM and KnS are designated for moving, the number of digits of S and $D$ has to be the same, i.e. their n has to be the same.


4, Program Example 3:

To avoid coincidence of the device numbers to be moved designated by the two operands and cause confusion, please be aware of the arrangement on the designated device numbers:
a. When S > D, the instruction is processed following the order: $1 \rightarrow 2 \rightarrow 3$

b. When $S<D$, the instruction is processed following the order: $3 \rightarrow 2 \rightarrow 1$


| $\begin{array}{\|l\|} \hline \mathrm{ZL} \\ 16 \end{array}$ |  | FMOV |  |  |  |  | $S$ D |  |  |  |  |  |  |  |  | Fill Move |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | FMOV: 7 steps 16-bit DFMOV : 13 steps 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS |  |  | T | C | D |  | E | F |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  | * | * | * |  |  |  |

1. Explanations

- S: Source of data D: Destination of data n: Number of data to be moved
- If $S$ is used in device $F$, only 16 -bit instruction is applicable
- The contents in n registers starting from the device designated by $S$ will be moved to n registers starting from the device designated by D. If $n$ exceeds the actual number of available source devices, only the devices that fall within the valid range will be used
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example:

When $\mathrm{X10}=\mathrm{On}$, K10 will be moved to the 5 consecutive registers starting from D10.



1. Explanations:

- D1: Data to be exchanged 1. D2: Data to be exchanged 2
- If D1 and D2 are used in device F, only 16-bit instruction is applicable.
- The contents in the devices designated by D1 and D2 will exchange.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example 1:
When X0 $=\mathrm{Off} \rightarrow \mathrm{On}$, the contents in D20 and D40 exchange with each other.

3. Program Example 2:

When XO = Off $\rightarrow$ On, the contents in D100 and D200 exchange with each other.



1, Explanations:

- S:Source of data D: Conversion result.
- If $S$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- The four arithmetic operations and applications in PLC and the execution of INC and DEC instructions are performed in BIN format. Therefore, if the user needs to see the decimal value display, simply use this instruction to convert the BIN value into BCD value.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example:

When $\mathrm{XO}=\mathrm{On}$, the binary value of D 10 will be converted into $B C D$ value, and the 1 s digit of the conversionresult will be stored in KIYO (YO ~ Y3, the 4 bit devices).

When D10 $=001 \mathrm{E}$ (hex) $=0030$ (decimal), the execution result will be: Y0 $\sim$ Y3 $=0000$ (BIN).



1, Explanations:

- S: Source of data D: Conversion result.
- If $S$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- The four arithmetic operations and applications in PLC and the execution of INC and DEC instructions are performed in BIN format. Therefore, if the user needs to see the decimal value display, simply use this instruction to convert the BIN value into BCD value.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example:

When $\mathrm{XO}=\mathrm{On}$, the binary value of D10 will be converted into BCD value, and the 1 s digit of the conversion result will be stored in K1YO (YO ~ Y3, the 4 bit devices).
When D10 $=001 \mathrm{E}$ (hex) $=0030$ (decimal), the execution result will be: YO $\sim$ Y3 $=0000$ (BIN).

| X0 | BIN | K1M0 |
| :--- | :--- | :--- |
| D10 |  |  |

Note: a. BCD and BIN instruction application instructions:

1) When the PLC wants to read a BCD type DIP switch from the outside, it must use the BIN command to convert the read data into a BIN value before storing it in the PLC.
2) When the PLC wants to display the internally stored data through an external BCD-type 7-segment display, it must use the BCD command to first convert the internal data to be displayed into BCD values and then send it to the 7 -segment display.
3) When $\mathrm{XO}=\mathrm{On}$, convert the K $4 X O B C D$ value to BIN value and transfer it to D 100 , then convert the BIN value of D 100 to $B C D$ value and transfer it to K4Y20.



## 6.3 (ZL 20-29) Four logical operations

| ZL <br> 20 |  | ADD |  |  |  |  | S1 S2 D |  |  |  |  |  |  |  |  | BIN Addition |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | ADD: 7 steps $\quad 16$-bitDADD: 13 steps $\quad 32$-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |  |

1. Explanations:

- S1:Summand S2: Addend D:Sum
- If $S 1, S 2$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- This instruction adds S 1 and S 2 in BIN format and store the result in D.
- The highest bit is symbolic bit $0(+)$ and $1(-)$, which is suitable for algebraic addition, e.g. $3+(-9)=-6$
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example 1:
In 16-bit BIN addition: When X0 $=$ On, the content in D0 will plus the content in D10 and the sum will be stored in D20.

3. Program Example 2:

In 32-bit BIN addition: When $\mathrm{X10}=\mathrm{On}$, the content in (D31, D30) will plus the content in (D41, D40) and the sum will be stored in (D51, D50). D30, D40 and D50 are low 16-bit data; D31, D41 and D51 are high 16-bit data.

| X10 | DADD | D30 | D40 | D50 |
| :--- | :--- | :--- | :--- | :--- |

$$
(\mathrm{D} 31, \mathrm{D} 30)+(\mathrm{D} 41, \mathrm{D} 40)=(\mathrm{D} 51, \mathrm{D} 50)
$$



1, Explanations:

- S1:Minuend S2: Subtrahend D: Remainder
- If $S 1, S 2$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- This instruction subtracts S1 and S2 in BIN format and stores the result in D.
- The highest bit is symbolic bit $0(+)$ and $1(-)$, which is suitable for algebraic subtraction
- For flag operations of SUB instruction and the positive/negative sign of the value, see the explanations in ADD instruction on the previous page.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2. Program Example 1:

In 16-bit BIN subtraction: When X0 $=$ On, the content in DO will minus the content in D10 and the remainder will be stored in D20.

3. Program Example 2:

In 32-bit BIN subtraction: When $\mathrm{X1}=\mathrm{On}$, the content in (D31, D30) will minus the content in (D41, D40) and the remainder will be stored in (D51, D50). D30, D40 and D50 are low 16-bit data; D31, D41 and D51 are high 16-bit data.



1, Explanations:

- S1:Multiplicand S2:Multiplicator D: Product
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- If $S 1$ and $S 2$ are used in device F, only 16 -bit instruction is applicable.
- If $D$ is used in device $E$, only 16 -bit instruction is applicable
- In 16-bit instruction, D occupies 2 consecutive devices.
- In 32-bit instruction, D occupies 4 consecutive devices.
- This instruction multiplies $S 1$ by $S 2$ in BIN format and stores the result in D. Be careful with the positive/negative signs of $\mathrm{S} 1, \mathrm{~S} 2$ and D when doing 16 -bit and 32 -bit operations.
- In 16-bit BIN multiplication:


Symbol bit $=0$ refers to a positive value .
Symbol bit = 1 refers to a negative value.
16 -bit value $\times 16$-bit value $=32$-bit value
When D serves as a bit device, it can designate K1 ~ K4 and construct a 16-bit result, occupying consecutive 2 groups of 16 -bit data.

If the product of a 16 -bit multiplication must be a 16 -bit value ( 16 -bit value $\times 16$-bit value $=16$-bit value), users have to use ZL 114 MULI6/MUL16P. Please refer to the explanation of ZL 114 MULI6/MUL16P for more information.

- 32-bit BIN multiplication:


Symbol bit $=0$ refers to a positive value. Symbol bit = 1 refers to a negative value.

When $D$ serves as a bit device, it can designate $K 1 \sim K 8$ and construct a 32-bit result, occupying consecutive 2 groups of 32-bit data.

2, Program Example:
The 16 -bit D0 is multiplied by the 16 -bit D10 and brings forth a 32 -bit product. The higher 16 bits are stored in D21 and the lower 16-bit are stored in D20. On/Off of the most left bit indicates the positive/negative status of the result value.


| $\begin{aligned} & \mathrm{ZL} \\ & 23 \end{aligned}$ |  | DIV |  |  |  |  | S 1 S2 D |  |  |  |  |  |  |  |  | BIN Division |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DIV: 7 steps $\quad 16$-bit DDIV: 13 steps 32-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |

1. Explanations:

- S1:Dividend
S2: Divisor
D: Quotient and remainder
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- If $S 1$ and $S 2$ are used in device $F$, only 16 -bit instruction is applicable.
- If $D$ is used in device $E$, only 16 -bit instruction is applicable.
- In 16-bit instruction, D occupies 2 consecutive devices.
- In 32-bit instruction, D occupies 4 consecutive devices.
- This instruction divides S1 and S2 in BIN format and stores the result in D. Be careful with the positive/negative signs of $S 1, S 2$ and $D$ when doing 16 -bit and 32-bit operations.
- This instruction will not be executed when the divisor is 0 .
- In 16-bit BIN division:


When $D$ serves as a bit device, it can designate $\mathrm{K} 1 \sim \mathrm{~K} 4$ and construct a 16 -bit result, occupying consecutive 2 groups of 16 -bit data and bringing forth the quotient and remainder.

- In 32-bit BIN division:

Quotient Remainder


When D serves as a bit device, it can designate K1 ~ K8 and construct a 32-bit result, occupying consecutive 2 groups of 32 -bit data and bringing forth the quotient and remainder.

2, Program Example:
When X0 = On, D0 will be divided by D10 and the quotient will be stored in D20 and remainder in D21. On/Off of the highest bit indicates the positive/negative status of the result value.

6 Application Instructions ZL 00-49

| DIV | D0 | D10 | D20 |
| :--- | :--- | :--- | :--- |
|  | DIV | D0 | D10 |


| ZL 24 | D | INC |  |  |  |  | D |  |  |  |  |  |  |  |  | Increment |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | INC:3steps 16-bit DINC: 5 steps 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |

1, Explanations:

- D: Destination device
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- If $D$ is used in device $F$, only 16 -bit instruction is applicable.
- If the instruction is not a pulse execution one, the content in the designated device $D$ will plus " 1 " in every scan period whenever the instruction is executed.
- This instruction adopts pulse execution instructions (INCP, DINCP).
- In 16-bit operation, 32,767 pluses 1 and obtains -32,768. In 32-bit operation, 2,147,483,647 pluses 1 and obtains $-2,147,483,648$.

2. Program Example:

When XO = Off $\rightarrow$ On, the content in DO pluses 1 automatically


| ZL 25 | D | DEC |  |  |  |  | D |  |  |  |  |  |  |  |  | Decrement |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DEC:3steps 16-bit DDEC: 5 steps 32-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |

1, Explanations:

- D: Destination device
- If $D$ is used in device F, only 16-bit instruction is applicable.
- If the instruction is not a pulse execution one, the content in the designated device D will minus "1" in every scan period whenever the instruction is executed.
- This instruction adopts pulse execution instructions (DECP, DDECP).
- In 16-bit operation, -32,768 minuses 1 and obtains 32,767. In 32-bit operation, -2,147,483,648 minuses 1 and obtains 2,147,483,647.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example:
When XO = Off $\rightarrow$ On, the content in DO minuses 1 automatically.
If the DEC instruction needs to be executed once, add the rising edge " $\uparrow$ " of NP before DEC.



1, Explanations:

- S1:Source data device 1 S2: Source data device 2 D: Operation result
- If $S 1, S 2$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- This instruction conducts logical AND operation of S1 and S2 and stores the result in D.
- Operation rule: The corresponding bit of the operation result in $D$ will be " 0 " if any of the bits in S 1 or S 2 is " 0 ".
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2. Program Example 1:

When X0 = On, the 16-bit D0 and D2 will perform WAND, logical AND operation, and the result will be stored in D4.

3. Program Example 2:

When $\mathrm{X1}$ = On, the 32-bit (D1 1, D10) and (D21, D20) will perform DAND, logical AND operation, and the result will be stored in (D41, D40).



1. Explanations:

- S1:Source data device 1 S2: Source data device 2 D: Operation result
- If $S 1, S 2$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- This instruction conducts logical OR operation of S1 and S2 and stores the result in D.
- Operation rule: The corresponding bit of the operation result in $D$ will be "1" if any of the bits in S1 or S2 is "1".
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example 1:

When X0 = On, the 16-bit D0 and D2 will perform WOR, logical OR operation, and the result will be stored in D4.


## 3. Program Example 2:

When $\mathrm{X} 1=$ On, the 32-bit (D11, D10) and (D21, D20) will perform DOR, logical OR operation, and the result will be stored in (D41, D40).

| D1 | DOR | D10 | D20 | D40 |
| :--- | :--- | :--- | :--- | :--- |



| $\begin{array}{l\|} \mathrm{ZL} \\ 28 \end{array}$ | W | XOR |  |  |  |  | S |  |  |  |  | D |  |  |  | Logical Exclusive OR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | WXOR: 7 steps 16 -bit DXOR: 13 steps 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  |  | KnM | KnS | T | C | D |  | E | F |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |

1, Explanations:

- S1:Source data device 1 S2: Source data device 2 D: Operation result
- If $S 1, S 2$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- This instruction conducts logical XOR operation of S1 and S2 and stores the result in D.
- Operation rule: If the bits in S1 and S2 are the same, the corresponding bit of the operation result in D will be " 0 " ; if the bits in S 1 and S 2 are different, the corresponding bit of the operation result in D will be " 1 ".
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example 1:

When X0 = On, the 16-bit D0 and D2 will perform WXOR, logical XOR operation, and the result will be stored in D4.


## 3, Program Example 2:

When $\mathrm{X1}$ = On, the 32-bit (D11, D10) and (D21, D20) will perform DXOR, logical XOR operation, and the result will be stored in (D41, D40).

| D1 | DXOR | D10 | D20 | D40 |
| :--- | :--- | :--- | :--- | :--- |


| b31 |  |  |  |  |  |  |  |  |  |  |  |  |  |  | b15 |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | (S1) D11 D10 | 11 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 1 | 11 | 1 | 1 | 1 | , | 1 | 1 | 1 | 0 | , | 0 | 1 |  | 1 |
| Before | DXOR |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | (S2) D21 D20 | $0 \mid 0$ | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 010 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 |  | 0 |


| $\begin{array}{c}\text { After } \\ \text { execution }\end{array}$ (D) D41 D40 |
| :--- |
| 1 |$\frac{1}{1}|1|$



1. Explanations:

- D: Device to store 2' s complement.
- If $D$ is used in device $F$, only 16 -bit instruction is applicable
- This instruction converts a negative BIN value into an absolute value
- This instruction adopts pulse execution instructions (NEGP, DNEGP).
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example 1:

When X0 $=$ Off $\rightarrow$ On, the phase of every bit of the content in D10 will be reversed $(0 \rightarrow 1,1 \rightarrow 0)$ and pluses 1. The result will then be stored in D10.

3. Program Example 2:

- Obtaining the absolute value of a negative value:
a) When the $15^{\text {th }}$ bit of DO is " $1^{\text {" }}, \mathrm{M} 0=\mathrm{On}$. ( D 0 is a negative value).
b) When $\mathrm{M} 0=\mathrm{Off} \rightarrow \mathrm{On}$, NEG instruction will obtain 2 's complement of D 0 and further its absolute value.


4. Program Example 3:

- Obtaining the absolute value by the remainder of the subtraction. When $\mathrm{XO}=\mathrm{On}$ :
a) If $D 0>D 2, M 0=O n$.
b) If $\mathrm{DO}=\mathrm{D} 2, \mathrm{Ml}=\mathrm{On}$.
c) If $\mathrm{DO}<\mathrm{D} 2, \mathrm{M} 2=\mathrm{On}$.
d) $D 4$ is then able to remain positive


Remarks:

- Negative value and its absolute value
a. The sign of a value is indicated by the highest (most left) bit in the register. 0 indicates that the value is a positive one and 1 indicates that the value is a negative one。
b. NEG instruction is able to convert a negative value into its absolute value.

```
(DO=2)
```



```
(D0=1)
```



```
(DO=0)
```




| ( $\mathrm{DO}=32,765$ ) |  |
| :---: | :---: |
|  |  |
| $(\mathrm{DO}=32,766)$ |  |
|  |  |
| $(\mathrm{D} 0=32,767) \quad \overline{\mathrm{D}})+1=32,767$ |  |
|  |  |
| $(\mathrm{DO}=32,768) \quad$ ( $\overline{\mathrm{DO}})+1=32,768$ |  |
|  |  |
|  | $x$ absolute value is 32,767 |

## 6.4 (ZL 30-39) Cyclic shift



1. Explanations:

- D: Device to be rotated n : Number of bits to be rotated in 1 rotation
- If $D$ is used in device $F$, only 16 -bit instruction is applicable.
- If D is designated as KnY, KnM, and KnS, only K4 (16-bit) and K8 (32-bit) are valid.
- Range of n: K1 ~K16 (16-bit); K1 ~ K32 (32-bit)
- This instruction rotates the device content designated by D to the right for n bits.
- This instruction adopts pulse execution instructions (RORP, DRORP)
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example

When $\mathrm{XO}=\mathrm{Off} \rightarrow$ On, the 16 bits ( 4 bits as a group) in D 10 will rotate to the right, as shown in the figure below.



1. Explanations:

- D: Device to be rotated n : Number of bits to be rotated in 1 rotation
- If $D$ is used in device $F$, only 16 -bit instruction is applicable.
- If D is designated as KnY , KnM , and KnS, only K4 (16-bit) and K8 (32-bit) are valid.
- Range of n : K1 ~K16 (16-bit); K1~K32 (32-bit).
- This instruction rotates the device content designated by D to the left for n bits.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example:

When $\mathrm{XO}=\mathrm{Off} \rightarrow$ On, the 16 bits ( 4 bits as a group) in D10 will rotate to the left, as shown in the figure below


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| $\begin{aligned} & \mathrm{ZL} \\ & 32 \end{aligned}$ |  | RCR |  |  |  |  | D |  |  | n |  |  |  |  |  | Rotation Right with Carry |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | RCR: 5 steps DRCR: 9 step | $\begin{aligned} & \text { 16-bit } \\ & \text { 32-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- D: Device to be rotated $n$ : Number of bits to be rotated in 1 rotation
- If $D$ is used in device $F$, only 16 -bit instruction is applicable.
- If D is designated as KnY , KnM , and KnS, only K4 (16-bit) and K8 (32-bit) are valid.
- Range of n : K1 ~K16 (16-bit); K1~K32 (32-bit).
- This instruction rotates the device content designated by $D$ to the right for n bits.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example:

When $\mathrm{XO}=\mathrm{Off} \rightarrow$ On, the 16 bits ( 4 bits as a group) in D10 will rotate to the right, as shown in the figure below。



1. Explanations:

- D: Device to be rotated $n$ : Number of bits to be rotated in 1 rotation
- If $D$ is used in device $F$, only 16 -bit instruction is applicable
- If D is designated as KnY , KnM , and KnS , only K 4 (16-bit) and K 8 (32-bit) are valid
- Range of n: K1~K16 (16-bit); K1~K32 (32-bit)
- This instruction rotates the device content designated by $D$ to the left for $n$ bits
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2. Program Example:

When $\mathrm{XO}=\mathrm{Off} \rightarrow$ On, the 16 bits ( 4 bits as a group) in D10 will rotate to the left, as shown in the figure below


16 bits
higher bit to the left lower bit



1, Explanations:

- S: Start No. of the shifted device D: Start No. of the device to be shifted
- nl: Length of data to be shifted n2: Number of bits to be shifted in 1 shift
- Range of nl: 1~ 1,024
- Range of $\mathrm{n} 2: 1 \sim \mathrm{n} 1$
- This instruction shifts the bit device of nl bits (desired length for shifted register) starting from D to the right for n 2 bits. S is shifted into D for n 2 bits to supplement empty bits.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example:

When XO = Off $\rightarrow$ On, MO ~M15 will form 16 bits and shifts to the right ( 4 bits as a group).
The figure below illustrates the right shift of the bits in one scan.
0 M3 ~M0 $\rightarrow$ carry
(2 M7~M4 $\rightarrow$ M3~M0
© M11~M8 $\rightarrow$ M7~M4
© M15 ~M12 $\rightarrow \mathrm{M} 11 \sim$ MB
© $\mathrm{X} 3 \sim \mathrm{X} 0 \rightarrow \mathrm{M} 15 \sim \mathrm{M} 12$ completed


4 bits as a group shifting to the right


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| $\begin{aligned} & \hline \mathrm{ZL} \\ & 35 \end{aligned}$ |  | SFTL |  |  |  |  |  |  | D | $\mathrm{n}_{1}$ |  | $\mathrm{n}_{2}$ |  |  |  | Bit Shift Left |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | SFFL: 9 step | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S | * | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{n}_{1}$ |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{n}_{2}$ |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S:Start No. of the shifted device D: Start No. of the device to be shifted
- n1: Length of data to be shifted n2: Number of bits to be shifted in 1 shift
- Range of n1: 1~ 1,024
- Range of $n 2: 1 \sim n 1$
- This instruction shifts the bit device of $n 1$ bits (desired length for shifted register) starting from D to the left for n 2 bits. S is shifted into D for n 2 bits to supplement empty bits
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example:

- When $\mathrm{X} 0=\mathrm{Off} \rightarrow \mathrm{On}, \mathrm{M0} \sim \mathrm{M} 15$ will form 16 bits and shifts to the left (4 bits as a group).
- The figure below illustrates the left shift of the bits in one scan.

> © M15~M12 $\rightarrow$ carry
> (3) M11~M8 $\rightarrow$ M15~M12
> © M7~M4 $\rightarrow$ M11~M8
> © M3 ~M0 $\rightarrow$ M7~M4
> © $\mathrm{X} 3 \sim \mathrm{X} 0 \rightarrow \mathrm{M} 3 \sim \mathrm{M} 0$ completed


| $\begin{aligned} & \mathrm{ZL} \\ & 36 \end{aligned}$ |  | WSFR |  |  |  |  |  |  | D | $\mathrm{n}_{1}$ |  | $\mathrm{n}_{2}$ |  |  |  | Word Shift Left |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | WSFR: 9 step | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| $\mathrm{n}_{1}$ |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{n}_{2}$ |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S: Start No. of the shifted device
- n 1 : Length of data to be shifted

D: Start No. of the device to be shifted
n 2 : Number of words to be shifted in 1 shift

- The type of devices designated by S and D has to be the same, e.g. KnX, KnY, KnM, and KnS as a category and $T, C$, and $D$ as another category
- Provided the devices designated by $S$ and $D$ belong to $K n$ type, the number of digits of $K n$ has to be the same
- Range of n1: 1~512
- Range of $\mathrm{n} 2: 1 \sim \mathrm{n} 1$
- This instruction shifts the stack data of $n 1$ words starting from $D$ to the right for $n 2$ words. $S$ is shifted into $D$ for n 2 words to supplement empty words.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example 1:

- When $\mathrm{XO}=\mathrm{Off} \rightarrow \mathrm{On}$, the 16 register stack data composed of D20~D35 will shift to the right for 4 registers.
- The figure below illustrates the right shift of the words in one scan.

```
(1) D23 ~ D20 \(\rightarrow\) carry
(2) \(\mathrm{D} 27 \sim \mathrm{D} 24 \rightarrow \mathrm{D} 23 \sim \mathrm{D} 20\)
(3) D31~D28 \(\rightarrow\) D27~D24
© D35 ~ D32 \(\rightarrow\) D31~D28
© D13 ~ D10 \(\rightarrow\) D35 ~ D32 completed
```

| X0 | WSFR | D10 | D20 | K16 |
| :--- | :--- | :--- | :--- | :--- |


(4)
(3)
(2)
(1)

## 3, Program Example 2:

- When $\mathrm{XO}=\mathrm{Off} \rightarrow \mathrm{On}$, the bit register stack data composed of $\mathrm{Y} 10 \sim \mathrm{Y} 27$ will shift to the right for 2 digits.
- The figure below illustrates the right shift of the words in one scan.
(1) $\mathrm{Y} 17 \sim \mathrm{Y} 10 \rightarrow$ carry
(2) $\mathrm{Y} 27 \sim \mathrm{Y} 20 \rightarrow \mathrm{Y} 17 \sim \mathrm{Y} 10$
(3) $\mathrm{X} 27 \sim \mathrm{X} 20 \rightarrow \mathrm{Y} 27 \sim \mathrm{Y} 20$ completed

When using Kn type device, please designate the same number of digits.

(3)

(2)
(1)

| ZL 37 |  | WSFL |  |  |  |  |  |  | D | $\mathrm{n}_{1}$ |  |  | $\mathrm{n}_{2}$ |  |  | Word Shift Left |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | WSFL: 9 step | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| $\mathrm{n}_{1}$ |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{n}_{2}$ |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- $S$ : Start No. of the shifted device D: Start No. of the device to be shifted
- n1: Length of data to be shifted n2: Number of words to be shifted in 1 shift
- The type of devices designated by S and D has to be the same, e.g. KnX, KnY, KnM, and KnS as a category and $T, C$, and $D$ as another category
- Provided the devices designated by S and D belong to Kn type, the number of digits of Kn has to be the same.
- Range of nl: 1~512
- Range of $n 2: 1 \sim n 1$
- This instruction shifts the stack data of n 1 words starting from D to the left for n 2 words. S is shifted into D for n 2 words to supplement empty words.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example:

- When $\mathrm{XO}=\mathrm{Off} \rightarrow \mathrm{On}$, the 16 register stack data composed of D20~D35 will shift to the left for 4 registers.
- The figure below illustrates the left shift of the words in one scan.

$$
\begin{aligned}
& \text { © D35 ~ D32 } \rightarrow \text { carry } \\
& \text { (2 D31 ~D28 } \rightarrow \text { D35 ~D32 } \\
& \text { © D27 ~D24 } \rightarrow \text { D31~D28 } \\
& \text { © D23 ~D20 } \rightarrow \text { D27 ~D24 } \\
& \boldsymbol{\Theta} \text { D13 ~D10 } \rightarrow \text { D23 ~D20 completed }
\end{aligned}
$$

| x0 | WSFL | D10 | D20 | K16 |
| :--- | :--- | :--- | :--- | :--- |



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| ZL <br> 38 |  | SFWR |  |  |  |  |  |  | D |  | n |  |  |  |  | Shift Register Write |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | SFWR: 7 steps 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- $S$ : Device of stack data written in $\quad D$ : Start No. of stack data $n$ : Length of stack data
- Range of $\mathrm{n}: 2 \sim 512$
- The stack data of $n$ words starting from D are defined as "first-in, first-out" stack data and designate the first device as the pointer. When the instruction is executed, the content in the pointer pluses 1 , and the content in the device designated by $S$ will be written into the designated location in the "first-in, first-out" stack data designated by the pointer. When the content in the pointer exceeds $n-1$, this instruction will not process any new value written.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example:

- Pointer D0 is reset as 0 . When $\mathrm{XO}=\mathrm{Off} \rightarrow \mathrm{On}$, the content in D20 will be sent to D 1 and the content in pointer D0 becomes 1. After the content in D20 is changed, make X0 $=$ Off $\rightarrow$ On again, and the content in D2 will be sent to D2 and the content in D0 becomes 2 .
- The figure below illustrates the shift and writing in 1~2 execution of the instruction.
- The content in D20 is sent to D1.
(2) The content in pointer D0 becomes 1 .


| $\begin{aligned} & \mathrm{ZL} \\ & 39 \end{aligned}$ |  | SFRD |  |  |  |  |  |  | D |  | n |  |  |  |  | Shift Register Read |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | SFRD: 7 steps 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S:Start No. of stack data D: Device of stack data read out $n$ : Length of stack data
- Range of $\mathrm{n}: 2 \sim 512$
- The stack data of n words starting from S are defined as "first-in, first-out" stack data. After the content in $S$ minuses 1 , the content in the device designated by $(S+1)$ will be written into the location designated by $D$, and $(S+n-1) \sim(S+2)$ will all right shift for one register while the content in $(S+n-1)$ remains the same. When the content in $S$ equals 0 , this instruction will not process any new value read out.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example:

- When $\mathrm{XO}=\mathrm{Off} \rightarrow \mathrm{On}$, the content in D1 will be sent to D21 and D9~D2 will shift to the right for 1 register (content in D9 remains unchanged) and the content in D0 minus 1 .
- The figure below illustrates the shift and reading in 1~3 execution of the instruction
- The instruction executes a shift read operation according to the following numbers 1~3.
© The content in D1 is sent to D21.
(2 D9 ~ D2 shift to the right for 1 register.
(s) The content in D0 minuses 1 .

| SO | SFRD | D0 | D21 | K10 |
| :---: | :---: | :---: | :---: | :---: |



## 6.5 (ZL40-49) Data processing



1, Explanations:

- D1: Start device of the range to be reset D2: End device of the range to be reset
- No. of operand D1 $\leqslant$ No. of operand D2.
- D1 and D2 have to designate devices of the same type.
- When the instruction is executed, area from D1 to D2 will be cleared.
- 16-bit counter and 32-bit counter can use ZRST instruction together.
- When D1 > D2, only operands designated by D2 will be reset.

2, Program Example:

- When X0 = On, auxiliary relays M300 ~ M399 will be reset to Off.
- When $\mathrm{X} 1=\mathrm{On}, 16$ counters $\mathrm{CO} \sim \mathrm{C} 127$ will all be reset (writing in 0 ; contact and coil being reset to Off).
- When X10 $=$ On, timers T0 $\sim$ T127 will all be reset (writing in 0 ; contact and coil being reset to Off).
- When $\mathrm{X} 2=$ On, steps $S 0 \sim S 127$ will be reset to Off.
- When $\mathrm{X} 3=$ On, data registers D0 ~D100 will be reset to 0 .
- When X4 = On, 32-bit counters C235 ~ C254 will all be reset. (writing in 0; contact and coil being reset to Off


Remarks:

1) Devices, e.g. bit devices $Y, M, S$ and word devices $T, C, D$, can use RST instruction.
 reset.


| $\begin{array}{\|l\|} \hline \mathrm{ZL} \\ 41 \end{array}$ |  | DECO |  |  |  |  |  |  |  | $S$ D |  |  | n |  |  |  | Decode |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DECO: 7 step | 16-bit |  |
|  | X | Y | M |  | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| S | * | * | * |  | * | * | * |  |  |  |  | * | * | * | * | * |  |  |  |
| D |  | * | * |  | * |  |  |  |  |  |  | * | * | * | * | * |  |  |  |
| n |  |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- $\quad$ : Source device to be decoded D: Device for storing the decoded result n: Length of decoded bits
- Range of $n$ when $D$ is a bit device: $1 \sim 8$
- Range of n when D is a word device: $1 \sim 4$
- See the specifications of each model for their range of use.
- The lower " n " bits of S are decoded and the results of " 2 " bits are stored in D .
- This instruction adopts pulse execution instructions (DECOP)
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2, Program Example 1:

- When D is used as a bit device, $\mathrm{n}=1 \sim 8$. Errors will occur if $\mathrm{n}=0$ or $\mathrm{n}>8$.
- When $\mathrm{n}=8$, the maximum points to decode is $28=256$ points. (Please be aware of the storage range of the devices after the decoding and do not use the devices repeatedly.)
- When X10 $=$ Off $\rightarrow$ On, this instruction will decode the content in X0 $\sim$ X2 to M100~M107.
- When the source of data is $1+2=3$, set M103, the 3rd bit starting from M100, as 1 .
- After the execution of this instruction is completed and X10 turns to Off, the content that has been decoded and output keeps acting.



## 3. Program Example 2:

- When $D$ is used as a word device, $n=1 \sim 4$. Errors will occur if $n=0$ or $n>4$.
- When $n=4$, the maximum points to decode is $2^{4}=16$ points.
- When X10 $=$ Off $\rightarrow$ On, this instruction will decode b2 ~b0 in D10 to b7 ~b0 in D20. b15 ~b8 that have not been used in D20 will all become 0 .
- The lower 3 bits of D10 are decoded and stored in the lower 8 bits of D20. The higher 8 bits of D20 are all 0.
- After the execution of this instruction is completed and X10 turns to Off, the content that has been decoded and output keeps acting.


| ZL <br> 42 |  | ENCO |  |  |  |  |  |  | $S$ D |  | n |  |  |  |  | Encode |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | + De | vice |  | Word Devices |  |  |  |  |  |  |  |  |  |  | ENCO: 7 step | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |  |
| S | * | * | * | * | * | * |  |  |  |  | * | * | * | * | * |  |  |
| D |  |  |  |  |  |  |  |  |  |  | * | * | * | * | * |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S: Source device to be encoded D: Device for storing the encoded result n : Length of encoded bits
- Range of $n$ when $S$ is a bit device: $1 \sim 8$
- Range of $n$ when $S$ is a word device: $1 \sim 4$
- The lower " $2^{n}$ " bits of $S$ are encoded and the result is stored in D .
- If several bits of $S$ are 1 , the first bit that is 1 will be processed orderly from high bit to low bit.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example 1:

- When S is used as a bit device, $\mathrm{n}=1 \sim 8$. Errors will occur if $\mathrm{n}=0$ or $\mathrm{n}>8$.
- When $\mathrm{n}=8$, the maximum points to encode is $28=256$ points.
- When X10 $=$ Off $\rightarrow$ On, this instruction will encode the 23 bits data (MO~M7) and store the result in the lower 3 bits (b2~b0) of D0. b15 ~b3 that have not been used in D0 will all become 0 .
- After the execution of this instruction is completed and X10 turns to Off, the content in D remains unchanged.



1. Explanations:

- $S$ : Source device $D$ : Destination device for storing counted value
- If $S$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- Among the bits of $S$, the total of bits whose content is "1" will be stored in D.
- When 32 - instruction is in use, D will occupy 2 registers.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example:

- When $\times 10=$ On, among the 16 bits of D0, the total of bits whose content is "1" will be stored in D2。



1, Explanations:

- S:Source device D: Device for storing check result n : Bits specified for check
- If $S$ is used in device $F$, only 16 -bit instruction is applicable
- Range of n: 0~15 (16-bit instruction); 0~31 (32-bit instruction)
- When the $n$th bit of $S$ is " 1 ", $D=O n$; when the $n$th bit of $S$ is " 0 ", $D=$ Off
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2. Program Example:

- When $X O=O$ O , assume the $15^{\text {th }}$ bit of $D O$ is " 1 ", and $M O=O$. Assume the 15 th bit of $D O$ is " 0 ", and $M O$ $=$ Off.
- When XO goes Off, MO will remains in its previous status.


| $\begin{aligned} & \hline \mathrm{ZL} \\ & 45 \end{aligned}$ |  | MEAN |  |  |  |  |  |  | S | D | n |  |  |  |  | Mean |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | MEAN: 7 steps 16 -bit <br> DMEAN: 13 steps 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E |  |  |
| S |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |
| n |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |  |

1. Explanations:

- $s$ : Start device to obtain mean value $D$ : Destination device for storing mean value n : The number of consecutive source devices used
- If $D$ is used in device $F$, only 16 -bit instruction is applicable.
- Range of n: $1 \sim 64$
- After the content of $n$ devices starting from $S$ are added up, the mean value of the result will be stored in D.
- Remainders in the operation will be left out.
- Provided the No. of designated device exceeds its normal range, only the No. within the normal range can be processed.
- If $n$ falls without the range of $1 \sim 64$, PLC will determine it as an "instruction operation error".
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example:

- When $\times 10=O n$, the contents in $3(\mathrm{n}=3)$ registers starting from DO will be summed and then divided by 3 . The obtained mean value will be stored in D10 and the remainder will be left out.



1. Explanations:

- S:Source device D: Device for storing the result
- This instruction performs a square root operation on $S$ and stores the result in D.
- $S$ can only be a positive value. If $S$ is negative, PLC will regard it as an "instruction operation error" and will not execute this instruction.
- The operation result $D$ should be integer only, and the decimal will be left out.


## 2, Program Example:

- When $\mathrm{X10}=\mathrm{On}$, the instruction performs a square root on D0 and stores the result in D12.



1. Explanations:

- S : Source device for conversion
D: Device for storing the conversion result.
- BIN integer is converted into binary floating point value. At this time, S of the 16 -bit instruction, FLT, occupies 1 register and D occupies 2 registers.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example 1:

- the BIN integer is converted into binary floating point value.
- When X10 = On, D0 (BIN integer) is converted into D13 and D12 (binary floating point value).
- When $\mathrm{X} 11=\mathrm{On}, \mathrm{D} 1$ and DO (BIN integer) are converted into D21 and D20 (binary floating point value).
- If $D 0=K 10, X 10$ will be On. The 32-bit value of the converted floating point will be H 41200000 and stored in 32- bit register D12 (D13).
- If 32-bit register $\mathrm{DO}(\mathrm{DI})=\mathrm{K} 100,000, \mathrm{X} 11$ will be On. The 32 -bit value of the converted floating point will be H47C35000 and stored in 32-bit register D20 (D21).


3. Program Example 2:

- Please use this instruction to complete the following operation.

note: 1) D10 (BIN integer) is converted to D101 and D102 (binary floating point value).

2) $X 7$ ~ X0 (BCD value) are converted to D200 (BIN value).
3) D200 (BIN integer) is converted to D203 and D202 (binary floating point value). 4The result of $\mathrm{K} 615 \div \mathrm{K} 10$ is stored in D301 and D300 (binary floating point value).
4) The result of binary decimal division (D101, D100) $\div($ D203, D202) is stored in D401 and D400 (binary floating point value).
5) The result of binary decimal multiplication (D401, D400) $\times(\mathrm{D} 301, D 300)$ is stored in D21 and D20 (binary floating point value).
6) D21 and D20 (binary floating point value) are converted to D31 and D30 (decimal floating point value).
7) D21 and D20 (binary floating point value) are converted to D41 and D40 (BIN integer).

## 7 Application Instructions ZL50~ZL99

## 7.1 (ZL 50-59) High-speed processing

| $\begin{aligned} & \mathrm{ZL} \\ & 50 \\ & \hline \end{aligned}$ |  | REF |  |  |  |  | D |  |  | n |  |  |  |  |  | Refresh |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | REF: 5 steps 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| D | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- D: Start device to be I/O refreshed $n$ : Number of items to be I/O refreshed
- D must designate $\mathrm{X} 0, \mathrm{X} 10, \mathrm{Y}, \mathrm{Y} 10 \cdots$ the points whose 1 s digit is " 0 ". See remarks for more details
- Range of $\mathrm{n}: 8 \sim 256$ (has to be the multiple of 8 ).
- The status of all PLC input/output terminals will be updated after the program scans to END. When the program starts to scan, the status of the external input terminal is read and stored into the memory of the input point. The output terminal will send the content in the output memory to the output device after END instruction is executed. Therefore, this instruction is applicable when the latest input/output data are needed for the operation.
- D has to be designated to be $\mathrm{XO}, \mathrm{X} 10, \mathrm{YO}, \mathrm{Y} 10 \cdots$ such forms whose 1 st digit is " 0 ". Range of $\mathrm{n}: 8 \sim 256$ (must be 8 ' s multiple); otherwise it will be regarded as an error. The range varies in different models. See Remarks for more details.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example 1:

- When XO = On, PLC will read the status of input points X0 ~ X17 immediately and refresh the input signals without any input delay.


3. Program Example 2:

- When $\mathrm{XO}=\mathrm{On}$, the 8 output signal from $\mathrm{YO} \sim Y 7$ will be sent to output terminals and refreshed without having to wait for the END instruction for output.

| X0 | REF | Y0 | K8 |
| :--- | :--- | :--- | :--- |

7 Application Instructions ZL50~ZL99


1. Explanation:

- n : Response time (unit: ms)
- Range of $\mathrm{n}: \mathrm{n}=\mathrm{K} 0 \sim \mathrm{~K} 60$
- To avoid interferences, X0 ~ X17 are equipped with digital filters on output terminals. Digital filters adjust the response time by REFF instruction.
- Rules for adjusting the reponse time of the filter at X0 $\sim \mathrm{XI} 7$ :
a) When the power of PLC turns from Off to On or the END instruction is being executed, the response time will be determined upon the contents in D1020 and D1021.
b) You can use MOV instruction in the program to move the time values to D1020 and D1021 and make adjustments in the next scan.
c) You can use REFF instruction to change the response time during the execution of the program. The changed response time will be move to D1020 and D1021 and you can make adjustments in the next scan.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example:

- When the power of PLC turns from Off to On, the response time of X0 $\sim X 17$ will be determined by the contents in D1020 and D1021.
- When X20 $=$ On, REFF K5 will be executed and the response time will be changed to 5 ms for the adjustment in the next scan.
- When X20 = Off, the REFF K20 will be executed and the response time will be changed to 20 ms for the adjustment in the next scan.


| ZL 52 |  | MTR |  |  |  |  | $S \quad D_{1} \quad D_{2}$ |  |  |  |  | n |  |  |  | Input Matrix |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | + De | vice |  |  |  |  |  | ord Devid | vices |  |  |  |  |  | MTR: 9 steps | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{D}_{1}$ |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{D}_{2}$ |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- $s$ : Start device of matrix input

D2: Corresponding start device for matrix scan

D1: Start device of matrix output
n : Number of arrays in matrix scan

- $S$ must designate $\times 0, \times 10 \cdots$ the X points whose 1 st digit is " 0 " and occupies 8 consecutive points.
- DI must designate YO, Y10 $\ldots$ the $Y$ points whose 1 st digit is " 0 " and occupies $n$ consecutive points.
- D2 must designate YO, MO. SO $\cdots$ the $Y$, M, S points whose 1st digit is " 0 " .
- Range of $\mathrm{n}: 2 \sim 8$.
- $S$ is the start device No. of all input terminals connected to the matrix. Once $S$ is designated, the 8 points following the No. will be the input terminals in the matrix.
- Dl designate the start device No. of transistor output Y in the matrix scan.
- This instruction occupies continuous 8 input devices starting from S. n external output terminals starting from D1 read the 8 switches of $n$ arrays by matrix scan, obtaining $8 \times n$ multiple-matrix input points. The status of scanned switches will be stored in the devices starting from D2.
- Maximum 8 input switches can be parallelly connected in 8 arrays and obtaining 64 input points $(8 \times 8=$ 64).
- When the 8 -point 8 -array matrix inputs are in use, the reading time of each array is approximately 25 ms , totaling the reading of 8 arrays 200 ms , i.e. the input signals with On/Off speed of over 200 ms are not applicable in a matrix input.
- Whenever this instruction finishes a matrix scan, M1029 will be On for one scan period.
- There is no limitation on the number of times using the instruction, but only one instruction can be executed in one scan cycle.

2. Program Example:

- When PLC RUN, MRT instruction will start to be executed. The statuses of the external 2 arrays of 16 switches, will be read in order and stored in the internal relays M10 ~M17, M20 ~M27.

- The figure below illustrates the external wiring of the 2-array matrix input loop constructed by X40~X47 andY40 ~Y41. The 16 switches correponds to the internal relays M10~M17, M20~M27. Should be used with MTR instruction.

- See the figure above. The 8 points starting from X40 start to perform a matrix scan from Y40~Y41 ( $\mathrm{n}=$ 2). $D_{2}$ designates that the start device No. of the read results is $M 10$, indicating that the first array is read to M10 ~ M17 and the second array is read to M20 ~ M27.

Read input signals in the 1st array


Read input signals in the 2nd array


Processing time of each array: approx. 25 ms

| ZL <br> 53 |  | HSCS |  |  |  |  |  |  |  |  | D |  |  |  |  | High Speed Counter Set |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DHSCS: 13 steps 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  |  | KnM | KnS | T | C | D |  | E | F |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S1: Comparative value S2: No. of high speed counter D: Comparison result
- D can designate IO $\square 0$;$=1 \sim 6$
- The high speed counter inputs counting pulses from the corresponding external input terminals X0 ~ X17 by inserting an interruption. When the high speed counter designated in S2 pluses 1 or minuses 1, DHSCS instruction will perform a comparison immediately. When the present value in the high speed counter equals the comparative value designated in $S 1$, device designated in D will turn On. Even the afterward comparison results are unequal, the device will still be On.
- If the devices specified as the device $D$ are $Y 0 \sim Y 17$, when the compare value and the present value of the high-speed counter are equal, the comparison result will immediately output to the external inputs YO ~ Y17, and other $Y$ devices will be affected by the scan cycle. However, $M, S$ devices are immediate output and will not be affected by the scan cycle.


## 2. Program Example 1:

- After PLC RUN and MO = On, DHSCS instruction will be executed. When the present value in C235 changes from 99 to 100 or 101 to 100 , Y 10 will be On constantly


3. Program Example 2:

- Differences between Y output of DHSCS instruction and general Y output:

1) When the present value in C249 changes from 99 to 100 or 101 to $100, Y 10$ outputs immediately to the external output point by interruption and has nothing to do with the PLC scan time. However, the time will still be delayed by the relay ( 10 ms ) or transistor (10us) of the output module.
2) When the present value in C249 changes from 99 to 100 , the drive contact of $C 249$ will be On immediately. When the execution arrives at SET Y17, Y17 will still be affected by the scan time and will output after END instruction.


4, Program Example 3:

- High speed counter interruption:

1) Operand $D$ of DHSCS instruction can designate $10 \square 0, \square=1 \sim 6$, as the timing of interruption when the counting reaches its target.
2) When the present value in C251 changes from 99 to 100 or 101 to 100 , the program will jump to 1010 and execute the interruption service subroutine.


## Remarks:

1) The output contact of the high speed counter and the comparative outputs of ZL 53 DHSCS, ZL 34 DHSCR and ZL 55 DHSZ instructions only perform comparison and contact outputs when there is a counting input. When using data operation instructions, e.g. DADD, DMOV, for changing the present value in the high speed
counter or making the present value equals the set value, there will not be comparisons or comparative outputs because there is no counting inputs.
2) Supports high speed counters. C235 ~ C240 are program-interruption 1-phase high speed counter with a total bandwidth of 20 kHz , can be used alone with a counting frequency of up to 10 kHz . C241 ~ C254 are hardware high speed counter (HHSC). There are four HHSC of HHSCO ~ 3. The pulse input frequency of HHSC0~4 can reach 200kHz (1 phase or A-B phase).
3) Every HHSC can only be designated to one counter by DCNT instruction.
4) There are three counting modes in every HHSC (see the table below):

- 1-phase 1 input refers to "pulse/direction" mode
- 2-phase 2 inputs refers to "A-B phase" mode.

|  | software high-speed counter |  |  |  |  |  | Hardware high-speed counter |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 phase 1 input |  |  |  |  |  | 1 phase 1 input |  |  |  | 2 phase 2 input |  |  |  |
|  | C235 | C236 | C237 | C238 | C239 | C240 | C241 | C242 | C243 | C244 | C251 | C252 | C253 | C254 |
| X0 | U/D |  |  |  |  |  | U/D |  |  |  | A |  |  |  |
| X1 |  | U/D |  |  |  |  |  |  |  |  | B |  |  |  |
| X2 |  |  | U/D |  |  |  |  | U/D |  |  |  | A |  |  |
| X3 |  |  |  | U/D |  |  |  |  |  |  |  | B |  |  |
| X4 |  |  |  |  | U/D |  |  |  | U/D |  |  |  | A |  |
| X5 |  |  |  |  |  | U/D |  |  |  |  |  |  | B |  |
| X6 |  |  |  |  |  |  |  |  |  | U/D |  |  |  | A |
| $\times 7$ |  |  |  |  |  |  |  |  |  |  |  |  |  | B |

Description:

| U | Progressively increasing input |
| :--- | :--- |
| D | Progressively decreasing input |
| A | A phase input |
| B | B phase input |

5) Counting modes:

Special D1225 ~ D1228 are for setting up different counting modes of the hardware high speed counters (HHSCO $\sim 3$ ). There are normal $\sim 4$ times frequency for the counting and the default setting is double frequency.

7 Application Instructions ZL50～ZL99

| Counting modes |  | Wave pattern |  |
| :---: | :---: | :---: | :---: |
| Type | Set value in special D | Counting up（＋1） | Counting down（－ |
| 1－phase 1 input | 1 <br> （Normal frequency） | U／D $\qquad$ <br> U／D FLAG $\qquad$ |  |
|  | 2 （Double frequency） | U／D $\qquad$ <br> U／D FLAG $\qquad$ |  |
| 1－phase 2 inputs | 1 <br> （Normal frequency） | $\mathrm{C}$ |  |
|  | 2 （Double frequency） | U |  |
| 2－phase <br> 2 inputs | 1 <br> （Normal frequency） |  |  |
|  | 2 （Double frequency） |  |  |
|  | 3 （Triple frequency） | A |  |
|  | 4 （4 times frequency） |  |  |

6）Special registers for relevant flags and settings of high speed counters：

| Flag | Function |
| :---: | :---: |
| M1235～M1244 | C235～C244 High speed counter counting direction specified． <br> When M12םロ＝Off ，C2口ロ：Count on。 <br> When M12ロロ＝On ，C2口ロ：Count off。 |
| D1225 | The counting mode of the 1st group counters（C251） |
| D1226 | The counting mode of the 2nd group counters（C252） |
| D1227 | The counting mode of the 3rd group counters（C253） |
| D1228 | The counting mode of the 4th group counters（C254） |
| D1225～D1228 | PLC hardware high speed counter HHSCO～HHSC3 counting mode setting，not the following setting values are preset for the double frequency counting mode． <br> 2：for the double frequency counting mode，（factory value）． <br> 3 ：it is the triple frequency counting mode． <br> 4：it is the quadruple frequency counting mode．（desired value） |



1. Explanations:

- S1:Comparative value S2: No. of high speed counter D: Comparison result
- $\quad$ S2 has to designate the No. of high speed counters C235 ~ C255. See remarks of ZL 53 DHSCS for more details.
- D of high speed counters C241~C254 that are the same as the counters designated by S2
- The high speed counter inputs counting pulses from the corresponding external input terminals X0~X7 by inserting an interruption. When the No. of high-speed counter designated in S2 "+1" or "-1", DHSCR will perform a comparison immediately. When the present value in the high speed counter equals the comparative value designated in S 1 , the device designated in D will turn Off and even the afterward comparison results are unequal, the device will still be Off.
- If the devices designated in $D$ are $Y 0 \sim Y 17$, when the comparative value equals the present value in the high speed counter, the comparison result will immediately output to the external output terminals YO~ Y17 (and clear the designated $Y$ output) and the rest of $Y$ devices will be affected by the scan cycle. Devices $M$ and $S$ act immediately without being affected by the scan cycle.


## 2, Program Example 1:

- When $M 0=O$ and the present value in the high speed counter C251 changes from 99 to 100 or 101 to 100, Y10 will be cleared and Off.
- When the present value in the high speed counter C251 changes from 199 to 200 , the contact of C251 will be On and make YO = On. However, the program scan time will delay the output.
- Y 10 will immediately reset the status when the counting reaches its target. D is also able to designate high speed counters of the same No. See Program Example 2.


3. Program Example 2:

- When DHSCR instruction designates the same high speed counter, and the present value in the high speed counter C251 changes from 999 to 1,000 or 1,001 to $1,000, \mathrm{C} 251$ will be reset to Off.


| $\begin{aligned} & \hline \mathrm{ZL} \\ & 55 \end{aligned}$ |  |  | HSZ |  |  |  | S 1 |  | S |  |  | D |  |  |  | High Speed Zone Compare |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DHSZ: 17 steps | 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  | KnM | Kns | T | C | D | E |  |  | F |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |  |
| S |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S1: Lower bound of the comparison zone $\quad$ 2: Upper bound of the comparison zone
S: No. of high speed counter
D: Comparison result
- $\quad$ S1 has to be euqal to or smaller than S 2 . $(\mathrm{S} 1 \leqslant \mathrm{~S} 2)$
- When $S 1>$ S2, the instruction will perform a comparison by using $S 1$ as the upper bound and $S 2$ as the lower bound.
- $\quad$ S has to designate high speed counters C235 ~ C255, See remarks of ZL 53 DHSCS for more details.
- D will occupy 3 consecutive devices.
- The output will not be affected by the scan time.
- The zone comparisons and outputs are all processed by inserting interruptions.


## 2, Program Example 1:

- Designate device Y0 and YO ~ Y2 will be automatically occupied.
- When DHSZ instruction is being executed and the counting of the high speed counter C246 reaches upper and lower bounds, one of $\mathrm{YO} \sim \mathrm{Y} 2$ will be On.


3, Program Example 2:

- Use DHSZ instruction for high/low speed stop control. C251 is an A-B phase high speed counter and DHSZ only performs comparison output when there is a C251 counting pulse input. Therefore, even when the present value in the counter is $0, \mathrm{Y} 10$ will not be On.
- When $\times 10=$ On, DHSZ will require that Y10 has to be On when the present value in the counter $\leqslant K 2,000$. To solve this requirement, you can execute DZCPP instruction when the program was first RUN and compare C251 with K2,000. When the present value in the counter $\leqslant K 2,000, Y 10$ will be On. DZCPP instruction is a pulse execution instruction and will only be executed once with Y10 being kept On.
- When the drive contact $\mathrm{X} 10=$ Off, $\mathrm{Y} 10 \sim \mathrm{Y} 12$ will be reset to Off.

- The timing diagram


4, Program Example 3:

- The multiple set values comparison mode: If D of DHSZ instruction designates a special auxiliary relay M1150, the instruction will be able to compare (output) the present value in the high speed counter with many set values.
- In this mode: S1: start device in the comparison table. S1 can only designate data register D and can be modified by E and F. Once this mode is enabled, S 1 will not be changed even the E and F has been changed.

S2: number of group data to be compared. S2 can only designate K1 ~ K255 or H1 ~ HFF and can be modified by E and F. Once this mode is enabled, S2 cannot be changed. If $S 2$ is not within its range, error code 01EA (hex) will display and the instruction will not be executed.
$S$ : No. of high speed counter (designated as C241~C254).
D: Designated mode (can only be M1150)

- The No. of start register designated in S1 and the number of rows (groups) designated in S2 construct a comparison table. Please enter the set values in every register in the table before executing the instruction.
- When the present value in the counter C251 designated in $S$ equals the set values in D1 and D0, the $Y$ output designated by D2 will be reset to Off $(D 3=K 0)$ or On $(D 3=K 1)$ and be kept. Output $Y$ will be processed as an interruption. No. of Y output pointss are in decimal (range: $0 \sim 255$ ). If the No. falls without the range,SET/RESET will not be enabled when the comparison reaches its target.
- When this mode is enabled, PLC will first acquire the set values in D0 and D1 as the target value for the first comparison section. At the same time, the index value displayed in D1 150 will be 0 , indicating that PLC
performs the comparison based on the group 0 data.
- When the group 0 data in the table have been compared, PLC will first execute the Y output set in group 0 data and determine if the comparison reaches the target number of groups. If the comparison reaches the target, M1151 will be On; if the comparison has not reached the final group, the content in D 1150 will plus 1 and continue the comprison for the next group
- M1151 is the flag for the completion of one execution of the table, can be Off by the user. Or when the next comparion cycle takes place and the group 0 data has been compared, PLC will automatically reset the flag.
- When the drive contact of the instruction X10 goes Off, the execution of the instruction will be interrupted and the content in D1150 (table counting register) will be reset to 0 . However, the On/Off status of all outputs will be remained.
- When the instruction is being executed, all set values in the comparison table will be regarded as valid values only when the scan arrives at END instruction for the first time.
- This mode can only be used once in the program.
- This mode can only be used on the hardware high speed counters C241~C254.
- When in this mode, the frequency of the input counting pulses cannot exceed 50 kHz or the neighboring two groups of comparative values cannot differ by 1 ; otherwise there will not be enough time for the PLC to react and result in errors.

- The comparison table:

| 32-bit data for comparison |  |  |  | No. of Y output |  | On/Off indication |  | Table counting register D1150 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| High word |  | Low word |  |  |  |  |  |  |
|  | (K0) | D0 | (K100) |  | (K10) |  | (K1) | 0 |
|  | (K0) |  | (K200) |  | (K11) | D7 | (K1) | 1 |
|  | (K0) |  | (K300) | D10 | (K10) | D1 | (K0) | 2 |
| D13 | (K0) |  | (K400) | D14 | (K11) |  | (K0) | 3 |
|  |  |  |  |  |  |  |  | $0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 0$ <br> Cyclic scan |



- Special registers for flags and relevant settings:

| Flag | Function |
| :--- | :--- |
| M1150 | DHSZ instruction in multiple set values comparison mode |
| M1151 | The execution of DHSZ multiple set values comparison <br> mode is completed |


| Special D | Function |
| :--- | :--- |
| D1150 | Table counting register for DHSZ multiple set values <br> comparison mode |

5, Program Example 4:

- DHSZ and DPLSY instructions are combined for frequency control. If $D$ of DHSZ instruction is a special auxiliary relay M1152, the present value in the counter will be able to control the pulse output frequency of DPLSY instruction.
- In this mode: S1: start device in the comparison table. S1 can only designate data register D and can be modified by E and F. Once this mode is enabled, $S 1$ will not be changed even the $E$ and $F$ has been changed

S2: number of group data to be compared. S2 can only designate K1 ~ K255 or H1 ~ HFF and can be modified by E and F. Once this mode is enabled, S 2 cannot be changed. If S 2 is not within its range, error code 01EA (hex) will display and the instruction will not be executed.

S: No. of high speed counter (designated as C241~C254).
D: Designated mode (can only be M1152)

- This mode can only be used once. this mode can only be used in the hardware high speed counter C241
~ C254. Please enter the set values in every register in the table before executing the instruction.
- When this mode is enabled, PLC will first acquire the set values in DO and DI as the target value for the first comparison section. At the same time, the index value displayed in D1 152 will be 0 , indicating that PLC performs the comparison based on the group 0 data.
- When the group 0 data in the table have been compared, PLC will first execute at the frequency set in group 0 data (D2, D3) and copy the data to D1152 and D1153, determining if the comparison reaches the target number of groups. If the comparison reaches the target, M1153 will be On; if the comparison has not reached the final group, the content in D 1151 will plus 1 and continue the comprison for the next group.
- M1153 is the flag for the completion of one execution of the table, can be Off by the user. Or when the next comparion cycle takes place and the group 0 data has been compared, PLC will automatically reset the flag.
- If you wish to use this mode with PLSY instruction, please preset the value in D1152.
- If you wish to stop the execution at the last row, please set the value in the last row KO .
- When the drive contact of the instruction X 10 goes Off, the execution of the instruction will be interrupted and the content in D1151 (table counting register) will be reset to 0 .
- When in this mode, the frequency of the input counting pulses cannot exceed 50 kHz or the neighboring two groups of comparative values cannot differ by 1 ; otherwise there will not be enough time for the PLC to react and result in errors.

- The comparison table:

| 32-bit data for comparison |  | Pulse output frequency$0 \sim 200 \mathrm{kHz}$ |  | Table counting register D1151 |
| :---: | :---: | :---: | :---: | :---: |
| High word | Low word |  |  |  |
| D1 (K0) | DO (K0) | D3, D2 | ( $\mathrm{K} 5,000$ ) | 0 |
| D5 (K0) | D4 (K100) | D7, D6 | (K10,000) | 1 |
| D9 (K0) | D8 (K200) | D11, D10 | (K15,000) | 2 |
| D13 (K0) | D12 (K300) | D15, D14 | $(\mathrm{K}, 000)$ | 3 |
| D17 (K0) | D16 (K400) | D19, D18 | (KO) | 4 |
|  |  |  |  | $0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 4$ <br> Cyclic scan |



- Special registers for flags and relevant settings:

| Flag | Function |
| :--- | :--- |
| M1152 | DHSZ instruction in frequency control mode |
| M1153 | The execution of DHSZ frequency control mode is completed |


| Special D | Function |
| :--- | :--- |
| D1151 | Table counting register for DHSZ multiple set values comparison mode |
| D1152 (low word) | In frequency control mode, DHSZ reads the upper and lower limits in the table <br> counting register D1153 and D1152. |
| D1153 (high word) | Current number of pulses output by DPLSY instruction |
| D1649 (high word) |  |



- During the execution of DHSZ instruction, do not modify the set values in the comparison table.
- The designated data will be arranged into the the above program diagram when the program executes to END instruction. Therefore, PLSY instruction has to be executed after DHSZ instruction has been executed once.


1, Explanations:

- S1: External pulse input terminal S2: Pulse receiving time (ms) D: Detected result
- The received number of pulses of the input terminal designated in S 1 is calculated within the time (in ms ) designated in S 2 . The result is stored in the register designated in D .
- D will occupy 5 consecutive devices. $\mathrm{D}+1$ and D are the detected value obtained from the previous pulses; $D+3$ and $D+2$ are the current accumulated number of values; $D+4$ is the counting time remaining (max. $32,767 \mathrm{~ms}$ ).
- This instruction is mainly used for obtaining a proportional value of rotation speed. The result D and rotation speed will be in proportion. The following equation is for obtaining the rotation speed of motor.

$$
\mathrm{N}=\frac{60(\mathrm{DO})}{\mathrm{nt}} \times 10^{3}(\mathrm{rpm})
$$

$\mathrm{N}: \quad$ Rotation speed
n : The number of pulses produced per rotation
t : Detecting time designated in $\mathbf{S}_{2}(\mathrm{~ms})$

- The $X$ input point designated by this instruction cannot be used again as the pulse input terminal of the high speed counter or as an external interruption signal.
- There is no limitation on the times of using this instruction in the program, but only one instruction will be executed at a time.

2, Program Example:

- When $\mathrm{X7}$ = On, D2 will calculate the high-speed pulses input by X 1 and stop the calculation automatically after $1,000 \mathrm{~ms}$. The result will be stored in D0.
- When the $1,000 \mathrm{~ms}$ counting is completed, D2 will be cleared to 0 . When $X 7$ is On again, D2 will start the calculation again.



1, Explanations:

- S1: Pulse output frequency module)
S2: Number of output pulses

D: Pulse output device (please use transistor outpu $\dagger$

- S2 designates the number of output pulses. The 16-bit instruction can designate $1 \sim 32,767$ pulses and the 32- bit instruction can designate $1 \sim 2,147,483,647$ pulses.
- When the PLSY instruction is used in the program, the output cannot be duplicated with the output of the ZL 58 PWM instruction and the ZL 59 PLSR instruction.
- Number of continuous pulses for all series:

| Group No |  |  | current number <br> of output pulses <br> (32-bit | Pulse complete flag |
| :--- | :--- | :--- | :--- | :--- |
| integer) |  |  |  |  |


| CH21 (Y52,Y53) | Y52 | Y53 | D1520 | M1119 |
| :--- | :--- | :--- | :--- | :--- |
| CH22 (Y54,Y55) | Y54 | Y55 | D1536 | M1205 |
| CH23 (Y56,Y57) | Y56 | Y57 | D1552 | M1206 |

- When PLSY instruction is executed, it will designate the number of output pulses (S2) output from the output device (D) at a pulse output frequency (S1).
- When PLSY instruction is used in the program, its outputs cannot be the same as those in ZL 58 PWM and ZL 59 PLSR.
- when PLSY and DPLSY instruction is disabled, the user will have to reset the pulse output completed flags.
- The user has to reset the pulse output completed flags after the pulse output is completed.
- After PLSY instruction starts to be executed, Y will start a pulse output. Modifying S2 at this moment will not affect the current output. If you wish to modify the number of output pulses, you have to first stop the execution of PLSY instruction and modify the number.
- S1 can be modified when the program executes to PLSY instruction.
- Off time : On time of the pulse output = 1:1.
- When the program executes to PLSY instruction, the current number of output pulses will be stored in the special data registers. See remarks for more details.


## 2, Program Example:

- When $\mathrm{XO}=\mathrm{On}$, there will be 200 pulses output from Y0 at 1 kHz . When the pulse output is completed, M1029 will be On and $Y 10$ will be On.
- When $\mathrm{XO}=\mathrm{Off}$, the pulse output from YO will stop immediately. When XO is On again, the output will start again ffrom the first pulse.
- If the frequency needs to be sent all the time, write 0 to $S 2$.


| Group No | PUL | DIR | current number of output pulses (32-bit integer) | Pulse complete flag | Pulse sending | Emergency stop without slowing down |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| CH0 (Y0,Y1) | YO | Y1 | D1648 | M1029 | M1344 | M1308 |
| CH1 (Y2,Y3) | Y2 | Y3 | D1664 | M1030 | M1345 | M1309 |
| CH2 (Y4,Y5) | Y4 | Y5 | D1680 | M1036 | M1346 | M1310 |
| CH3 (Y6,Y7) | Y6 | Y7 | D1696 | M1037 | M1347 | M1311 |
| CH4 (Y10,Y11) | Y10 | Y11 | D1712 | M1 102 | M1348 | M1312 |
| CH5 (Y12,Y13) | Y12 | Y13 | D1728 | M1103 | M1349 | M1313 |
| CH6 (Y14,Y15) | Y14 | Y15 | D1744 | M1104 | M1350 | M1314 |
| CH7 (Y16,Y17) | Y16 | Y17 | D1760 | M1105 | M1351 | M1315 |
| CH8 (Y20,Y21) | Y20 | Y21 | D1776 | M1106 | M1352 |  |
| CH9 (Y22,Y23) | Y22 | Y23 | D1792 | M1107 | M1353 |  |
| CH10 (Y24,Y25) | Y24 | Y25 | D1808 | M1108 | M1354 |  |
| CH11 (Y26,Y27) | Y26 | Y27 | D1824 | M1109 | M1355 |  |
| CH12 (Y30,Y31) | Y30 | Y31 | D1840 | M1110 | M1356 |  |
| CH13 (Y32,Y33) | Y32 | Y33 | D1856 | M1111 | M1357 |  |
| CH14 (Y34,Y35) | Y34 | Y35 | D1872 | M1112 | M1358 |  |
| CH15 (Y36,Y37) | Y36 | Y37 | D1888 | M1113 | M1359 |  |
| CH16 (Y40,Y41) | Y40 | Y41 | D1904 | M1114 | M1360 |  |
| CH17 (Y42,Y43) | Y42 | Y43 | D1920 | M1115 | M1361 |  |
| CH18 (Y44,Y45) | Y44 | Y45 | D1472 | M1116 | M1362 |  |
| CH19 (Y46, Y47) | Y46 | Y47 | D1488 | M1117 | M1363 |  |
| CH20 (Y50,Y51) | Y50 | Y51 | D1504 | M1118 | M1364 |  |
| CH21 (Y52,Y53) | Y52 | Y53 | D1520 | M1119 | M1365 |  |
| CH22 (Y54,Y55) | Y54 | Y55 | D1536 | M1205 | M1366 |  |
| CH23 (Y56,Y57) | Y56 | Y57 | D1552 | M1206 | M1367 |  |
| Remarks |  |  | D1648:Low word of the current number of output pulses from CHO. D1649:High word of the current number of output pulses from CHO . | After CHOCH23 pulse output is completed, the corresponding flag bit is ON | Only when the pulse is being sent, the flag bit corresponding to $\mathrm{CHO}-\mathrm{CH} 23$ is ON | Off->On: The high-speed pulse output pauses immediately. On->Off: Continuing to output the pulses which have not been output |


| $\begin{aligned} & \mathrm{ZL} \\ & 58 \end{aligned}$ |  | PWM |  |  |  |  |  |  |  |  | D |  |  |  |  | Pulse Width Modulation |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | PWM: 7 step | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- s : Pulse output width S : Pulse output period D: Pulse output device (please use transistor output module)
- $\mathrm{S} 1 \leqslant \mathrm{~S} 2$
- Range of $S 1$ : ( $\dagger$ ) $0 \sim 32,767 \mathrm{~ms}$. (refer to the remarks for more information about the time unit settings.)
- Range of S2: (T) $1 \sim 32,767 \mathrm{~ms}(b u t S 1 \leqslant \$ 2)$.
- Pulse output device

| Output point | Y0, Y2, Y4, Y6 $\cdots \cdots$ Y24, Y26 |
| :--- | :--- |

- When PWM instruction is used in the program, its outputs cannot be the same as those of API 57 PLSY, API 59 PLSR or other positioning instructions.
- PWM instruction designates the pulse output width in S1 and pulse output period in S2 and outputs from output device D.
- When, $\mathrm{S} 1<0$ or $S 2 \leqslant 0$ or $S 1>S 2$, there will be operational errors, and there will be no output from the pulse output device. When $\$ 1=0$, there will be no output from the pulse output device. When $\$ 1=\$ 2$, the the pulse output device will keep being On.
- $\quad S 1$ and $S 2$ can be changed when PWM instruction is being executed.

2. Program Example:

- When $\mathrm{XO}=\mathrm{On}, \mathrm{YO}$ will output the pulses as below. When X0 $=$ Off, YO output will also be Off


| $\begin{array}{\|l\|} \hline \mathrm{ZL} \\ 59 \end{array}$ |  | PLSR |  |  |  |  |  |  | S2 |  | S3 |  | D |  |  | Pulse Ramp |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  | + De | vic |  | Word Devices |  |  |  |  |  |  |  |  |  |  | PLSR: 9 steps DPLSR: 17 step |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  | 16-bit |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| S3 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S1:Maximum speed of pulse output

S3: Acceleration/deceleration time (ms) D: Pulse output device (please use transistor output module PLC)

- Range of S1: $10 \sim 32,767 \mathrm{~Hz}$ ( $16-\mathrm{bit}$ ); $10 \sim 200,000 \mathrm{~Hz}$ ( $32-\mathrm{bit}$ ). The maximum speed has to be 10 ' s multiple; if not, the 1s digit will be left out. $1 / 10$ of the maximum speed is the variation of one acceleration or deleration. Please be aware if the variation reponds to the acceleration/deceleration demand from the step motor, in case the step motor may crash.
- Range of S2: $110 \sim 32,767$ (16-bit); $110 \sim 2,147,483,647$ (32-bit). If $\$ 2$ is less than 110 , the pulet output will be abnormal.
- Range of 53 : below $5,000 \mathrm{~ms}$. The acceleration time and deceleration time have to be the same.
- Refer to the related section in explanation of PLSY instruction for D devices and maximum frequency.
- PLSR instruction is a pulse output instruction with acclerating and decelerating functions. The pulses accelerate from the static status to target speed and decelerates when the target distance is nearly reached. The pulse output will stop when the target distance is reached. S 2 and $\$ 3$ can be changed when PLSR instruction is being executed.

7 Application Instructions ZL50~ZL99

| Group No | PUL | DIR | current number of output pulses (32-bit integer) | Pulse complete flag | Pulse sending | Emergency stop without slowing down |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| CH0 (Y0,Y1) | YO | Y1 | D1648 | M1029 | M1344 | M1308 |
| CH1 (Y2,Y3) | Y2 | Y3 | D1664 | M1030 | M1345 | M1309 |
| CH2 (Y4,Y5) | Y4 | Y5 | D1680 | M1036 | M1346 | M1310 |
| CH3 (Y6,Y7) | Y6 | Y7 | D1696 | M1037 | M1347 | M1311 |
| CH4 (Y10,Y11) | Y10 | Y11 | D1712 | M1102 | M1348 | M1312 |
| CH5 (Y12,Y13) | Y12 | Y13 | D1728 | M1103 | M1349 | M1313 |
| CH6 (Y14,Y15) | Y14 | Y15 | D1744 | M1104 | M1350 | M1314 |
| CH7 (Y16,Y17) | Y16 | Y17 | D1760 | M1105 | M1351 | M1315 |
| CH8 (Y20,Y21) | Y20 | Y21 | D1776 | M1106 | M1352 | M1316 |
| CH9 (Y22,Y23) | Y22 | Y23 | D1792 | M1107 | M1353 | M1317 |
| CH10 (Y24,Y25) | Y24 | Y25 | D1808 | M1108 | M1354 | M1318 |
| CH11 (Y26,Y27) | Y26 | Y27 | D1824 | M1109 | M1355 | M1319 |
| CH12 (Y30,Y31) | Y30 | Y31 | D1840 | M1110 | M1356 | M1320 |
| CH13 (Y32,Y33) | Y32 | Y33 | D1856 | M1111 | M1357 | M1321 |
| CH14 (Y34,Y35) | Y34 | Y35 | D1872 | M1112 | M1358 | M1322 |
| CH15 (Y36,Y37) | Y36 | Y37 | D1888 | M1113 | M1359 | M1323 |
| CH16 (Y40,Y41) | Y40 | Y41 | D1904 | M1114 | M1360 | M1324 |
| CH17 (Y42, Y43) | Y42 | Y43 | D1920 | M1115 | M1361 | M1325 |
| CH18 (Y44,Y45) | Y44 | Y45 | D1472 | M1116 | M1362 | M1326 |
| CH19 (Y46,Y47) | Y46 | Y47 | D1488 | M1117 | M1363 | M1327 |
| CH20 (Y50,Y51) | Y50 | Y51 | D1504 | M1118 | M1364 | M1328 |
| CH21 (Y52,Y53) | Y52 | Y53 | D1520 | M1119 | M1365 | M1329 |
| CH22 (Y54,Y55) | Y54 | Y55 | D1536 | M1205 | M1366 | M1330 |
| CH23 (Y56,Y57) | Y56 | Y57 | D1552 | M1206 | M1367 | M1331 |
| Remarks |  |  | D1648:Low word of the current number of output pulses from CHO . D1649:High word of the current number of output pulses from CHO . | After CHOCH23 pulse output is completed, the corresponding flag bit is ON | Only when the pulse is being sent, the flag bit corresponding to $\mathrm{CH} 0-\mathrm{CH} 23$ is ON | Off->On: The high-speed pulse output pauses immediately. On->Off: <br> Continuing to output the pulses which have not been output |

- when all the $\mathrm{CHO}(\mathrm{YO}, \mathrm{Y} 1)$ pulses have been sent, M 1029 will be On; when all the $\mathrm{CH} 1(\mathrm{Y} 2, \mathrm{Y} 3)$ pulses have
been sent, M1030 will be On; when CH2 (Y4, Y5) pulses have been sent, M1036 will be On; when CH3 (Y6, Y7) pulses have been sent, M1037 will be On. When all the CH4 (Y10, Y11) pulses have been sent, M1102 will be On. When all the CH5 (Y12, Y13) pulses have been sent, M1 103 will be On. Next time when PLSR instruction is enabled, M1029, M1030, M1036, M1037, M1 102 and M1103 will be 0 again and after the pulse output is completed, they will become 1 again. Other pulse output can be deduced by analogy, the flag bit is detailed in the table above.
- During every acceleration section, the number of pulses ( frequency $\times$ time) may not all be integers. PLC will round up the number to an integer before the output. Therefore, the acceleration time of every section may not be exactly the same. The offset is determined upon the frequency and the decimal after rounding up. In order to ensure the correct number of output pulses, PLC will supplement insufficient pulses in the last section.
- For the limitation on the times using this instruction in the program, refer to PLSY instruction for more information.
- Range of $\mathbf{S}_{3}$ : below $5,000 \mathrm{~ms}$. The acceleration time and deceleration time have to be the same.
a. The acceleration and deceleration time must be more than 10 times the maximum scan period (the contents of D1012), and if the value is set to less than 10x, the slope of the acceleration and deceleration will be incorrect.
b. The minimum setting value for the acceleration and deceleration time can be determined by the following formula.

$$
S_{3} \geq \frac{90,000}{S_{1}}
$$

If the setting value is less than the result of the above calculation formula, the acceleration and deceleration time will become larger, and if the setting value is less than $90000 / \mathrm{S} 1$, the result value of $90000 / \mathrm{S} 1$ will be used as the setting value.
c. The maximum setting value for the acceleration and deceleration time can be determined by the following formula.

$$
\mathrm{S}_{3} \leq \frac{\mathrm{S}_{2}}{\mathrm{~S}_{1}} \times 818
$$

d. The number of variable speed segments for acceleration and deceleration is fixed at 10 segments. If
the input acceleration and deceleration time is greater than the maximum setting value, the maximum setting time will prevail, and if it is less than the minimum setting value, the minimum setting value will be the main value.

D pulse output device, additional deceleration pulse output device

| Refer to the output of the modulation pulse table |  |
| :--- | :--- |
| PLSR output | Y0, Y2, Y4, Y6 $\cdots \cdots$ Y54, Y56 |

2, Example:

- When $M 0=O n$, the PLSR instruction is executed with the maximum frequency value of pulse output $1,000 \mathrm{~Hz}$, the total pulse number of all pulse output D10, and the acceleration/deceleration time 3,000ms, then Y0 outputs pulses. Start outputting pulses at a frequency of $1,000 / 10 \mathrm{~Hz}$ each time. The time of each frequency output pulse is fixed 3,000/9 (ms).
- When MO turns Off, the output is interrupted. When XO turns On again, the pulse count starts from 0 .



## Remarks:

- Based on the number of pulses. If the output cannot reach the maximum acceleration frequency within the acceleration/deceleration time offered, the instruction will automatically adjust the acceleration/deceleration time and the maximum frequency.The operands have to be set before the execution of the instruction PLSR. You cannot change the acceleration/deceleration during the instruciton execution.
- All acceleration/deceleration instructions are included with the brake function. The brake function will be enabled when PLC is performing acceleration and the switch contact is suddenly Off. The deceleration will operate at the slope of the acceleration。

Frequency F


## 7.2 (ZL 60-69) Convenience instructions



1. Explanations:

- SI: Start device for data stack comparison

S2: Data to be compared
D: Start device for storing comparison result
n : Length of data to be compared

- When $S 2$ are used in device $F$, only 16 -bit instruction is applicable.
- D will occupy 5 consecutive points.
- Range of n : for 16 -bit instruction $1 \sim 256$; for 32-bit instruction $1 \sim 128$.
- The n data in the registers starting from S1 are compared with S2 and the results are stored in the registers starting from D .
- In the 32-bit instruction, S1, S2, D and n will designate 32-bit registers.
- For D, the 16-bit counters and 32-bit counters cannot be mixed when being used

2. Program Example:

- When $\mathrm{XO}=\mathrm{On}$, the data stack consist of D10 ~ D19 will be compared against D0 and the result will be stored in D50 ~ D52. If there are equivalent values appearing during the comparison, D50 ~ D52 will all be 0.
- The data are compared algebraically. ( $-10<2$ ).
- The No. of the register with the smallest value among the compared data will be recorded in D53; the biggest will be recorded in $D 54$. When there are more than one smallest value or biggest value, device $D$ will record the No. of the register with bigger value.


| $\begin{aligned} & \mathrm{ZL} \\ & 62 \end{aligned}$ |  | ABSD |  |  |  |  |  |  | S2 D |  |  | $n$ |  |  |  | Absolute Drum Sequencer |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | ABSD: 9 steps <br> DABSD: 17 steps | $\begin{aligned} & \text { 16-bit } \\ & \text { 32-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY |  |  |  | KnM | Kns | T | C |  |  | D | E | F |
| S1 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S1: Start device in the data table S2: No. of counter

D: Start No. of the devices for the comparison results $n$ : Number of data for comparison

- When S 1 designates $\mathrm{KnX}, \mathrm{KnY}, \mathrm{KnM}$ and KnS , the 16 -bit instruction has to designate K 4 and 32-bit instruction has to designate K 8 .
- Range of n : $1 \sim 64$
- ABSD instruction is for the absolute control of the multiple output pulses generated by the present value in the counter.
- $\quad$ S2 of DABSD instruction can designate high speed counters. However, when the present value in the high speed counter is compared with the target value, the result cannot output immediately owing to the scan time. If an immediate output is required, please use DHSZ instruction that is exclusively for high speed counters.


## 2, Program Example:

- Before the execution of ABSD instruction, use MOV instruction to write all the set values into D100 ~ D107 in advance. The even-number $D$ is for lower bound value and the odd-number $D$ is for upper bound value.
- When $\mathrm{X10}=\mathrm{On}$, the present value in counter Cl 10 will be compared with the four groups of lower and upper bound values in D100 ~ D107. The comprison results will be stored in M10~M13.
- When $\mathrm{X10}=$ Off, the original On/Off status of M10~M13 will be remained.

- M10~M13 will be On when the present value in C10 $\leqq$ upper bound value or $\geqq$ lower bound value.

| Lower bound value | Upper bound value | Present value in C10 | Output |
| :---: | :---: | :---: | :---: |
| $\mathrm{D} 100=40$ | $\mathrm{D} 101=100$ | $40 \leqq \mathrm{C} 10 \leqq 100$ | $\mathrm{M} 10=\mathrm{On}$ |
| $\mathrm{D} 102=120$ | $\mathrm{D} 103=210$ | $120 \leqq \mathrm{C} 10 \leqq 210$ | $\mathrm{M} 11=\mathrm{On}$ |
| Lower bound value | Upper bound value | Present value in C10 | Output |
| $\mathrm{D} 104=140$ | $\mathrm{D} 105=170$ | $140 \leqq \mathrm{C} 10 \leqq 170$ | $\mathrm{M} 12=$ On |
| $\mathrm{D} 106=150$ | $\mathrm{D} 107=390$ | $150 \leqq \mathrm{C} 10 \leqq 390$ | $\mathrm{M} 13=\mathrm{On}$ |

- If the lower bound value > upper bound value, when C10 < upper bound value (60) or > upper bound value (140), M12 will be On.

| Lower bound value | Upper bound value | Present value in C10 | Output |
| :---: | :---: | :---: | :---: |
| $D 100=40$ | $D 101=100$ | $40 \leqq C 10 \leqq 100$ | $M 10=$ On |
| $D 102=120$ | $D 103=210$ | $120 \leqq C 10 \leqq 210$ | $M 11=$ On |
| $D 104=140$ | $D 105=60$ | $60 \leqq C 10 \leqq 140$ | $M 12=$ On |
| $D 106=150$ | $D 107=390$ | $150 \leqq C 10 \leqq 390$ | $M 13=$ On |



| $\begin{aligned} & \mathrm{ZL} \\ & 63 \end{aligned}$ |  | INCD |  |  |  |  | S 1 |  | S2 D |  |  | n |  |  |  | Incremental Drum Sequencer |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | it De | Devic |  | Word Devices |  |  |  |  |  |  |  |  |  |  | INCD: 9 steps | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S1:Start device in the data table S2: No. of counter

D: Start No. of the devices for the comparison results $n$ : Number of data for comparison

- When SI designates $\mathrm{KnX}, \mathrm{KnY}, \mathrm{KnM}$ and KnS , it has to designate K4.
- In the 16-bit instruction, S 2 has to designate CO ~ C198 and will occupy 2 consecutive No. of counters.
- Range of $\mathrm{n}: 1 \sim 64$
- INCD instruction is for the relative control of the multiple output pulses generated by the present value in the counter.
- The present value in $S 2$ is compared with $S 1$. $\$ 2$ will be reset to 0 whenever a comparison is completed. The current number of data processed in temporarily stored in $\mathrm{S} 2+1$.


## 2, Program Example:

- Before the execution of INCD instruction, use MOV instruction to write all the set values into D100~D104 in advance. D100 = 15, D101 = 30, D102 = 10, D103 = 40, D104 $=25$.
- The present value in C10 is compared against the set values in D100~D104. The present value will be reset to 0 whenever a comparison is completed.
- The current number of data having been processed is temporarily stored in C11.
- The number of times of reset is temporarily stored in C 11 .
- Whenever the content in C11 pluses 1, M10 ~ M14 will also correspondingly change. See the timing diagram below.
- When XO goes from On to Off, ClO and $\mathrm{Cl1}$ will both be reset to 0 and M10~M14 will all be Off. When XO is On again, the instruction will start its execution again from the beginning.

| X0 | M1013 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
|  |  | CNT | C10 | K100 |  |
|  | INCD | D100 | C10 | M10 | K5 |



| ZL 64 |  | TTMR |  |  |  |  | D |  |  | n |  |  |  |  |  | Teaching Timer |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | TTMR: 5 steps | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- D: Device No. for storing the "On" time of button switch n: Multiple setting
- D will occupy 2 consecutive devices
- Range of $\mathrm{n}: 0 \sim 2$
- TTMR instruction can be used 8 times in the program.
- The "On" time (unit: 100 ms ) of the external button switch is stored in device No. D + 1. The "On" time (unit: second) of the switch is multiplied by n and stored in D .
- Multiple setting:

When $\mathrm{n}=0$, unit of $\mathrm{D}=$ second
When $\mathrm{n}=1$, unit of $\mathrm{D}=100 \mathrm{~ms}(\mathrm{D} \times 10)$
When $\mathrm{n}=2$, unit of $\mathrm{D}=10 \mathrm{~ms}(\mathrm{D} \times 100$

2, Program Example 1:

- The "On" (being pressed) time of button switch XO is stored in D1. The setting of n is stored in DO. Therefore, the button switch will be able to adjust the set value in the timer.
- When X0 goes Off, the content in D1 will be cleared to 0 , but the content in DO will remain.

- Assume the "On" time of XO is T (sec.), see the relation between DO, D1 and n in the table below.

| n | DO | Dl (unit: 100 ms ) |
| :--- | :--- | :--- |
| K0 (unit: s) | $1 \times \mathrm{T}$ | $\mathrm{Dl}=\mathrm{D} 0 \times 10$ |
| K1 (unit: 100 ms ) | $10 \times \mathrm{T}$ | $\mathrm{Dl}=\mathrm{D} 0$ |
| K2 (unit: 10 ms ) | $100 \times \mathrm{T}$ | $\mathrm{Dl}=\mathrm{D} 0 / 10$ |

## 3. Program Example 2:

- Use TMR instruction to write in 10 groups of set time.
- Write the set values into D100 ~ D109 in advance.
- The timing unit for timer T0 $\sim$ T9 is 0.1 sec . The timing unit for the teaching timer is 1 sec .
- Connect the 1-bit DIP switch to XO ~ X3 and use BIN instruction to convert the set value of the switch into a bin value and store it in E .
- Store the "On" time (sec.) of X10 in D200.
- M0 refers to the pulses generated from one scan period after the button switch of the teaching timer X10 is released.
- Use the set number of the DIP switch as the indirectly designated pointer and send the content in D200 to DIOOE (D100~D109).


Remarks:

- There is no limitation on the times using this instruction in the program and 8 instructions can be executed at the same time.

| ZL <br> 65 |  | STMR |  |  |  |  | S m D |  |  |  |  |  |  |  |  | Special Timer |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | STMR: 7 step | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| s |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |  |  |  |
| m |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S: No. of timer m: Set value in timer (unit: 100 ms ) D: No. of start output device
- Range of S: for TO ~T183.
- Range of m: $1 \sim 32,767$
- D will occupy 4 consecutive devices.
- STMR instruction is used for Off-delay, one shot timer and flashing sequence.
- The No. of timers designated by STMR instructions can be used only once.

2, Program Example:

- When $\mathrm{X10}=\mathrm{On}$, STMR instruction will designate timer TO and set the set value in TO as 5 seconds.
- Y0 is the contact of Off-delay. When X10 goes from Off to On, YO will be On. When X10 goes from On to Off, YO will be Off after a five seconds of delay.
- When X10 goes from On to Off, there will be a five seconds of $\mathrm{Y} 1=$ On output.
- When X10 goes from Off to On, there will be a five seconds of Y2 = On output.
- When X10 goes from Off to On, Y3 will be On after a five seconds of delay. When X10 goes from On to Off, Y3 will be Off after a five seconds of delay.

- Add $a b$ contact of $Y 3$ after X 10 , and Y 1 and Y 2 can operate for flashing sequence output. When X 10 goes Off, YO , Y 1 and Y 3 will be Off and the content in T 10 will be reset to 0 .


| ZL 66 |  | ALT |  |  |  |  | D |  |  |  |  |  |  |  |  | Alternate State |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | ALT: 3 steps | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- D: Destination device
- When ALT instruction is executed, "On" and "Off" of $D$ will switch.
- This instruction adopts pulse execution instructions (ATLP).
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example 1

- When X0 goes from Off to On, Y0 will be On. When X0 goes from Off to On for the second time, Y0 will be Off.


3. Program Example 2:

- Using a single switch to enable and disable control. At the beginning, $M 0=O f f$, so $\mathrm{YO}=\mathrm{On}$ and $\mathrm{Y} 1=\mathrm{Off}$. When X 10 switches between On/Off for the first time, MO will be On, so Y1 = On and YO = Off. For the second time of On/Off switching, $M 0$ will be Off, so $\mathrm{YO}=\mathrm{On}$ and $\mathrm{Y} 1=$ Off.


4. Program Example 3:

- Generate flashing. When $\mathrm{X10}=\mathrm{On}, \mathrm{TO}$ will generate a pulse every 2 seconds and YO output will switch between On and Off following the TO pulses.


| $\begin{aligned} & \mathrm{ZL} \\ & 67 \end{aligned}$ |  | RAMP |  |  |  |  | S 1 |  | S2 |  | D |  | n |  |  | Ramp Variable Value |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | it De | evice |  |  |  |  |  | ord D | vices |  |  |  |  |  | RAMP: 9 steps | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S1:Start of ramp signal S2: End of ramp signal

D: Duration of ramp signal $n$ : Scan times

- Range of n: $1 \sim 32,767$
- D will occupy 2 consecutive points.
- This instruction is for obtaining slope (the relation between linearity and scan time). Before using this instruction, you have to preset the scan time.
- The set value of start ramp signal is pre-written in D10 and set value of end ramp signal in D11. When X10 $=$ On, D10 increases towards D1 1 through $\mathrm{n}(=100$ ) scans (the duration is stored in D12). The times of scans are stored in D13.
- In the program, first drive M1039 = On to fix the scan time. Use MOV instruction to write the fixed scan time to the special data register D1039. Assume the scan time is 30 ms and take the above program for example, $n=K 100$, the time for D10 to increase to D1 1 will be 3 seconds (30ms $\times 100$ ).
- When X10 goes Off, the instruction will stop its execution. When $\times 10$ goes On again, the content in D12 will be reset to 0 for recalculation.
- When M1026 = Off, M1029 will be On and the content in D12 will be reset to the set value in D10.
- When this instruction is used with analog signal outputs, it will be able to buffer START and STOP.

2, Program Example:

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| X10 | RAMP | D10 | D11 | D12 | K100 |
| :--- | :--- | :--- | :--- | :--- | :--- |



- Remarks:

D12 for enabling On/Off of M1026:


| $\begin{aligned} & \mathrm{ZL} \\ & 69 \end{aligned}$ |  | SORT |  |  |  |  |  | $S_{1}$ | $\mathrm{m}_{1}$ | $\mathrm{m}_{2}$ |  | D |  | n |  |  | Sort Tabulated Data |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  | De | evice |  | Word Devices |  |  |  |  |  |  |  |  |  |  |  | SORT: 11 steps DSORT: 21 step | $\begin{aligned} & \text { 16-bit } \\ & 32 \text {-bit } \end{aligned}$ |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T |  |  | C |  | D |  |  |  | E | F |
| $\mathrm{S}_{1}$ |  |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |  |
| ml |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| m2 |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  | * |  |  |  |  |  |

1. Explanations:

- S : Start device for the original data m 2 : Number of columns of data
n : Reference value for data sortin
- Range of ml: $1 \sim 32$
- Range of m2: $1 \sim 6$
- Range of $\mathrm{n}: 1 \sim \mathrm{~m} 2$
- The sorted result is stored in $\mathrm{ml} \times \mathrm{m} 2$ registers starting from the device designated in D . Therefore, if S and $D$ designate the same register, the sorted result will be the same as the data designated in $S$.
- It is better that the start No. designated in S is 0 .
- There is no limitation on the times of using this instruction. However, only one instruction can be executed at a time.
- The function of sorting one-dimensional data is added. If users set ml and m 2 to 1 , the function will be enabled. The operand $n$ represents the number of data. It must be in the range of 1 to 32 . The data in the $n$ devices starting from $S$ is sorted. The sorting result is stored in the devices starting from $D$. This function only needs one scan time. After data is sorted. M1 029 will be ON.

2. Program Example:

- When $\mathrm{XO}=$ On, the sorting will start. When the sorting is completed, M1029 will be On. DO NOT change the datato be sorted during the execution of the instruction. If you wish to change the data, please make X0 go from Off to On again.

| X0 |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| SORT | D0 | K5 | K5 | D50 | D100 |

- Example table of data sorting

|  |  | Columns of data: $\mathrm{m}_{2}$ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Data Column |  |  |  |  |
|  | Column <br> Row | 1 | 2 | 3 | 4 | 5 |
|  |  | Students No. | Physics | English | Math | Chemistry |
|  | 1 | (DO) 1 | (D5) 90 | (D10) 75 | (D15) 66 | (D20) 79 |
|  | 2 | (D1) 2 | (D6) 55 | (D11) 65 | (D16) 54 | (D21) 63 |
|  | 3 | (D2) 3 | (D7) 80 | (D12) 98 | (D17) 89 | (D22) 90 |
|  | 4 | (D3) 4 | (D8) 70 | (D13) 60 | (D18) 99 | (D23) 50 |
|  | 5 | (D4) 5 | (D9) 95 | (D14) 79 | (D19) 75 | (D24) 69 |

- Sorted data when D100 = K3

|  |  | Columns of data: $\mathrm{m}_{2}$ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Data Column |  |  |  |  |
|  | Column | 1 | 2 | 3 | 4 | 5 |
|  | Row | Students No. | Physics | English | Math | Chemistry |
| 4 | 1 | (D50) 4 | (D55) 70 | (D60) 60 | (D65) 99 | (D70) 50 |
| \% | 2 | (D51) 2 | (D56) 55 | (D61) 65 | (D66) 54 | (D71) 63 |
| ¢ | 3 | (D52) 1 | (D57) 90 | (D62) 75 | (D67) 66 | (D72) 79 |
| 言 | 4 | (D53) 5 | (D58) 95 | (D63) 79 | (D68) 75 | (D73) 69 |
| 1 | 5 | (D54) 3 | (D59) 80 | (D64) 98 | (D69) 89 | (D74) 90 |

- $\quad$ Sorted data when D100 $=$ K5

|  |  | Columns of data: $\mathrm{m}_{2}$ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Data Column |  |  |  |  |
|  | Column | 1 | 2 | 3 | 4 | 5 |
|  | Row | Students No. | Physics | English | Math | Chemistry |
| - | 1 | (D50) 4 | (D55) 70 | (D60) 60 | (D65) 99 | (D70) 50 |
| \% | 2 | (D51) 2 | (D56) 55 | (D61) 65 | (D66) 54 | (D71) 63 |
| $\stackrel{\text { ¢ }}{0}$ | 3 | (D52) 5 | (D57) 95 | (D62) 79 | (D67) 75 | (D72) 69 |
| O | 4 | (D53) 1 | (D58) 90 | (D63) 75 | (D68) 66 | (D73) 79 |
| $\dagger$ | 5 | (D54) 3 | (D59) 80 | (D64) 98 | (D69) 89 | (D74) 90 |

## 7.3 (ZL 70-79) External I/O device

| ZL 70 |  | TKY |  |  |  |  |  |  |  | D2 |  | 2 |  | n |  | Ten Key Inpu† |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | TKY: 7 steps DTKY: 13 step | $\begin{aligned} & \text { 16-bit } \\ & 32 \text {-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | $T$ |  |  | C | D |  |  | E | F |
| S | * | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{D}_{1}$ |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |  |
| $\mathrm{D}_{2}$ |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S: Start device for key input D1: Device for storing keyed-in value D2: Key output signal
- $S$ will occupy 10 consecutive points; D2 will occupy 11 consecutive points.
- This instruction designates 10 external input points (representing decimal numbers $0 \sim 9$ ) starting from $S$. The 10 points are respectively connected to 10 keys. By pressing the keys, you can enter a 4-digit decimal figure 0 ~ 9,999 (16-bit instruction) or a 8-digit figure 0 ~ 99,999,999 (32-bit instruction) and store the figure in D1. D2 is used for storing key status.
- There is no limitation on the times of using this instruction. However, only one instruction can be executed at a time.


## 2. Program Example:

- Connect the 10 input points starting from X0 to the 10 keys ( $0 \sim 9$ ). When $X 20=O n$, the instruction will be executed and the keyed-in values will be stored in D0 in bin form. The key status will be stored in M10 ~ M19.


- As shown in the timing chart below, the 4 points $X 5, X 3, X 0$, and $X 1$ connected to the keys are entered in order and you can obtain the result 5,301. Store the result in D0. 9,999 is the maximum value allowed to stored in D0. Once the value exceeds 4 digits, the highest digit will overflow.
- M12 = On when from X2 is pressed to the other key is pressed. Same to other keys.
- When any of the keys in X0~X11 is pressed, one of M10~19 will be On correspondingly.
- M20 = On when any of the keys is pressed.
- When X20 goes Off, the keyed-in value prior to D0 will remain unchanged, but M10~M20 will all be Off.


7 Application Instructions ZL50~ZL99

| ZL <br> 71 |  | HKY |  |  |  |  |  | S | $\mathrm{D}_{1}$ |  |  | $D_{3}$ |  |  |  | Hexadecimal Key Input |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | HKY: 9 steps <br> DHKY: 17 step |  | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E |  |  |  |  |
| S | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{D}_{1}$ |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 32-bit |
| $\mathrm{D}_{2}$ |  |  |  |  |  |  |  |  |  |  | * | * | * | * |  |  |  |  |
| $\mathrm{D}_{3}$ |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S: Start device for key scan input

D2: Device for storing keyed-in value

D1: Start device for key scan output
D3: Key output signal

- S will occupy 4 consecutive points.
- D1 will occupy 4 consecutive points.
- D3 will occupy 8 consecutive points.
- This instruction designates 4 continuous external input points starting from $S$ and 4 continuous external input points starting from D1 to construct a 16-key keyboard by a matrix scan. The keyed-in value will be stored in D2 and D3 is used for storing key status. If several keys are pressed at the same time, the first key pressed has the priority.
- The keyed-in value is termporarily stored in DO. When the 16-bit instruction HKY is in use, 9,999 is the maximum value DO is able to store. When the value exceeds 4 digits, the highest digit will overflow. When the 32-bit instruction DHKY is in use, $99,999,999$ is the maximum value D0 is able to store. When the value exceeds 8 digits, the highest digit will overflow.
- There is no limitation on the times of using this instruction. However, only one instruction can be executed at a time.

2, Program Example:

- Designate 4 input points $\mathrm{X} 10 \sim \mathrm{X} 13$ and the other 4 input points $\mathrm{Y} 10 \sim \mathrm{Y} 13$ to construct a 16 -key keyboard. When $\mathrm{X} 4=\mathrm{On}$, the instruction will be executed and the keyed-in value will be stored in D0 in bin form. The key status will be stored in MO ~M7

| X4 | HKY | X10 | Y10 | D0 | M0 |
| :--- | :--- | :--- | :--- | :--- | :--- |

- Key in numbers:

- Function keys input:

1) When $A$ is pressed, MO will be On and retained. When $D$ is pressed next, M0 will be Off, M3 will be On and retained.
2) When many keys are pressed at the same time, the first key pressed has the priority.


- Key output signal:

1) When any of $A \sim F$ is pressed, M6 will be On for once.
2) When any of $0 \sim 9$ is pressed, M7 will be On for once.

- When X4 goes Off, the keyed-in value prior to D0 will remain unchanged, but M0~M7 will all be Off.
- External wiring:


Remarks:

- When this instruction is being executed, it will require 8 scans to obtain one valid keyed-in value. A scan period that is too long or too short may result in poor keyed-in effect, which can be avoided by the following methods:
a) If the scan period is too short, I/O may not be able to respond in time, resulting in not being able to read the keyed-in value correctly. In this case, please fix the scan time.
b) If the scan period is too long, the key may respond slowly. In this case, write this instruction into the time interruption subroutine to fix the time for the execution of this instruction.
- Functions of M1167:
a) When M1167 = On, HKY instruction will be able to input the hexadecimal value of $0 \sim F$.
b) When M1167 = Off, HKY instruction will see A ~ F as function keys.
- Functions of D1037 :

Write D1037 to set the overlapping time for keys (unit: ms). The overlapping time will vary upon different program scan time and the settings in D1037.

| $\begin{aligned} & \mathrm{ZL} \\ & 72 \\ & \hline \end{aligned}$ |  | DSW |  |  |  |  | S |  | D1 | D |  | 2 n |  |  |  | Digital Switch |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DSW: 9 step | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| S | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{D}_{1}$ |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{D}_{2}$ |  |  |  |  |  |  |  |  |  |  | * | * | * | * | * |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S:Start device for switch scan input

D2: Device for storing the set value of switch

DI: Start device for switch scan output
n : Groups of switches

- Range of $\mathrm{n}: 1 \sim 2$.
- This instruction designates 4 or 8 consecutive external input points starting from $S$ and 4 consecutive external input points starting from D1 to scan read 1 or 24 -digit DIP switches. The set values of DIP switches are stored in D2. n decides to read 1 or 24 -digit DIP switches.
- There is no limitation on the times of using this instruction in the program. However, two instructions are allowed to be executed at a time.

2. Program Example:

- The first group of DIP switches consist of X20~X23 and Y20~Y23. The second group of switches consist of X24 ~X27 and Y20 ~Y23. When X10 = On, the instruction will be executed and the set values of the first group switches will be read and converted into bin values before being stored in D20. The set values of the second group switches will be read, converted into bin values and stored in D21.

| X10 | DSW | X20 | Y20 | D20 | K2 |
| :---: | :--- | :--- | :--- | :--- | :--- |

- When X10 = On, the Y20 ~ Y23 auto scan cycle will be On. Whenever a scan cycle is completed, M1029 will be On for a scan period.
- Please use transistor output for Y20 ~ Y23. Every pin 1, 2, 4, 8 shall be connected to a diode ( $0.1 \mathrm{~A} / 50 \mathrm{~V}$ ) before connecting to the input terminals on PLC.

- Wiring for DIP swich input:


Remarks:

- When $n=K 1$, D2 will occupy one register. When $n=K 2, D 2$ will occupy 2 consecutive registers.
- Follow the methods below for the transistor scan output:
a) When $\mathrm{X10}=\mathrm{On}$, DSW instruction will be executed. When X 10 goes Off, M10 will keep being On until the scan output completes a scan cycle and go Off.
b) When X 10 is used as a button switch, whenever X 10 is pressed once, M 10 will be reset to Off when the scan output designated by DSW instruction completes a scan cycle. The DIP switch data will be read completely and the scan output will only operate during the time when the button switch is pressed. Therefore, even the scan output is a transistor type, the life span of the transistor can be extended because it does not operate too frequently.


7 Application Instructions ZL50~ZL99

| $\begin{aligned} & \text { ZL } \\ & 73 \end{aligned}$ |  | SEGD |  |  |  |  |  |  |  | D |  |  |  |  |  | Seven Segment Decoder |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | SEGD: 5 steps | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |  |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |  |  |

1. Explanations:

- $s$ : Source device to be decoded
D: Output device after the decoding
- When $\mathrm{X10}=\mathrm{On}$, the contents ( $0 \sim \mathrm{~F}$ in hex) of the lower 4 bits ( $\mathrm{b} 0 \sim \mathrm{~b} 3$ ) of D10 will be decoded into a 7 segment display for output. The decoded results will be stored in Y10~Y17. If the content exceeds 4 bits, the lower 4 bits are still used for the decoding.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

| X10 |  |  |  |
| :--- | :--- | :--- | :--- |
|  | SEGD | D10 | K2Y10 |

- Decoding table of the 7 -segment display:

| Hex | Bit combination | Composition of the 7 segment display | Status of each segment |  |  |  |  |  |  | Data displayed |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | B0(a) | B1(b) | B2(c) | B3(d) | B4(e) | B5(f) | B6(g) |  |
| 0 | 0000 |  | ON | ON | ON | ON | ON | ON | OFF | [1 |
| 1 | 0001 |  | OFF | ON | ON | OFF | OFF | OFF | OFF | \| |
| 2 | 0010 |  | ON | ON | OFF | ON | ON | OFF | ON | $\underline{\square}$ |
| 3 | 0011 |  | ON | ON | ON | ON | OFF | OFF | ON | こ |
| 4 | 0100 |  | OFF | ON | ON | OFF | OFF | ON | ON | 4 |
| 5 | 0101 |  | ON | OFF | ON | ON | OFF | ON | ON | E |
| 6 | 0110 | a | ON | OFF | ON | ON | ON | ON | ON | $E$ |
| 7 | 0111 | $\dagger \mathrm{g} b$ | ON | ON | ON | OFF | OFF | ON | OFF | $F$ |
| 8 | 1000 | J | ON | ON | ON | ON | ON | ON | ON | El |
| 9 | 1001 | d | ON | ON | ON | ON | OFF | ON | ON | 딘 |
| A | 1010 |  | ON | ON | ON | OFF | ON | ON | ON | Fl |
| B | 1011 |  | OFF | OFF | ON | ON | ON | ON | ON | EI |
| C | 1100 |  | ON | OFF | OFF | ON | ON | ON | OFF | F |
| D | 1101 |  | OFF | ON | ON | ON | ON | OFF | ON | 딘 |
| E | 1110 |  | ON | OFF | OFF | ON | ON | ON | ON | $E$ |
| F | 1111 |  | ON | OFF | OFF | OFF | ON | ON | ON | $F$ |


| ZL 74 |  | SEGL |  |  |  |  |  |  | S | D | n |  |  |  |  | Seven Segment with Latch |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | SEGL: 7 steps | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S: Source device to be displayed in 7-segment display D: Start device for 7-segment display scan output
n : Polarity setting of output signal and scan signal
- Range of n : $0 \sim 7$. See remarks for more details.
- The instruction can be used twice in the program.
- This instruction occupies 8 or 12 continuous external input points starting from $D$ for displaying 1 or 24 -digit 7-segment display data and outputs of scanned signals. Every digit carries a 7 -segment display drive (to convert the BCD codes into 7 -segment display signal). The drive also carries latch control signals to retain the 7 -segment display.
- n decides there be 1 group or 2 groups of 4 -digit 7 -segment display and designates the polarity for the output.
- When there is 1 group of 4 -digit output, 8 output points will be occupied. When there are 2 groups of 4digit output, 12 output points will be occupied.
- When this instruction is being executed, the scan output terminals will circulate the scan in sequence. When the drive contact of the instruction goes from Off to On again, the scan output terminal will restart the scan again.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example:

- When $\mathrm{X10}=\mathrm{On}$, this instruction starts to be executed, Y10~Y17 construct a 7-segment display scan circuit. The value in D10 will be converted into BCD codes and sent to the first group 7 -segment display. The value in D11 will be converted into BCD codes as well and sent to the second group 7 -segment display. If the values in D10 and D11 exceed 9,999, operational error will occur

| X 10 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| H | SEGL | D10 | Y10 | K4 |

- When X10 = On, Y14~Y17 will circulate the scan automatically. Every cycle requires 12 scan period. Whenever a cycle is completed, M1029 will be On for a scan period.
- When there is 1 group of 4 -digit 7 -segment display, $\mathrm{n}=0 \sim 3$.
a) Connect the already decoded 7 -segment display terminals 1, 2, 4, 8 in parallel an connect them to Y10 ~ Y13 on the PLC. Connect the latch terminals of each digit to Y14~Y17 on the PLC.
b) When $\mathrm{X10}=\mathrm{On}$, the instruction will be executed and the content in D10 will be sent to the 7 -segment displays in sequence by the circulation of Y14~Y17.
- When there is 2 groups of 4 -digit 7 -segment display, $n=4 \sim 7$.
a) Connect the already decoded 7 -segment display terminals $1,2,4,8$ in parallel an connect them to $\mathrm{Y} 20 \sim$ Y23 on the PLC. Connect the latch terminals of each digit to Y14~Y17 on the PLC.
b) The contents in D10 are sent to the first group 7-segment display. The contents in D11 are sent the the second group 7-segment display. If D10 $=$ K1234 and D $11=K 4321$, the first group will display 1234 , and the second group will display 4321 .
- Wiring of the 7 -segment display scan output.


Remarks:

- When this instruction is executed, the scan time has to be longer than 10 ms . If the scan time is shorter than 10 ms , please fix the scan time at 10 ms .
- n is for setting up the polarity of the transistor output and the number of groups of the 4 -digit 7 -segment display.
- The output point must be a transistor module of NPN output type with open collector outputs. The output
has to connect to a pull-up resistor to VCC (less than 30VDC). Therefore, when output point Y is On, the signal output will be in low voltage.

- Positive logic (negative polarity) output of BCD code

| BCD value |  |  |  | Y output (BCDcode) |  |  |  | Signal output |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| b3 | b2 | b1 | bo | 8 | 4 | 2 | 1 | A | B | C | D |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 |

- Negative logic (positive polarity) output of BCD code.

| BCD value |  |  |  | Y output (BCDcode) |  |  |  | Signal output |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{b}_{3}$ | $\mathrm{~b}_{2}$ | $\mathrm{~b}_{1}$ | $\mathrm{~b}_{0}$ | 8 | 4 | 2 | 1 | A | B | C | D |
| 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 |
| 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |

- Scan latched signal display

| Positive logic (negative polarity) |  | Negative logic (positive polarity) |  |
| :---: | :---: | :---: | :---: |
| Y output (latch) | Output signal | Y output (latch) | Output signal |
| 1 | 0 | 0 | 1 |

- Settings of n :

| Groups of 7-segment display | 1 group |  |  |  | 2 groups |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Y output of BCD code | + |  | - |  | + | - |  |  |
| Scan latched signal display | + | - | + | - | + | - | + | - |
| $\mathbf{n}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |

+ : Positive logic (negative polarity) output -: Negative logic (positive polarity) output
- The polarity of transistor output and the polarity of the 7-segment display input can be the same or different by the setting of $n$.

| ZL 75 |  | ARWS |  |  |  |  |  |  | $D_{1} D_{2}$ |  |  | n |  |  |  |  | Arrow Switch |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  |  | ARWS: 9 steps | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D |  | E | F |  |  |
| S | * | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\mathrm{D}_{1}$ |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |  |  |
| $\mathrm{D}_{2}$ |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S:Start device for key input

D2: Start device for 7-segmentdisplay scan output

DI: Device to be displayed in 7 -segment display n : Polarity setting of output signal and scan signal

- $S$ will occupy 4 consecutive points.
- Range of $\mathrm{n}: 0 \sim 3$. See remarks of API 74 SEGL for more details.
- There no limitation on the times of using this instruction in the program. However, only one instruction is allowed to be executed at a time.
- The output points designated by this instruction shall be transistor output.
- When using this instruction, please fix the scan time, or place this instruction in the time interruption subroutine (16~18 ㅁ).
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example:

- When this instruction is executed, X20 is defined as down key, X21 is defined as up key, X22 is defined as right key and X 23 is defined as left key. The keys are used for setting up and displaying external set values. The set values (range: $0 \sim 9,999$ ) are stored in D20.
- When $\mathrm{X10}=$ On, digit 103 will be the valid digit for setup. If you press the left key at this time, the valid digit will circulate as $10^{3} \rightarrow 10^{0} \rightarrow 10^{1} \rightarrow 10^{2} \rightarrow 10^{3} \rightarrow 10^{0}$
- If you press the right key at this time, the valid digit will circulate as $10^{3} \rightarrow 10^{2} \rightarrow 10^{1} \rightarrow 10^{0} \rightarrow 10^{3} \rightarrow 10^{2}$. During the circulation, the digit indicators connected Y24~Y27 will also be On interchangeably following the circulation.
- If you press the up key at this time, the valid digit will change as $0 \rightarrow 1 \rightarrow 2 \cdots \rightarrow 8 \rightarrow 9 \rightarrow 0 \rightarrow 1$. If you press the down key, the valid digit will change as $0 \rightarrow 9 \rightarrow 8 \cdots \rightarrow 1 \rightarrow 0 \rightarrow 9$. The changed value will
also be displayed in the 7-segment display.


| ZL 76 |  | ASC |  |  |  |  |  |  |  | D |  |  |  |  |  | ASCII Code Conversion |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | ASC: 11 steps | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| S |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |  |  |

1. Explanations:

- S: English letter to be converted into ASCII code
D: Device for storing ASCII code
- If the execution of this instruction is connected to a 7 -segment display, the error message can be displayed by English letters.

2. Program Example:

- When $\mathrm{XO}=\mathrm{On}$, convert A $\sim$ H into ASCII code and stored it in DO ~ D3.



## 7.4 (ZL 80-89) External SER equipment

| $\begin{aligned} & \text { ZL } \\ & 80 \end{aligned}$ |  | RS |  |  |  |  |  |  | $m$ D |  |  | n |  |  |  | Serial Communication Instruction |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | RS: 9 steps | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |  |
| S |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |
| m |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |

1. Explanations:

- $S$ : Start device for the data to be transmitted $m$ : Length of data to be transmitted

D: Start device for receiving data $n$ : Length of data to be received

- Range of m: $0 \sim 256$
- Range of $\mathrm{n}: 0 \sim 256$
- The instruction RS supports COM2 (RS-485)
- This instruction is a handy instruction exclusively for MPU to use RS-485 serial communication interface. The user has to pre-store word data in $S$ data register, set up data length $m$ and the data receiving register D and received data length $n$. If $E, F$ index registers are used to modify $S$ and $D$, the user cannot change the set values of E and F when the instruction is being executed; otherwise errors may cause in data writing or reading.
- Designate $m$ as $K 0$ if you do not need to send data. Designate $n$ as $K O$ if you do not need to receive data.
- There is no limitation on the times of using this instruction in the program, but only one instruction is allowed to be executed at a time.
- During the execution of RS instruction, changing the data to be transmitted will be invalid.
- If the peripheral devices, e.g. AC motor drive, are equipped with RS-485 serial communication and its communication format is open, you can use RS instruction to design the program for the data transmission between PLC and the peripheral device.
- If the communication format of the peripheral device is Modbus, PLC offers handy communication instructions API 100 MODRD, API 101 MODWR, and API 150 MODRW, to work with the device. See explanations of the instructions in this application manual.
- For the special auxiliary relays M1120~M1161 and special data registers D1120~D1131 relevant to RS-485 communication, see remarks for more details.

2, Program Example 1:

- Use COM2 on the PLC to carry out RS-485 communication.
- Write the data to be transmitted in advance into registers starting from D100 and set M1 122 (sending request flag) as On.
- When $\mathrm{X10}=\mathrm{On}$, RS instruction will be executed and PLC will start to wait for the sending and receiving of data. D100 starts to continuousl send out 10 data and when the sending is over, M1122 will be automatically reset to Off (DO NOT use the program to execute RST M1122). After 1 ms of waiting, PLC will start to receive the 10 data. Store the data in consecutive registers starting from D120.
- When the receiving of data is completed, M1123 will automatically be On. After the program finishes processing the received data, M1123 has to be reset to Off and the PLC will start to wait for the sending and receiving of data again. DO NOT use the program to continuously execute RST M1123.


3, Program Example 2:

- Use COM2 on the PLC to carry out RS-485 communication

Switching between 8-bit mode (M1161 = On) and 16-bit mode (M1161=Off)
8-bit mode:

1) The head code and tail code of the data are set up by M1126 and M1130 together with D1124 ~ D1126. When PLC is executing RS instruction, the head code and tail code set up by the user will be
sent out automatically. M1161 = On indicates PLC in 8-bit conversion mode. The 16-bit data will be divided into the higher 8 bits and lower 8 bits. The higher 8 bits are ignored and only the lower 8 bits are valid for data transmission.


Sending data: (PLC -> external equipment)

| STX | D100L | D101L | D102L | D103L | ETX1 | ETX2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Head code | (S) source data register, starting from |  |  |  | Tail code 1 | Tail code 2 |

Receiving data: (External equipment -> PLC)
Receiving data: (External equipment -> PLC)

| D120L | D121L | D122L | D123L | D124L | D125L |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Head <br> code | Sreceived data register, starting from <br> the lower 8 bits of D120 | Tail code <br> 1 | Tail code <br> 2 |  |  |
|  | n length $=7$ |  |  |  |  |

When receiving data, PLC will receive the head code and tail code of the data from the external equipment; therefore, the user has to be aware of the setting of data length $n$.
2) 16 -bit mode:

The head code and tail code of the data are set up by M1126 and M1130 together with D1124 ~ D1126. When PLC is executing RS instruction, the head code and tail code set up by the user will be sent out automatically. M1161 = Off indicates PLC in 16-bit conversion mode. The 16-bit data will be divided into the higher 8 bits and lower 8 bits for data transmission.


Sending data: (PLC -> external equipment)

| STX | D100L | D100L | D101L | D101L | ETX1 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Head <br> code | ETX2 <br> source data register, starting from <br> the lower 8 bits of D100 | Tail code <br> 1 | Tail code <br> 2 |  |  |
|  | (m) length $=4$ |  |  |  |  |

Receiving data: (External equipment -> PLC)

| D120L | D120H | D121L | D121H | D122L | D122H | D123L |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Head code |  | ived dat ower 8 | gister, st <br> D120 | g from | Tail code 1 | Tail code 2 |

When receiving data, PLC will receive the head code and tail code of the data from the external equipment; therefore, the user has to be aware of the setting of data length $n$

## 4, Program Example 3:

Use COM2 on the PLC to carry out RS-485 communication.
Connect PLC to AC motor drives (AC motor drive in ASCII Mode; PLC in 16-bit mode and M1161 $=$ Off). Write in the 6 data starting from parameter address H 2101 in AC motor drive in advance as the data to be transmitted


PLC $\rightarrow$ AC motor drive, PLC sends ": 010321010006 D4 CR LF"

AC motor drive $\rightarrow$ PLC, PLC receives ": $01030 C 010017660000000001360000$ 3B CR
LF "Registers for sent data (PLC sends out message)

| Register | Data |  | Explanation |  |
| :---: | :---: | :---: | :---: | :---: |
| D100 low | ' $\because$ | 3 AH | STX |  |
| D100 high | '0' | 30 H | ADR 1 | Address of AC motor drive: ADR$(1,0)$ |
| D101 low | '1' | 31 H | ADR 0 |  |
| D101 high | '0' | 30 H | CMD 1 | Instruction code: CMD (1,0) |
| D102 low | '3' | 33 H | CMD 0 |  |
| D102 high | '2' | 32 H | Start data address |  |
| D103 low | '1' | 31 H |  |  |  |
| D103 high | '0' | 30 H |  |  |  |
| D104 low | '1' | 31 H |  |  |  |
| D104 high | '0' | 30 H | Number of data (counted by words) |  |
| D105 low | '0' | 30 H |  |  |  |
| D105 high | '0' | 30 H |  |  |  |
| D106 low | '6' | 36 H |  |  |  |
| D106 high | 'D' | 44 H | LRC CHK 1 | Error checksum: $\operatorname{LRC} \operatorname{CHK}(0,1)$ |
| D107 low | '4' | 34 H | LRC CHK 0 |  |
| D107 high | CR | D H | END |  |
| D108 low | LF | A H |  |  |  |

Registers for received data (AC motor drive responds with messages)

| Register | Data |  | Explanation |
| :---: | :---: | :---: | :---: |
| D120 low | ':' | 3AH | STX |
| D120 high | '0' | 30 H | ADR 1 |
| D121 low | '1' | 31 H | ADR 0 |
| D121 high | '0' | 30 H | CMD 1 |
| D122 low | '3' | 33 H | CMD 0 |
| D122 high | '0' | 30 H | Number of data (counted by byte) |
| D123 low | 'C' | 43 H |  |
| D123 high | '0' | 30 H | Content of address 2101 H |
| D124 low | '1' | 31 H |  |
| D124 high | '0' | 30 H |  |
| D125 low | '0' | 30 H |  |
| D125 high | '1' | 31 H | Content of address 2102 H |
| D126 low | '7' | 37 H |  |
| D126 high | '6' | 36 H |  |
| D127 low | '6' | 36 H |  |
| D127 high | '0' | 30 H | Content of address 2103 H |
| D128 low | '0' | 30 H |  |
| D128 high | '0' | 30 H |  |
| D129 low | '0' | 30 H |  |
| D129 high | '0' | 30 H | Content of address 2104 H |
| D130 low | '0' | 30 H |  |
| D130 high | '0' | 30 H |  |
| D131 low | '0' | 30 H |  |
| D131 high | '0' | 30 H | Content of address 2105 H |
| D132 low | '1' | 31 H |  |
| D132 high | '3' | 33 H |  |
| D133 low | '6' | 36 H |  |
| D133 high | '0' | 30 H | Content of address 2106 H |
| D134 low | '0' | 30 H |  |
| D134 high | '0' | 30 H |  |
| D135 low | '0' | 30 H |  |
| D135 high | '3' | 33 H | LRC CHK 1 |
| D136 low | 'B' | 42 H | LRC CHK 0 |
| D136 high <br> D137 low | CR | $\frac{\mathrm{DH}}{\mathrm{AH}}$ | END |

5, Program Example 4:
Use COM2 on the PLC to carry out RS-485 communication.
Connect PLC to AC motor drives (AC motor drive in RTU Mode; PLC in 16-bit mode and M1161=On).
Write in H 12 to parameter address H2000 in VFD-B in advance as the data to be transmitted.


PLC $\rightarrow$ AC motor drive, PLC sends: 0106200000120207
AC motor drive $\rightarrow$ PLC, PLC receives: 0106200000120207
Registers for sent data (PLC sends out messages)

| Register | Data |  |
| :--- | :--- | :--- |
| D100 low | 01 H | Address |
| D101 low | 06 H | Function |
| D102 low | 20 H | Data address |
| D103 low | 00 H |  |
| D104 low | 00 H | Data content |
| D105 low | 12 H |  |
| D106 low | 02 H | CRC CHK Low |
| D107 low | 07 H | CRC CHK High |

Registers for received data (VFD-B responds with messages)

| Register | Data | Explanation |
| :--- | :--- | :--- |
| D120 low | 01 H | Address |
| D121 low | 06 H | Function |
| D122 low | 20 H | Data address |
| D123 low | 00 H |  |
| D124 low | 00 H | Data content |
| D125 low | 12 H |  |
| D126 low | 02 H | CRC CHK Low |
| D127 low | 07 H | CRC CHK High |

1. PLC COM2 RS-485: Associated flags (Auxiliary relays) and special registers (Special D) for communication instructions RS / MODRD / MODWR / FWD / REV / STOP / MODRW.
2. How to set up RS-485 communication protocol in D1 120

7 Application Instructions ZL50~ZL99

| - | Content | 0 | 1 |
| :---: | :---: | :---: | :---: |
| b0 | Data length | 7 | 8 |
| $\begin{aligned} & \text { b1 } \\ & \text { b2 } \end{aligned}$ | Parity bits | 00: None <br> 01: Odd <br> 11: Even |  |
| b3 | Stop bits | 1 bit | 2 bits |
| $\begin{aligned} & \text { b4 } \\ & \text { b5 } \\ & \text { b6 } \\ & \text { b7 } \end{aligned}$ | 0001 $(\mathrm{H} 1)$ $:$ <br> 0010 $(\mathrm{H} 2)$ $\vdots$ <br> 0011 $(\mathrm{H} 3)$ $\vdots$ <br> 0100 $(\mathrm{H} 4)$ $\vdots$ <br> 0101 $(\mathrm{H} 5)$ $\vdots$ <br> 0110 $(\mathrm{H} 6)$ $\vdots$ <br> 0111 $(\mathrm{H} 7)$ $\vdots$ <br> 1000 $(\mathrm{H} 8)$ $\vdots$ <br> 1001 $(\mathrm{H} 9)$ $\vdots$ <br> 1010 $(\mathrm{HA})$ $\vdots$ <br> 1011 $(\mathrm{HB})$ $\vdots$ <br> 1100 $(\mathrm{HC})$  | 110 150 300 600 1200 2400 4800 9600 19200 38400 57600 115200 |  |
| b8 | Start word | None | D1124 |
| b9 | First end word | None | D1125 |
| b10 | Second end word | None | D1126 |
| b15 ~ b11 | Not defined |  |  |

3, When RS instruction is in use, the frequently used communication format in the peripheral device will define the start word and end word of the control string. Therefore, you can set up the start word and end word in D1124 ~ D1126 for COM2 or use the start word and end word defined by the PLC. When you use M1126, M1130 and D1124 ~ D1126 to set up the start word and end word, b8 ~ b10 of D1120 have to be set as 1 to make valid the RS-485 communication protocol. See the table below for how to set up.

|  |  | M1130 |  |
| :---: | :---: | :---: | :---: |
|  |  | 0 | 1 |
| $\stackrel{\sim}{\stackrel{N}{c}}$ | 0 | D1124: user defined | D1124: H 0002 |
|  |  | D1125: user defined | D1125: H 0003 |
|  |  | D1126: user defined | D1126: H 0000 (no setting) |
|  | 1 | D1124: user defined | D1124: H 003A ( ${ }^{\prime} \cdot{ }^{\prime}$ ) |
|  |  | D1125: user defined | D1125: H 000D (CR) |
|  |  | D1126: user defined | D1126: H 000A (LF) |

4, Example of how to set up the communication format of COM2:
Assume there is a communication format: Baud rate 9600 7, N, 2
STX : "."
ETX1 : "CR"
ETX2 : "LF"

Check the table and obtain the communication format H788 and write it into D1120.

| b1 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | bo |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| D1120 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |


| M1002 | MOV | H788 | D1120 |
| :---: | :--- | :--- | :--- |

When STX, ETX1 and EXT2 are in use, please be aware of the On and Off of the special auxiliary relays M1126 and M1130.
5, M1143 is for the selection of ASCII mode or RTU mode. On = RTU mode; Off = ASCII mode.
Take the standard Modbus format for example:
In ASCII mode (M1143 = Off)

| STX | Start word = ' ${ }^{\prime}$ (3AH) |
| :---: | :---: |
| Address Hi | Communication address: <br> The 8 -bit address consists of 2 ASCII codes |
| Address Lo |  |
| Function Hi | Function code: <br> The 8 -bit function code consists of 2 ASCII codes |
| Function Lo |  |
| DATA ( n -1) | Data: <br> The $\mathrm{n} \times 8$-bit data consists of 2 n ASCII codes |
| ....... |  |
| DATA 0 |  |
| LRC CHK Hi | LRC checksum: <br> The 8 -bit checksum consists of 2 ASClI code |
| LRC CHK Lo |  |
| END Hi | End word: <br> $\mathrm{END} \mathrm{Hi}=\mathrm{CR}(0 \mathrm{DH}), \mathrm{END} \mathrm{Lo}=\mathrm{LF}(0 \mathrm{AH})$ |
| END Lo |  |

The communication protocol is in Modbus ASCII mode, i.e. every byte is composed of 2 ASCII characters. For example, 64 Hex is ' 64 ' in ASCII, composed by ' 6 ' ( 36 Hex ) and ' 4 ' ( 34 Hex ). Every hex ' 0 '.... $9^{\prime}$ ', ' $\mathrm{A}^{\prime} . .$. ' F ' corresponds to an ASCII code.

| Character | $' 0 '$ | $' 1 '$ | $' 2 '$ | $' 3 '$ | $' 4 '$ | $' 5 '$ | $' 6 '$ | $' 7 '$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ASCII code | 30 H | 31 H | 32 H | 33 H | 34 H | 35 H | 36 H | 37 H |


| Character | $' 8$ ' | ' 9 ' | ' $A^{\prime}$ | 'B' | 'C' | ' $D^{\prime}$ | ' $E$ ' | ' $\mathrm{F}^{\prime}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ASCII code | 38 H | 39 H | 41 H | 42 H | 43 H | 44 H | 45 H | 46 H |

Start word (STX):
Fixed as ':' (3AH)
Address:
' 0 ' ' 0 ': Broadcasting to all drivers
' 0 ' ' 1 ': To the driver at address 01
'0' ' F ': To the driver at address 15
' 1 ' ' 0 ': To the driver at address 16
....and so on, maximum to the driver at address 254 (' $F$ ' ' $E$ ')
Function code:
' 0 ' ' 1 ': Reading several bit devices
'0' '2': Reading several bit devices (read-only devices)
' 0 ' ' 3 ': Reading several word devices
'0' '4': Reading several word devices (read-only devices)
' 0 ' '5': Writing a state in a single bit device
'0' '6': Writing data in a single word device
' 0 ' ' $F$ ': Writing states in bit devices
' 1 ' ' 0 ': Writing data in word devices
'1' '7': Reading word devices and writing data in word devices
Data characters: The data sent by the user.
LRC checksum:
LCR checksum is 2's complement of the value added from Address to Data Content.
For example: $01 \mathrm{H}+03 \mathrm{H}+21 \mathrm{H}+02 \mathrm{H}+00 \mathrm{H}+02 \mathrm{H}=29 \mathrm{H}$. 2 's complement of $29 \mathrm{H}=\mathrm{D} 7 \mathrm{H}$
End word (END):
Fixed as END Hi $=\mathrm{CR}(0 \mathrm{DH})$, END Lo $=\mathrm{LF}(0 \mathrm{AH})$
End word (END):
Fixed as END Hi $=\mathrm{CR}(0 \mathrm{DH})$, END Lo $=\mathrm{LF}(0 \mathrm{AH})$
For example: Read 2 continuous data stored in the registers of the driver at address 01 H (see the table below).
The start register is at address 2102 H .


Responding message:

| STX | '. |
| :---: | :---: |
| Slave station address | '0' |
|  | '1' |
| Function code | '0' |
|  | 3 |
| Number of data (counted by byte) | '0' |
|  | '4' |
| Content in start address2102 H | '1' |
|  | 7 |
|  | 7 |
|  | '0' |
| Content of address$2103 \mathrm{H}$ | '0' |
|  | '0' |
|  | '0' |
|  | '0' |
| LRC check | '7' |
|  | '1' |
| END | CR |
|  | LF |

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In RTU mode (M1143 = On)

| Name | Data (hexadecimal system) |
| :---: | :---: |
| START | See the following explanation |
| Address | Communication address: In 8-bit binary |
| Function | Function code: In 8-bit binary |
| DATA ( $\mathrm{n}-1$ ) | Data: <br> $\mathrm{n} \times 8$-bit data |
| DATA 0 |  |
| CRC CHK Low | CRC checksum: 16 -bit CRC consists of 28 -bit binary |
| CRC CHK High |  |
| END | See the following explanation |

Address:
00 H : Broadcasting to all drivers
01 H : To the driver at address 01
OFH: To the driver at address 15
10H: To the driver at address $16 \ldots$. And so on, maximum to the driver at address 254 (FE H)
Function code:
02 H : Reading several bit devices
03H: Reading several word devices
04H: Reading several word devices (read-only devices)
05 H : Writing a state in a single bit device
06 H : Writing data in a single word device
OFH: Writing states in bit devices
10 H : Writing data in word devices
17H: Reading word devices and writing data in word devices
Data characters: The data sent by the user.
CRC checksum: Starting from Address and ending at Data Content.
Step 1: Make the 16 -bit register (CRC register) $=$ FFFFH
Step 2: Exclusive OR the first 8 -bit message and the low 16 -bit CRC register. Store the result in the CRC register.
Step 3: Right shift CRC register for a bit and fill " 0 " into the high bit.
Step 4: Check the value shifted to the right. If it is 0 , fill in the new value obtained in step 3 and store the value in CRC register; otherwise, Exclusive OR A001H and CRC register and store the result in the CRC register.
Step 5: Repeat step 3-4 and finish operations of all the 8 bits.
Step 6: Repeat step 2-5 for obtaining the next 8-bit message until the operation of all the messages are completed. The final value obtained in the CRC register is the CRC checksum. The CRC checksum has to be placed interchangeably in the checksum of the message.

START and END:
See the table below :

| Baud rate(bps) | RTU timeout timer (ms) | Baud rate (bps) | RTU timeout timer (ms) |
| :---: | :---: | :---: | :---: |
| 300 | 40 | 9,600 | 2 |
| 600 | 21 | 19,200 | 1 |
| 1,200 | 10 | 38,400 | 1 |
| 2,400 | 5 | 57,600 | 1 |
| 4,800 | 3 | 115,200 | 1 |

For example: Read 2 continuous data stored in the registers of the driver at address 01 H (see the table below).
The start register is at address 2102 H .

Inquiry message:

| Name | Data <br> (Hexadecimal <br> value) |
| :---: | :---: |
| Address | 01 H |
| Function code | 03 H |
| Start data address | 21 H |
| Number of data | 02 H |
| (counted by words) | 00 H |
| CRC CHK Low | 62 H |
| CRC CHK High | F 7 H |

Responding message:

| Name | Data <br> (Hexadecimal <br> value) |
| :---: | :---: |
| Address | 01 H |
| Function | 03 H |
| Number of data <br> (counted by byte) | 04 H |
| Content in data address <br> 2102 H | 17 H |
| Content in data address <br> 2103 H | 70 H |
| (CRC CHK Low | 00 H |
| CRC CHK High | FE H |
|  | 5 C H |

6, Timing diagram of the RS-485 communication flag for COM2:




1. Explanations:

- $s$ : Source device
D: Destination device
- The most right digit of $\mathrm{X}, \mathrm{Y}$ and M of $\mathrm{KnX}, \mathrm{KnY}$ and KnM has to be 0 .
- When $S$ designates $K n X$, $D$ has to designate $K n M$; when $S$ designates $K n M$, $D$ has to designate $K n Y$.
- This instruction sends the content in S to $D$ in the form of octal system.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example 1:
When $\mathrm{X} 3=\mathrm{On}$, the content in K 4 X 10 will be sent to K 4 M 10 in octal form.


2, Program Example 2:
When $\mathrm{X} 2=\mathrm{On}$, the content in K 4 M 10 will be sent to K 4 Y 10 in octal form.


| $\begin{aligned} & \text { ZL } \\ & 82 \end{aligned}$ |  | ASCl |  |  |  |  |  |  | S | D | n |  |  |  |  | Converts Hex to ASCll |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | ASCI, ASCIP: 7 steps | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S:Start device for source data

D: Start device for storing the converted result
n: Number of bits to be converted

- Range of n: 1~256
- 16-bit conversion mode: When M1161 = Off, the instruction converts every bit of the hex data in S into ASCII codes and send them to the 8 high bits and 8 low bits of $\mathrm{D} . \mathrm{n}=$ the converted number of bits.
- 8-bit conversion mode: When M1161 = On, the instruction converts every bit of the hex data in S into ASCII codes and send them to the 8 low bits of $\mathrm{D} . \mathrm{n}=$ the number of converted bits. (All 8 high bits of $\mathrm{D}=0$ ).
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2. Program Example 1:

- M1161 = Off: The 16-bit conversion mode.
- When X0 $=$ On, convert the 4 hex values in D10 into ASCII codes and send the result to registers starting from D20.

- Assume:

| (D10) | $=0123 \mathrm{H}$ | '0' $=30 \mathrm{H}$ | '4' $=34 \mathrm{H}$ | ' 8 ' $=38 \mathrm{H}$ |
| :---: | :---: | :---: | :---: | :---: |
| (D11) | $=4567 \mathrm{H}$ | '1' $=31 \mathrm{H}$ | ${ }^{\prime} 5$ ' $=35 \mathrm{H}$ | ' 9 ' $=39 \mathrm{H}$ |
| (D12) | $=89 \mathrm{AB} \mathrm{H}$ | '2' $=32 \mathrm{H}$ | '6' $=36 \mathrm{H}$ | 'A' $=41 \mathrm{H}$ |
| (D13) | $=$ CDEF H | '3' $=33 \mathrm{H}$ | '7' $=37 \mathrm{H}$ | 'B' $=42 \mathrm{H}$ |

- When $n=4$, the bit structure will be as:

- When $\mathrm{n}=6$, the bit structure will be as:

| b1 | $5 \quad \mathrm{D} 10=\mathrm{H} 0123$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | b0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 0 | 0 | 0 | 0 | - 0 | 1 | 10 | 0 | $0 \mid 1$ | 1 | 0 | 0 | 0 | 1 |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| b15 D11 = H 4567 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 0 | 1 | 0 | 0 | 0 | 1 | 0 | D 1 | 10 | 0 | $1{ }^{1} 1$ | 1 | 0 | $0 \mid 1$ | 1 | 1 | 1 |
|  |  |  |  |  |  | 5 |  |  |  | 6 |  |  |  | 7 |  |  |
|  |  |  |  |  |  |  | onve | verte | ed to |  |  |  |  |  |  |  |
| b15 D20 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 0 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 10 | 0 | $0{ }^{0} 1$ | 1 | 1 | $0{ }^{0} 1$ | 1 | 1 | 0 |
| $7 \rightarrow \mathrm{H} 37 \quad \mid \quad 6 \rightarrow \mathrm{H} 36$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| b15 D21 ${ }^{\text {D }}$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 0 | 0 | 1 | 1 | 0 | 0 | 0 | 1 | 10 | 0 | $0 \mid 1$ | 1 | $1{ }^{1} 0$ | 0 | 0 | 0 | 0 |
| ${ }^{\prime}{ }^{\circ} \rightarrow \mathrm{H} 31 \quad \mid \quad 0^{\prime} \rightarrow \mathrm{H} 30$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| b15 D22 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 11 | 10 | 0 | $0 \mid 1$ | 1 | $1{ }^{1} 0$ | 0 | 0 | 1 | 0 |
|  |  |  | $\stackrel{ }{\rightarrow}$ | H |  |  |  |  |  |  |  | $\rightarrow \mathrm{H}$ | H32 |  |  |  |

- When $\mathrm{n}=1 \sim 16$ :

| $\mathrm{D}^{\mathrm{n}}$ | K1 | K2 | K3 | K4 | K5 | K6 | K7 | K8 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| D20 Low byte | "3" | 2" | "1" | "0" | "7" | "6" | "5" | "4" |
| D20 High byte |  | "3" | "2" | "1" | "0" | "7" | "6" | "5" |
| D21 Low byte |  |  | "3" | "2" | "1" | "0" | ${ }^{6}{ }^{\prime \prime}$ | "6" |
| D21 High byte |  |  |  | "3" | "2" | "1" | "0" | ${ }^{6}{ }^{7}$ |
| D22 Low byte |  |  |  |  | "3" | "2" | "1" | "0" |
| D22 High byte |  |  |  |  |  | " ${ }^{\text {n }}$ | "2" | "1" |
| D23 Llow byte |  |  |  |  |  |  | ${ }^{3}{ }^{\prime \prime}$ | "2" |
| D23 High byte |  |  |  |  |  |  |  | ${ }^{3}$ |
| D24 Low byte |  |  |  |  |  |  |  |  |
| D24 High byte |  |  |  |  |  |  |  |  |
| D25 Low byte |  |  |  | chang |  |  |  |  |
| D25 High byte |  |  |  |  |  |  |  |  |
| D26 Low byte |  |  |  |  |  |  |  |  |
| D26 High byte |  |  |  |  |  |  |  |  |
| D27 Low byte |  |  |  |  |  |  |  |  |
| D27 High byte |  |  |  |  |  |  |  |  |

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|  | K9 | K10 | K11 | K12 | K13 | K14 | K15 | K16 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| D20 Low byte | "B" | "A" | "9" | "8" | "F" | "E" | "D" | "C" |
| D20 High byte | "4" | "B" | ${ }^{\text {a }}$ " | "9" | "8" | "F" | E" | D" |
| D21 Low byte | "5" | "4" | "B" | "A" | "9" | "8" | "F" | "E" |
| D21 High byte | "6" | "5" | "4" | "B" | " ${ }^{\text {" }}$ | "9" | "8" | "F" |
| D22 Low byte | "7" | "6" | "5" | "4" | "B" | "A" | "9" | "8" |
| D22 High byte | "0" | "7" | "6" | "5" | "4" | "B" | "A" | "9" |
| D23 Llow byte | "1" | "0" | "7" | "6" | ${ }^{4} 5$ | "4" | "B" | " ${ }^{\text {" }}$ |
| D23 High byte | "2" | "1" | "0" | "7" | ${ }^{5} 6{ }^{\text {" }}$ | "5" | "4" | "B" |
| D24 Low byte | "3" | "2" | "1" | "0" | "7" | ${ }^{4} 6{ }^{\text {a }}$ | "5" | "4" |
| D24 High byte |  | " 3 " | "2" | "1" | "0" | "7" | "6" | "5" |
| D25 Low byte |  |  | " ${ }^{\text {" }}$ | " 2 " | "1" | "0" | ${ }^{4}{ }^{\text {¹ }}$ | "6" |
| D25 High byte |  |  | change | "3" | "2" | "1" | "0" | "7" |
| D26 Low byte |  |  |  |  | "3" | "2" | "1" | "0" |
| D26 High byte |  |  |  |  |  | "3" | "2" | "1" |
| D27 Low byte |  |  |  |  |  |  | ${ }^{3 \prime}$ | 2" |
| D27 High byte |  |  |  |  |  |  |  | "3" |

3, Program Example 2:

- M1161 = On: The 8-bit conversion mode.
- When X0 $=$ On, convert the 4 hex values in D10 into ASCII codes and send the result to registers starting from D20.

- Assume:

| (D10) | $=0123 \mathrm{H}$ | ' 0 ' $=30 \mathrm{H}$ | '4' $=34 \mathrm{H}$ | ' 8 ' $=38 \mathrm{H}$ |
| :---: | :---: | :---: | :---: | :---: |
| (D11) | $=4567 \mathrm{H}$ | '1' $=31 \mathrm{H}$ | $' 5$ ' $=35 \mathrm{H}$ | ' 9 ' $=39 \mathrm{H}$ |
| (D12) | $=89 \mathrm{AB} \mathrm{H}$ | '2' $=32 \mathrm{H}$ | '6' $=36 \mathrm{H}$ | ' $\mathrm{A}^{\prime}=41 \mathrm{H}$ |
| (D13) | $=\mathrm{CDEF} \mathrm{H}$ | ${ }^{\prime} 3$ ' $=33 \mathrm{H}$ | ${ }^{\prime} 7$ ' $=37 \mathrm{H}$ | ' $\mathrm{B}^{\prime}=42 \mathrm{H}$ |

- When $\mathrm{n}=2$, the bit structure will be as:

- When $\mathrm{n}=4$, the bit structure will be as:

- When $\mathrm{n}=1 \sim 16$ :

|  | K1 | K2 | K3 | K4 | K5 | K6 | K7 | K8 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| D20 | "3" | "2" | "1" | "0" | "7" | "6" | "5" | "4" |
| D21 |  | "3" | "2" | "1" | "0" | "7" | "6" | "5" |
| D22 |  |  | "3" | "2" | "1" | "0" | "7" | "6" |
| D23 |  |  |  | "3" | "2" | "1" | "0" | ${ }^{6}{ }^{7}$ |
| D24 |  |  |  | $\begin{gathered} \text { no } \\ \text { chang } \end{gathered}$ | "3" | "2" | "1" | "0" |
| D25 |  |  |  |  |  | "3" | "2" | "1" |
| D26 |  |  |  |  |  |  | "3" | "2" |
| D27 |  |  |  |  |  |  |  | ${ }^{\prime \prime} 3^{\prime \prime}$ |
| D28 |  |  |  |  |  |  |  |  |
| D29 |  |  |  |  |  |  |  |  |
| D30 |  |  |  |  |  |  |  |  |
| D31 |  |  |  |  |  |  |  |  |
| D32 |  |  |  |  |  |  |  |  |
| D33 |  |  |  |  |  |  |  |  |
| D34 |  |  |  |  |  |  |  |  |
| D35 |  |  |  |  |  |  |  |  |


|  | K9 | K10 | K11 | K12 | K13 | K14 | K15 | K16 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| D20 | "B" | "A" | "9" | "8" | "F" | "E" | "D" | "C" |
| D21 | "4" | "B" | "A" | "9" | "8" | "F" | "E" | "D" |
| D22 | "5" | "4" | "B" | "A" | "9" | "8" | "F" | "E" |
| D23 | "6" | "5" | "4" | "B" | "A" | "9" | "8" | "F" |
| D24 | "7" | "6" | "5" | "4" | "B" | "A" | "9" | "8" |
| D25 | "0" | "7" | "6" | "5" | "4" | "B" | " ${ }^{\text {" }}$ | "9" |
| D26 | ${ }^{\text {"1 }}$ | "0" | "7" | "6" | "5" | "4" | "B" | "A" |
| D27 | "2" | "1" | "0" | "7" | "6" | "5" | "4" | "B" |
| D28 | ${ }^{\prime} 3$ " | "2" | "1" | "0" | ${ }^{\text {c }} 7$ " | "6" | "5" | "4" |
| D29 |  | "3" | "2" | "1" | "0" | "7" | "6" | "5" |
| D30 |  |  | "3" | "2" | "1" | "0" | "7" | "6" |
| D31 |  |  | no change | "3" | "2" | "1" | "0" | "7" |
| D32 |  |  |  |  | "3" | "2" | "1" | "0" |
| D33 |  |  |  |  |  | "3" | "2" | "1" |
| D34 |  |  |  |  |  |  | "3" | "2" |
| D35 |  |  |  |  |  |  |  | "3" |


| $\begin{aligned} & \text { ZL } \\ & 83 \end{aligned}$ |  | HEX |  |  |  |  | S |  |  | D |  | n |  |  |  | Converts ASCll to Hex |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | HEX, HEXP: 7 steps | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S:Start device for source data D: Start device for storing the converted result
n : Number of bits to be converted
- Range of $\mathrm{n}: 1 \sim 256$
- 16-bit conversion mode: When M1161 = Off, the instruction is in 16-bit conversion mode. ASCII codes of the 8 high bits and 8 low bits of the hex data in $S$ are converted into hex value and sent to $D$ (every 4 bits as a group). $\mathrm{n}=$ the number of bits converted into ASCII codes.
- 8-bit conversion mode: When M1161 = On, the instruction is in 8-bit conversion mode. Every bit of the hex data in S are converted into ASCII codes and sent to the 8 low bits of $\mathrm{D} . \mathrm{n}=$ the number of converted bits. (All 8 high bits of $\mathrm{D}=0$ ).
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2, Program Example 1:

- M1161 = Off: The 16-bit conversion mode.
- When XO = On, convert the ASCII codes stored in the registers starting from D20 into hex value and send the result (every 4 bits as a group) to registers starting from D10. $n=4$.

- Assume:

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| S | ASCII code | Converted to hex | S | ASCII code | Converted to hex |
| :---: | :---: | :---: | :---: | :---: | :---: |
| D20 low byte | H 43 | "C" | D24 low byte | H 34 | " 4 " |
| D20 high byte | H 44 | "D" | D24 high byte | H 35 | "5" |
| D21 low byte | H 45 | "E" | D25 low byte | H 36 | ${ }^{\text {" } 6}$ |
| D21 high byte | H 46 | "F" | D25 high byte | H 37 | "7" |
| D22 low byte | H 38 | "8" | D26 low byte | H 30 | "0" |

- When $n=4$, the bit structure will be as:

- When $\mathrm{n}=1 \sim 16$ :

|  | D13 | D12 | D11 | D10 |
| :---: | :---: | :---: | :---: | :---: |
| 1 | The undesignated parts in the registers in use are all 0 . |  |  | *** H |
| 2 |  |  |  | **CD H |
| 3 |  |  |  | ${ }^{*} \mathrm{CDE} \mathrm{H}$ |
| 4 |  |  |  | CDEF H |
| 5 |  |  | ${ }^{* * *} \mathrm{CH}$ | DEF8 H |
| 6 |  |  | ${ }^{* *} \mathrm{CDH}$ | EF89 H |
| 7 |  |  | ${ }^{*} \mathrm{CDE} \mathrm{H}$ | F89A H |
| 8 |  |  | CDEF H | 89AB H |
| 9 |  | *** CH | DEF8 H | 9AB4 H |
| 10 |  | **CD H | EF89 H | AB45 H |
| 11 |  | ${ }^{*} \mathrm{CDE} \mathrm{H}$ | F89A H | B456 H |
| 12 |  | CDEF H | 89AB H | 4567 H |
| 13 | ${ }^{* * *} \mathrm{CH}$ | DEF8 H | 9AB4 H | 5670 H |
| 14 | **CD H | EF89 H | AB45 H | 6701 H |
| 15 | *CDE H | F89A H | B456 H | 7012 H |
| 16 | CDEF H | 89AB H | 4567 H | 0123 H |

3. Program Example 2:

- M1161 = On: The 8-bit converstion mode.

- Assume:

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| S | ASCII code | Converted to hex | S | ASCII code | Converted to hex |
| :---: | :---: | :---: | :---: | :---: | :---: |
| D20 | H 43 | "C" | D28 | H 34 | "4" |
| D21 | H 44 | "D" | D29 | H 35 | "5" |
| D22 | H 45 | "E" | D30 | H 36 | "6" |
| D23 | H 46 | *F* | D31 | H 37 | "7" |
| D24 | H 38 | "8" | D32 | H 30 | "0" |
| D25 | H 39 | "9" | D33 | H 31 | "1" |
| D26 | H 41 | " A | D34 | H 32 | -2' |
| D27 | H 42 | "B" | D35 | H 33 | ${ }^{3}{ }^{\prime}$ |

- When $\mathrm{n}=2$, the bit structure will be as:


D21 $\square$
| $44 \mathrm{H} \rightarrow{ }^{\prime} \mathrm{D}^{\prime} \quad$ |
D10


- When $\mathrm{n}=1 \sim 16$ :

|  | D13 | D12 | D11 | D10 |
| :---: | :---: | :---: | :---: | :---: |
| 1 | The used registers which are not specified are all 0 |  |  | *** CH |
| 2 |  |  |  | ${ }^{* *} \mathrm{CD} \mathrm{H}$ |
| 3 |  |  |  | ${ }^{*} \mathrm{CDE} \mathrm{H}$ |
| 4 |  |  |  | CDEF H |
| 5 |  |  | *** CH | DEF8 H |
| 6 |  |  | ${ }^{* *} \mathrm{CD} \mathrm{H}$ | EF89 H |
| 7 |  |  | ${ }^{*} \mathrm{CDE} \mathrm{H}$ | F89A H |
| 8 |  |  | CDEF H | 89 AB H |
| 9 |  | *** CH | DEF8 H | 9AB4 H |
| 10 |  | ${ }^{* *} \mathrm{CD} \mathrm{H}$ | EF89 H | AB45 H |
| 11 |  | *CDE H | F89A H | B456 H |
| 12 |  | CDEF H | 89AB H | 4567 H |
| 13 | ${ }^{* * *} \mathrm{CH}$ | DEF8 H | 9AB4 H | 5670 H |
| 14 | ${ }^{* *} \mathrm{CD} \mathrm{H}$ | EF89 H | AB45 H | 6701 H |
| 15 | ${ }^{*}$ CDE H | F89A H | B456 H | 7012 H |
| 16 | CDEF H | 89AB H | 4567 H | 0123 H |



1. Explanations:

- D: Device of the absolute value.
- This instruction obtains the absolute value of the content in the designated in $D$.
- This instruction adopts pulse execution instructions (ABSP, DABSP).
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.

2. Program Example:

- When $\mathrm{XO}=\mathrm{Off} \rightarrow \mathrm{On}$, obtain the absolute value of the content in DO.



1. Explanations:

- S1: Set value (SV) S2: Present value (PV) S3: Parameter D: Output value (MV)
- In the 16-bit instruction, S 3 will occupy 20 consecutive devices; in the 32-bit instruction, S 3 will occupy 21 consecutive devices.
- This instruction is specifically for PID control. PID operation will be executed by the scan only when the sampling time is reached. PID refers to "proportion, integration and differential" . PID control is widely applied to many machines, pneumatic and electronic equipments.
- For the 16-bit instruction, the parameters are S3 ~ S3+19; for the 32-bit instruction, the parameters are S3 ~ $S 3+20$. After all the parameters are set up, PID instruction will start to be executed and the results will be stored in D. D has to be the data register area without latched function. (If you wish to designate a latched data register area, place the data register in the latched area at the beginning of the program and clear it as 0. )
- Application examples

Use PID commands in temperature control system.
Control purpose: make the control system reach the temperature target value.
The program example is as follows:

Temperature heating program

| $\begin{gathered} \text { M1000 } \\ H \end{gathered}$ | PID | D214 | D220 | D420 | D40 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Normally o pen contac |  | set temper ature | measure te mperature | S3:TC0 sam TC0 output pling time value |  |
|  |  | MUL | D420 | K10 | D426 |
|  |  |  | S3:TC0 sam pling time | TC0 sampli ng time |  |
|  |  |  | MOV | D426 | D41 |
|  |  |  | TCO sampli ng time |  | TCO cycle |
|  |  | GPW/M | D40 | D41 | Y34 |
|  |  | TCO output TCO cycle value |  |  |  |

Note: The sampling time is set to 2.8 S , the default D424 is K4 when power on, when M1 is ON, D424 is K3, and the system enters the temperature auto-tuning mode. After auto-tuning, the value of D424 changes from K3 $\rightarrow$ K4.

2, Other Example:

- Complete the parameter setting before executing PID instruction.
- When $\mathrm{X0}=\mathrm{On}$, the instruction will be executed and the result will be stored in D150. When X0 goes Off, the instruction will not be executed and the data prior to the instruction will stay intact.


Remarks:

- There is no limitation on the times of using this instruction. However, the register No. designated in S3 cannot be repeated.
- For the 16-bit instruction, S3 will occupy 20 registers. In the program example above, the area designated in S3 is D100~D1 19. Before the execution of PID instruction, you have to transmit the setting value to the designated register area by MOV instruction, If the designated registers are latched, use MOVP instruction to transmit all setting value at a time.
- Settings of S3 in the 16-bit instruction:

| Device No. | Function | Setup Range | Explanation |
| :---: | :---: | :---: | :---: |
| $S_{3}$ | Sampling time ( $\mathrm{T}_{\mathrm{s}}$ ) <br> (unit: 10ms) | $\begin{aligned} & \text { 1~2,000 } \\ & \text { (unit: 10ms) } \end{aligned}$ | If $T S$ is less than 1 program scan time, PID instruction will be executed for 1 program scan time. If TS = 0, PID instruction will not be enabled. The minimum TS has to be longer than the program scan time |
| $S_{3}+1$ | Proportional gain ( $\mathrm{K}_{\mathrm{P}}$ ) | 0~30,000(\%) | The magnified error proportional value between SV - PV |
| $\mathrm{S}_{3}+2$ : | Integral gain ( $\mathrm{K}_{1}$ ) | 0~30,000(\%) | For control mode KO~K8 |
| $\mathrm{S}_{3}+3$ : | Differential gain ( $\mathrm{K}_{\mathrm{D}}$ ) | -30,000~30,000(\%) | For control mode K0~K8 |
| $S_{3}+4$ : | Control mode | 0: automatic control <br> 1: forward control ( $\mathrm{E}=\mathrm{SV}-\mathrm{PV}$ ) <br> 2: inverse control ( $\mathrm{E}=\mathrm{PV}-\mathrm{SV}$ ) <br> 3: Auto-tuning of parameter exclusively for the temperature control. The device will automatically become K 4 when the autotuning is completed and be filled in with the appropriate parameter KP, KI and KD (not avaliable in the 32-bit instruction). <br> 4: Exclusively for the adjusted temperature control (not avaliable in the 32-bit instruction) <br> 5: Auto direction control (limited integrall upper/lower limit) <br> 7: Manual control 1: Users set an MV. The accumulated integral value increases according to the error. It is suggested that the control mode should be used in a control environment which changes more slowly. <br> 8: Manual control 2: Users set an MV. The accumulated integral value will stop increasing. When the control mode becomes the |  |

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|  |  | automatic mode (the control mode K5 is used), the instruction PID outputs an appropriate accumulated integral value according to the last MV. |  |
| :---: | :---: | :---: | :---: |
| $\mathrm{S}_{3}+5$ : | The range that error value (E) doesn' $\dagger$ work | 0~32,767 | $E=$ the error of $S V-P V$. When $S 3+5=K 0$, the function will not be enabled, e.g. when $\$ 3+5$ is set as $5, \mathrm{MV}$ of E between -5 and 5 will be 0 . |
| $\mathrm{S}_{3}+6$ : | Upper bound of output value (MV) | -32,768~32,767 | Ex: if $S 3+6$ is set as 1,000 , the output will be 1,000 when MV is bigger than 1,000 . S3 +6 has to be bigger or equal $\$ 3+7$; otherwise the upper bound and lower bound will switch. |
| $S_{3}+7$ : | Lower bound of output value (MV) | -32,768~32,767 | Ex: if $S 3+7$ is set as $-1,000$, the output will be $-1,000$ when MV is smaller than $-1,000$. |
| $\mathrm{S}_{3}+8$ : | Upper bound of integral value | -32,768~32,767 | Ex: if $\$ 3+8$ is set as 1,000 , the output will be 1,000 when the integral value is bigger than 1,000 and the integration will stop. $\mathrm{S} 3+8$ has to be bigger or equal S3 +9; otherwier the upper bound and lower bound will switch. |
| $\mathrm{S}_{3}+9$ : | Lower bound of integral value | -32,768~32,767 | Ex: if $S 3+9$ is set as $-1,000$, the output will be 1,000 when the integral value is smaller than 1,000 and the integration will stop. If $S 3+8$ and S3+9 are set to 0 , there will be no upper limit for integration. |
| $S_{3}+10, ~ 11:$ | Accumulated integral value | 32-bit floating point | The accumulated integral value is only for reference. You can still clear or modify it (in 32bit floating point) according to your need. |
| $\mathrm{S}_{3}+12$ : | The previous PV | -32,768~32,767 | The previous PV is only for reference. You can still modify it according to your need. |
| $\begin{aligned} & S_{3}+13: \\ & \sim \\ & S_{3}+19: \end{aligned}$ | For system use only. |  |  |

- When parameter setting exceeds its range, the upper bound and lower bound will become the setting value. However, if the motion direction (DIR) exceeds the range, it will be set to 0 .
- PID instruction can be used in interruption subroutines, step points and CJ instruction.
- The maximum error of sampling time $T S=-(1$ scan time $+1 \mathrm{~ms}) \sim+(1$ scan time $)$. When the error affects the output, please fix the scan time or execute PID instruction in the interruption subroutine of the timer.
- PV of PID instruction has to be stable before the execution of PID instruction.
- For the 32-bit instruction, If S3 designates the parameter setting area of PID instruction as D100~ D120, S3 occupies 21 registers. Before the execution of PID instruction, you have to use MOV instrction first to send the setting value to the register area for setup. If the designated registers are latched one, use MOVP instruction to send all the setting value at a time.
- Settings of S 3 in the 32-bit instruction:

| Device No. | Function | Setup range | Explanation |
| :---: | :---: | :---: | :---: |
| $S_{3}$ | Sampling time (TS) <br> (unit: 10ms) | $\begin{aligned} & \text { 1~2,000 } \\ & \text { (unit: } 10 \mathrm{~ms} \text { ) } \end{aligned}$ | If $T S$ is less than 1 program scan time, PID instruction will be executed for 1 program scan time. If $T S=0$, PID instruction will not be enabled. The minimum TS has to be longer than the program scan time. |
| $\mathrm{S}_{3}+1$ | Proportional gain $\left(K_{p}\right)$ | 0~30,000(\%) | The magnified error proportional value between SV - PV |
| $S_{3}+2$ : | Integral gain ( $\mathrm{K}_{1}$ ) | 0~30,000(\%) | For control mode K0~K2, K5 |
| $\mathrm{S}_{3}+3$ : | Differential gain ( $\mathrm{K}_{\mathrm{D}}$ ) | -30,000~30,000(\%) | For control mode K0~K2, K5 |
| $S_{3}+4$ : | Control direction (DIR) | automatic control <br> forward control (E=SV-PV) <br> 2: inverse control (E=PV-SV) <br> 5: Automatic mode with MV upper/lower bound control. When MV reaches upper/lower bound, the accumulation of integral value stops. |  |
| $S_{3}+5,6:$ | The range that 32-bit error value (E) doesn' † work | 0~2,147,483,647 | $E=$ the error of SV -PV. When $S 3+5,6=K 0$, the function will not be enabled, e.g. when $\$ 3+5,6$ is set as $5, \mathrm{MV}$ of E between -5 and 5 will be 0 . |
| $\mathrm{S}_{3}+$ :7, 8: | Upper bound of 32bit output value (MV) | $\begin{array}{\|l} -2,147,483,648 ~ \\ 2,147,483,647 \end{array}$ | Ex: if $\$ 3+7,8$ is set as 1,000 , the output will be 1,000 when MV is bigger than $1,000.53+7,8$ has to be bigger or equal $\$ 3+9,10$; otherwise the upper bound and lower bound will switch. |
| S + 9, 10: | Lower bound of 32bit output value (MV) | $\begin{array}{\|l} -2,147,483,648 ~ \\ 2,147,483,647 \end{array}$ | Ex: if $S 3+9,10$ is set as $-1,000$, the output will be 1,000 when MV is smaller than $-1,000$. |
| $S_{3}+11,12:$ | Upper bound of 32bit integral value | $\begin{array}{\|l} -2,147,483,648 ~ \\ 2,147,483,647 \end{array}$ | Ex: if $S 3+11,12$ is set as 1,000 , the output will be 1,000 when the integral value is bigger than 1,000 and the integration will stop. $53+11,12$ has to be bigger or equal S3 $+13,14$; otherwier the |

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|  |  |  | upper bound and lower bound will switch. |
| :---: | :---: | :---: | :---: |
| $\mathrm{S}_{3}+13,14:$ | Lower bound of 32 bit integral value | $\begin{aligned} & -2,147,483,648 \sim \\ & 2,147,483,647 \end{aligned}$ | Ex: if $S 3+13,14$ is set as $-1,000$, the output will be 1,000 when the integral value is smaller than 1,000 and the integration will stop. |
| $\mathrm{S}_{3}+15,16:$ | 32-bit accumulated integral value | 32-bit floating point | The accumulated integral value is only for reference. You can still clear or modify it (in 32bit floating point) according to your need. |
| $\mathrm{S}_{3}+17,18:$ | 32-bit previous PV | -- | The previous PV is only for reference. You can still modify it according to your need. |
| $\mathrm{S}_{3}+19, ~ 20:$ | For system use only. |  |  |

The explanation of 32 -bit S3 and 16 -bit S3 are almost the same. The difference is the capacity of S3+5 ~ S3+20.

## 3. PID Equations:

- The PID operation is conducted according to the speed and the differential PV.
- The PID operation has three control directions: automatic, foreward and inverse. Forward or inverse are, designated in $\mathrm{S} 3+4$. Other relevant settings of PID operation are set by the registers designated in $\mathrm{S3} \sim \mathrm{S3}$ +5 .
- Basic PID equation:

$$
M V=K_{P}{ }^{*} E(t)+K_{I} * E(t) \frac{1}{S}+K_{D}{ }^{*} P V(t) S
$$

| Control direction | PID equation |
| :--- | :--- |
| Forward, automatic | $\mathrm{E}(\mathrm{t})=\mathrm{SV}-\mathrm{PV}$ |
| Inverse | $\mathrm{E}(\mathrm{t})=\mathrm{PV}-\mathrm{SV}$ |

$P V(t) S$ is the differential value of $P V(t) ; E(t) \frac{1}{S}$ is the integral value of $E(t)$.
When $\mathrm{E}(\mathrm{t})$ is less than 0 as the control direction is selected as forward or inverse, $\mathrm{E}(\mathrm{t})$ will be regarded as " 0 ".

The equation above illustrates that this instruction is different from a general PID instruction by the variable use of the differential value. To avoid the flaw that the transient differential value is too big when a general PID instruction is executed for the first time, our PID instruction monitors the differentiation status of the PV. When the variation of PV is too big, this instruction will reduce the output of MV.

- Symbol explanation:
$M V$ : Output value
$K_{p}$ : Proprotional gain
$E(t)$ : Error value
$P V$ : Present measured value
SV : Target value
$K_{D}$ : Differential gain
$P V(t) S$ : Differential value of $\mathrm{PV}(\mathrm{t})$
$K_{I}$ : Integral gain
$E(t) \frac{1}{S}$ : Integral value of $\mathrm{E}(\mathrm{t})$
- Temperature Control Equation:

When $S_{3}+4$ is $K 3$ and $K 4$, the equation used in diagram 2 (see below) will be changed as:

$$
M V=\frac{1}{K_{P}}\left[E(t)+\frac{1}{K_{I}}\left(E(t) \frac{1}{S}\right)+K_{D} * P V(t) S\right]
$$

In which the error value is fixed as $E(t)=S V-P V$
This equation is exclusively designed for temperature control. Therefore, when the sampling time (TS) is set as 4 seconds (K400), the range of output value (MV) will be KO ~K4,000 and the cycle time of GPWM instruction used together has to be set as 4 seconds (K4000) as well. If you have no idea how to adjust the parameters, you can select K3 (auto-tuning) and after all the parameters are adjusted (the control direction will be automatically set as K4), you can modify your parameters to better ones according to the result of the control.

- Control diagrams:


Diagram 1: $\mathrm{S}_{3}+4=\mathrm{K} 0 \sim \mathrm{~K} 2$
In Diagram 1, S is differentiation, referring to "PV - previous PV / sampling time" . $1 / \mathrm{S}$ is integration, referring to "(previous integral value + error value) $\times$ sampling time".$G(S)$ refers to the device being controlled.


In Diagram 2, $1 / \mathrm{KI}$ and $1 / K P$ refer to "divided by KI" and "divided by KP". Due to that this is exclusively for temperature control, you have to use PID instruction together with GPWM instruction. See Application 3 more details.

- Notes:

1) There are a lot of circumstances where PID instruction can be applied; therefore, please choose the control functions appropriately. For example, when you select parameter auto-tuning for the temperature ( $\mathrm{S} 3+4=\mathrm{K} 3$ ), you cannot use it in a motor control environment in case improper control may occur.
2) When you adjust the three main parameters, $K P, K I$ and $K D(S 3+4=K 0 \sim K 2)$, you have to adjust $K P$ first (according to your experiences) and set KI and KD as 0 . When you can roughly handle the control, you then adjust KI (increasingly) and KD (increasingly) (see example 4 below for how to adjust). $\mathrm{KP}=100$ refers to $100 \%$, i.e. the gain of the error is $1 . \mathrm{KP}<100 \%$ will decrease the error and $K P>100 \%$ will increase the error.
3) When you select the parameter exclusively for temperature control ( $S_{3}+4=K 3, K 4$ ), it is suggested that you store the parameter in D register in the latched area in case the automatically adjusted parameter will disappear after the power is cut off. There is no guarantee that the adjusted parameter is suitable for every control. Therefore, you can modify the adjusted parameter according to your actual need, but it is suggested that you modify only $K_{1}$ or $K_{D}$
4) PID instruction can to work with many parameters; therefore please do not randomly modify the parameters in case the control cannot be executed normally.
3. Example 1: Diagram of using PID instruction in position control ( $S 3+4=0$ )


Example 2: Diagram of using PID instruction with AC motor drive on the control ( $S 3+4=0$ )


Example 3: Diagram of using PID instruction in temperature control ( $\mathrm{S} 3+4=1$ )


Example 4: How to adjust PID parameters
Assume that the transfer function of the controlled device $\mathrm{G}(\mathrm{S})$ in a control system is a first-order function $G(s)=\frac{b}{s+a}$ (most models of motors are first-order function), $\mathrm{SV}=1$, and sampling time $\left(\mathrm{T}_{\mathrm{s}}\right)=10 \mathrm{~ms}$, we suggest you to follow the steps below for adjusting the parameters.
Step 1. Set $K_{l}$ and $K_{D}$ as 0 and $K_{P}$ as $5,10,20$ and 40 . Record the $S V$ and $P V$ respectively and the results are as the figure below.


Step 2. From the figure, we can see that when $K_{p}=40$, there will be over-reaction, so we will not select it. When $K_{P}=20$, the PV reaction curve will be close to SV and there will not be over-reaction, but due to its fast start-up with big transient MV, we will consider to put it aside. When $K_{P}=10$, the $P V$ reaction curve will get close to SV value more smoothly, so we will use it. Finally when $K_{p}=5$, we will not consider it due to the slow reaction.
Step 3. Select $K_{p}=10$ and adjust $K_{1}$ from small to big (e.g. 1, 2, 4 to 8). $K_{1}$ should not be bigger than $K_{p}$. Adjust $K_{D}$ from small to big (e.g. $0.01,0.05,0.1$ and 0.2 ). $K_{D}$ should not exceed $10 \%$ of $K_{p}$. Finally we obtain the figure of PV and SV below.


Note: This example is only for your reference. Please adjust your parameters to proper ones according to your actual condition of the control system.

## 8 Application Instructions ZL100~ZL149

## 8.1 (ZL 100-109) Communication instructions



1. Explanations:

- S1: Address of communication device S2: Address of data to be read $n$ : Length of read data
- Range of S1: K0~K254

Range of $\mathrm{n}: \mathrm{Kl} \leqslant \mathrm{n} \leqslant \mathrm{K} 6$

- MODRD is a drive instruction exclusively for peripheral communication equipment in MODBUS ASCII mode/RTU mode.
- If the address of S 2 is illegal to the designed communication device, the device will respond with an error, PLC will records the error code in D1130 and M1141 will be On.
- The feedback (returned) data from the peripheral equipment will be stored in D1070 ~ D1085. After receiving the feedback data is completed, PLC will auto-check if all data are correct. If there is an error, M1140 will be On.
- In ASCII mode, due to that the feedback data are all in ASCII, PLC will convert the feedback data into numerals and store them in D1050 ~ D1055. D1050 ~ D1055 will be invalid in RTU mode.
- After M1140 or M1141 turn On, the program will send a correct datum to the peripheral equipment. If the feedback datum is correct, M1140 and M1141 will be reset.

2. Program Example 1:

Communication between PLC and AC motor drives (ASCII Mode, M1143 = Off)


PLC $\rightarrow$ AC motor drives, PLC sends: "01 0321010006 D4"
AC motor drives $\rightarrow$ PLC, PLC receives: "01 $030 C 010017660000000001360000$ 3B"
Registers for sent data (sending messages)

| Register | DATA |  |  | Explanation |
| :---: | :---: | :---: | :---: | :---: |
| D1089 low | '0' | 30 H | ADR 1 | Address of AC motor drive: ADR $(1,0)$ |
| D1089 high | '1' | 31 H | ADR 0 |  |
| D1090 low | '0' | 30 H | CMD 1 | $\begin{aligned} & \text { Instruction code: CMD } \\ & (1,0) \end{aligned}$ |
| D1090 high | '3' | 33 H | CMD 0 |  |
| D1091 low | '2' | 32 H | Starting data address |  |
| D1091 high | '1' | 31 H |  |  |  |
| D1092 Iow | '0' | 30 H |  |  |  |
| D1092 high | '1' | 31 H |  |  |  |
| D1093 low | '0' | 30 H |  |  |
| D1093 high | '0' | 30 H | Number of d | (counted by words) |
| D1094 low | '0' | 30 H | Number of d | (counted by words) |
| D1094 high | '6' | 36 H |  |  |
| D1095 low | 'D' | 44 H | LRC CHK 1 | Checksum: LRC CHK |
| D1095 high | '4' | 34 H | LRC CHK 0 | $(0,1)$ |

Registers for received data (responding messages)

| Register | DATA |  | Explanation |  |
| :---: | :---: | :---: | :---: | :---: |
| D1070 low | '0' | 30 H | ADR 1 |  |
| D1070 high | '1' | 31 H | ADR 0 |  |
| D1071 low | '0' | 30 H | CMD 1 |  |
| D1071 high | '3' | 33 H | CMD 0 |  |
| D1072 low | '0' | 30 H | Number of data (counted by byte) |  |
| D1072 high | 'C' | 43 H |  |  |
| D1073 low | '0' | 30 H | Content of address 2101 H | PLC automatically convert ASCII codes to numerals and store the numeral in$\text { D1050 }=0100 \mathrm{H}$ |
| D1073 high | '1' | 31 H |  |  |
| D1074 low | '0' | 30 H |  |  |
| D1074 high | '0' | 30 H |  |  |
| D1075 low | '1' | 31 H | Content of address 2102 H | PLC automatically convert ASCII codes to numerals and store the numeral in$\text { D1051 }=1766 \mathrm{H}$ |
| D1075 high | '7' | 37 H |  |  |
| D1076 low | '6' | 36 H |  |  |
| D1076 high | '6' | 36 H |  |  |
| D1077 low | '0' | 30 H | Content of address 2103 H | PLC automatically convert ASCll codes to numerals and store the numeral in$\text { D1052 }=0000 \mathrm{H}$ |
| D1077 high | '0' | 30 H |  |  |
| D1078 low | '0' | 30 H |  |  |
| D1078 high | '0' | 30 H |  |  |
| D1079 low | '0' | 30 H | Content of address 2104 H | PLC automatically convert ASCII codes to numerals and store the numeral in$\mathrm{D} 1053=0000 \mathrm{H}$ |
| D1079 high | '0' | 30 H |  |  |
| D1080 low | '0' | 30 H |  |  |
| D1080 high | '0' | 30 H |  |  |
| D1081 low | '0' | 30 H | Content of address 2105 H | PLC automatically convert ASCll codes to numerals and store the numeral in D1054 $=0136 \mathrm{H}$ |
| D1081 high | '1' | 31 H |  |  |
| D1082 low | '3' | 33 H |  |  |
| D1082 high | '6' | 36 H |  |  |


| D1083 low | ' 0 ' | 30 H |  | PLC automatically convert <br> ASCII codes to numerals <br> and store the numeral in |
| :--- | :---: | :---: | :--- | :--- |
| D1083 high | ' ${ }^{\prime}$ | 30 H | Content of |  |
| address 2106 H | D1055 $=0000 \mathrm{H}$ |  |  |  |
| D1084 low | ' 0 | 30 H |  |  |
| D1084 high | ' 0 ' | 30 H |  |  |
| D1085 low | '3' | 33 H | LRC CHK 1 |  |
| D1085 high | 'B' | 42 H | LRC CHK 0 |  |

## 3. Program Example 2:

Communication between PLC and AC motor drives (RTU Mode, M1143=On)


PLC $\rightarrow$ AC motor drives, PLC sends: "01 0321020002 6F F7"
AC motor drives $\rightarrow$ PLC ,PLC receives: "01 030417700000 FE 5C"

Registers for sent data (sending messages)

| Register | DATA |  |
| :--- | :---: | :--- |
| D1089 low | 01 H | Address |
| D1090 low | 03 H | Function |
| D1091 low | 21 H | Starting data address |
| D1092 low | 02 H |  |
| D1093 low | 00 H | Number of data (counted by words) |
| D1094 low | 02 H |  |
| D1095 low | 6 F H | CRC CHK Low |
| D1096 low | F7 H | CRC CHK High |

Registers for received data (responding messages)

| Register | DATA | Explanation |
| :--- | :---: | :--- |
| D1070 low | 01 H | Address |
| D1071 low | 03 H | Function |
| D1072 low | 04 H | Number of data (counted by bytes) |
| D1073 low | 17 H | Content of address 2102 H |
| D1074 low | 70 H |  |
| D1075 low | 00 H | Content of address 2103 H |
| D1076 low | 00 H |  |
| D1077 low | FE H | CRC CHK Low |
| D1078 low | 5 C H | CRC CHK High |



1, Explanations:

- $\quad$ S1: Address of communication device $\quad$ S2: Address of data to be read $n$ : Data to be written
- Range of S1: K0~K254
- MODWR is a drive instruction exclusively for peripheral communication equipment in MODBUS ASCII mode/RTU mode.
- The feedback (returned) data from the peripheral equipment will be stored in D1070 ~ D1076. After receiving the feedback data is completed, M1127 will be On.

2, Program Example 1:
Communication between PLC and AC motor drives (ASCII Mode, M1143 = Off)


PLC $\rightarrow$ AC motor drives, PLC sends: "01 060100177071 "
AC motor drives $\rightarrow$ PLC , PLC receives:"01 $060100177071 "$
Registers for sent data (sending messages)

| Register | DATA |  | Explanation |  |
| :---: | :---: | :---: | :---: | :---: |
| D1089 low | '0' | 30 H | ADR 1 | Address of AC motor drive: |
| D1089 high | '1' | 31 H | ADR 0 | ADR (1,0) |
| D1090 low | '0' | 30 H | CMD 1 |  |
| D1090 high | '6' | 36 H | CMD 0 | Instruction code: CMD (1,0) |
| D1091 low | '0' | 30 H | Data address |  |
| D1091 high | '1' | 31 H |  |  |
| D1092 low | '0' | 30 H |  |  |
| D1092 high | '0' | 30 H |  |  |
| D1093 low | '1' | 31 H | Data contents |  |
| D1093 high | ${ }^{\prime} 7$ ' | 37 H |  |  |
| D1094 low | '7' | 37 H |  |  |
| D1094 high | '0' | 30 H |  |  |
| D1095 low | ${ }^{\prime} 7$ ' | 37 H | LRC CHK 1 | Error checksum: LRC CHK $(0,1)$ |
| D1095 high | '1' | 31 H | LRC CHK 0 |  |

PLC receiving data register (response messages)


| $\begin{gathered} \mathrm{ZL} \\ 102 \end{gathered}$ |  | RS 1 |  |  |  |  |  | $S_{2}$ |  | $S_{3}$ | S 4 |  | n |  |  | COM1: RS232 data read and write |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | + | vic |  | Word Devices |  |  |  |  |  |  |  |  |  |  | RSI: 7 step | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |  |
| $\mathrm{S}_{1}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| $\mathrm{S}_{2}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |

1. Explanations:

- This command has the functions of MODRD, MODWR and MODRW commands at the same time, which is more convenient to use.
- S1: Address of communication device (Unit Address) S2: Communication function code (Function Code).

| Function code | Command description |
| :--- | :--- |
| K1 | Reading several bit devices (same function as MODRD instruction) |
| K3 | Reading single or several word devices (same function as MODRD instruction) |
| K6 | Writing data in a single word device (same function as the MODWR instruction) |
| K15 | Writing states in bit devices (same function as MODRW instruction) |
| K16 | Writing data in word devices (same function as MODRW instruction) |

- S3, S4, n: Its functions vary according to different function codes.
- $\quad$ S3: The address of the data to be read and written.
- S4: Data to be read and written.
- n : read and write data length.
- S3, S4, n operands have the following functions according to different function codes:

| Function code | $\mathrm{S}_{3}$ | $\mathrm{~S}_{4}$ | n |
| :--- | :--- | :--- | :--- |
| K1 (Reading several bit devices) | Address from which data <br> is read | Register where data read is <br> stored | Number of data read |
| K3 (Reading single or several | Address from which data <br> is read | Register where data read is <br> stored | Number of data read |
| K6 (Writing data in a single word | Address from which data <br> is read | Data register where data <br> written is stored | None |

8 Application Instructions ZL100~ZL149

| K15 (Writing states in bit devices) | Address into which data <br> is written | Data register where data <br> written is stored | Number of <br> written |
| :--- | :--- | :--- | :--- | :--- |
| K16 (Writing data in word | Address into which data |  |  |
| devices) | Data register where data <br> written is stored | Number of <br> written |  |

- There is no limit to the number of times this command can be used in the program, and multiple commands can be executed at the same time.

Eg.


- pecial M

| special M | function |
| :--- | :--- |
| M1200 | ON when the RS1/RS2/RS3 command communication is successful, the system will <br> automatically OFF every time it is ON |
| M1201 | ON when the RS 1/RS2/RS3 command communication fails, the system will automatically OFF <br> every time it is ON |


| $\begin{gathered} \text { ZL } \\ 103 \end{gathered}$ |  | RS2 |  |  |  |  | $S_{1}$ | $S_{2}$ |  | $S_{3}$ | $S_{4}$ |  | n |  |  | COM2: RS485data read and write |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | RS2: 7 step | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| $\mathrm{S}_{1}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| $\mathrm{S}_{2}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |

1. Instruction description: refer to ZL102 RS 1 instruction description

| $\begin{gathered} \text { ZL } \\ 104 \end{gathered}$ |  | RS3 |  |  |  |  | S | $S_{2}$ | $S_{3}$ |  | $S_{4}$ |  | n |  |  | COM3: RS485 data read and write |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | RS3: 7 steps | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| $S_{1}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| $\mathrm{S}_{2}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |

1. Instruction description: refer to ZL102 RS1 instruction description

| $\begin{gathered} \text { ZL } \\ 108 \end{gathered}$ |  | CRC |  |  |  |  |  |  | n |  | D |  |  |  |  | Checksum CRC Mode |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | CRC, CRCP: 7 steps | 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E |  |  |  |
| S |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |

1, Explanations:

- S: Start operation device for RTU mode checksum

D: Start device for storing the operation result
n : Number of calculated bits
CRC checksum: See remarks

- Range of $\mathrm{n}: \mathrm{Kl} \sim \mathrm{K} 256$
- If n does not fall within its range, an operation error will occur, the instruction will not be executed, M1067, M1068 $=$ On and D1067 will record the error code H' OEIA.
- In 16-bit conversion mode: When M1161 = Off, S divides its hex data area into higher 8 bits and lower 8 bits and performs CRC checksum operation on each bit. The data will be sent to the higher 8 bits and lower 8 bits in $\mathrm{D} . \mathrm{n}=$ the number of calculated bits.
- In 8-bit conversion mode: When M1161 $=$ On, S divides its hex data area into higher 8 bits (invalid data) and lower 8 bits and performs CRC checksum operation on each bit. The data will be sent to the lower 8 bits in D and occupy 2 registers. $\mathrm{n}=$ the number of calculated bits. (All higher 8 bits in D are " 0 ".)

2, Program Example:

- When PLC communicates with AC motor drives (In RTU mode, M1143=On), (In 16-bit mode, M1161 = On), the sent data write in advance H 12 into H 2000 of AC motor drives.


PLC $\rightarrow$ AC motor drives, PLC sends: "01 062000001202 07"
Registers for sent data (sending messages)

| Register | DATA | Explanation |
| :--- | :---: | :--- |
| D100 low | 01 H | Address |
| D101 low | 06 H | Function |
| D102 low | 20 H | Data address |
| D103 low | 00 H |  |
| D104 low | 00 H | Data content |
| D105 low | 12 H |  |
| D106 low | 02 H | CRC CHK 0 |
| D107 low | 07 H | CRC CHK 1 |

The error checksum $\operatorname{CRC} \operatorname{CHK}(0,1)$ can be calculated by CRC instruction (in 8-bit mode, M1161 = On).

| M1000 | CRC | D100 | K6 | D106 |
| :--- | :--- | :--- | :--- | :--- |

CRC checksum: 02 H is stored in the lower 8 bits of D 106 and 07 H in the lower 8 bits of D107

Remarks:

1) The format of RTU mode with a communication datum:

| START | Time interval |
| :--- | :--- |
| Address | Communication address: 8-bit binary |
| Function | Function code: 8-bit binary |
| DATA $(\mathrm{n}-1)$ | Data content: |


| $\ldots . .$. | $\mathrm{n} \times 8$-bit data |
| :--- | :--- |
| DATA 0 |  |
| CRC CHK Low | 16 -bit CRC checksum consists of 2 8-bit binaries |
| CRC CHK High | Time interval |
| END |  |

2) CRC checksum starts from Address and ends at Data content.

The operation of CRC checksum:
Step 1: Make the 16-bit register (CRC register) $=$ FFFFH
Step 2: Exclusive OR the first 8-bit byte message instruction and the low-bit 16-bit CRC register. Store the result in CRC register.

Step 3: Shift the CRC register one bit to the right and fill 0 in the higher bit.
Step 4: Check the value that shifts to the right. If it is 0 , store the new value from Step 3 into the CRC register, otherwise, Exclusive OR A 001 H and the CRC register, and store the result in the CRC register.

Step 5: Repeat Step $3 \sim 4$ and finish calculating the 8 bits.
Step 6: Repeat Steps $2 \sim 5$ for obtaining the next 8-bit message instruction until all the message instructions are calculated. In the end, the obtained CRC register value is the CRC checksum. Be aware that CRC checksum should be placed in the checksum of the message instruction.

| $\begin{gathered} \text { ZL } \\ 109 \end{gathered}$ |  | SWRD |  |  |  |  | D |  |  |  |  |  |  |  |  | Read Digital Switch |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | SWRD: 3 steps | 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |  |  |

1. Explanations:

- D: Device for storing the read value.
- This instruction stores the value read from digital switch function card into D.
- The read value is stored in the low byte in D. Every switch has a corresponding bit.
- When there is no digital function card inserted, the error message C400 (hex) will appear in grammar check.

2. Program Example:

- There are 18 DIP switches on the digital switch function card. After the switches are read by SWRD instruction, the status of each switch will correspond to MO ~M7.

- The status of MO ~M7 can be executed by each contact instruction.
- The execution of END instruction indicates that the process of input is completed. REF (I/O refresh) instruction will be invalid.
- When SWRD instruction uses the data in digital switch function card, it can read minimum 4 bits (KIY*, K1M* or K1S*).


## Remarks:

When digital switch function card is inserted, the status of the 8 DIP switches will correspond to M1104~M1111。

## 8.2 (ZL 110-119) Floating point arithmetic

| $\begin{gathered} \text { ZL } \\ 110 \end{gathered}$ | D | ECMP |  |  |  |  | $S_{1} S_{2} \quad D$ |  |  |  |  |  |  |  |  | Floating Point Compare |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DECMP, DECMPP: 13 step | 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| $\mathrm{S}_{1}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| $\mathrm{S}_{2}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S 1 : Binary floating point comparison value 1

S2: Binary floating point comparison value 2
D: Comparison result

- D occupies 3 consecutive devices.
- The binary floating point values S 1 and S 2 are compared with each other. The comparison result ( $>,=,<$ ) is stored in D .
- If S 1 or S 2 is an designated constant K or H , the instruction will convert the constant into a binary floating point value before the comparison.


## 2, Program Example:

- Designated device M10 and M10~M12 are automatically occupied.
- When XO = On. DECMP instruction will be executed and one of M10 ~M12 will be On. When XO $=$ Off, DECMP instruction will not be executed and M10 ~ M12 will remain their status before X0 $=$ Off.
- To obtain results $\geqq, \leqq, \neq$ serial-parallel M10~M12.
- Use RST or ZRST instruction to clear the result.


| $\begin{gathered} \mathrm{ZL} \\ 111 \end{gathered}$ |  | EZCP |  |  |  |  |  |  | $S_{2}$ |  | D |  |  |  |  | Floating Point Zone Compare |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DEZCP: 17 step | 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  | KnM | Kns | T | C | D | E |  |  | F |
| $\mathrm{S}_{1}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| $\mathrm{S}_{2}$ |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| S |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S1: Lower bound of binary floating point

S: Binary floating point comparison result

S2: Upper bound of binary floating point
D: Comparison result

- D occupied 3 consecutive devices.
- $S 1 \leqslant S 2$. See the specifications of each model for their range of use.
- $S$ is compared with S1 and S2 and the result ( $(>,=,<)$ is stored in D.
- If S 1 or S 2 is andesignated constant K or H , the instruction will convert the constant into a binary floating point value before the comparison.
- When S1 > S2, S1 will be used as upper/lower bound for the comparison.

2, Program Example:

- Designated device $M 0$ and $M 0 \sim M 2$ are automatically occupied.
- When $X 0=$ On. DEZCP instruction will be executed and one of $M 0 \sim M 2$ will be On. When X0 = Off, EZCP instruction will not be executed and $M 0 \sim M 2$ will remain their status before $X 0=O f f$.
- Use RST or ZRST instruction to clear the result.



1. Explanations:

- $S$ : Source floating point data
D: Destination device
- $S$ can only be a floating-point constant value.
- This instruction is able to enter floating point values directly in S .
- When the instruction is executed, the content in $S$ is moved directly into $D$. When the instruction is not executed, the content in $D$ will not be modified.
- If users want to move the floating-point value in registers, they have to use DMOV.
- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.


## 2. Program Example:

- User DMOVR instruction to move 32-bit floating point data.
- When XO = Off, the content in (DII, D10) remains unchanged. When XO $=$ On, the present value F1. 20000004768372 will be moved to data registers (D11, D10).


| $\begin{gathered} \text { ZL } \\ 114 \end{gathered}$ | D | MUL16 <br> MUL32 |  |  |  |  | S |  | S2 |  | D |  |  |  |  | 16-bit Multiplication <br> 32-bit Multiplication |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | MUL : 7 steps <br> DMUL: 13 steps |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S1:Multiplicand S2:Multiplicator D: Product
- In 16-bit instruction, D occupies one device.
- In 32-bit instruction, D occupies 2 consecutive devices.
- This instruction multiplies S1 by S2 in BIN format and stores the result in D. Be careful with the positive/negative signs of $\mathrm{S} 1, \mathrm{~S} 2$ and D when doing 16 -bit and 32 -bit operations.
- In 16-bit BIN multiplication.
(S1)
(S2)
(D)
b15 .............b0 $\times$ b15.............b0 $=\mathrm{b} 15 \ldots \ldots \ldots \ldots \ldots \mathrm{c} 0$
b15 is a symbol bit. b15 is a symbol bit. b15 is a symbolbit.
16 bits $\times 16$ bits $=16$ bits
Symbol bit $=0$ refers to a positive value.
Symbol bit = 1 refers to a negative value.
When D serves as a bit device, it can designate K1 ~K4 and construct a 16-bit result, occupying 16-bit data.
- In 32-bit BIN multiplication.


32 bits $\times 32$ bits $=32$ bits
Symbol bit $=0$ refers to a positive value.
Symbol bit = 1 refers to a negative value.
2. Program Example 1:

- If MO is On, the 16-bit DO is multiplied by the 16 -bit D10 and a 16 -bit product is produced. The 16 -bit data is
stored in D20. On/Off of the most left bit indicates the positive/negative status of the result value.


16 bits $\times 16$ bits $=16$ bits
$D 0 \times D 10=D 20$
$D 0=K 100, D 10=K 200, D 20=K 10,000$

## 3. Program Example 2:

If XO is On, the 32-bit value K10,00 in (D1, D0) is multiplied by the 32-bit value K20,000 in (D1 1, D10) and a 32bit product is produced. The 32-bit data is stored in (D21, D20). On/Off of the most left bit indicates the positive/negative status of the result value.


32 bits $\times 32$ bits $=32$ bits
$(\mathrm{D} 1, \mathrm{D} 0) \times(\mathrm{D} 11, \mathrm{D} 10)=(\mathrm{D} 21, \mathrm{D} 20)$
$(D 1, D 0)=K 10,000,(D 11, D 10)=K 20,000,(D 21, D 20)=K 200,000,000$

Remarks:

1) If the value gotten from the 16 -bit multiplication can not be represented by a 16 -bit signed value, and is greater than the maximum 16-bit positive value K 32767 or less than the minimum 16-bit negative value K-32768, the low 16-bit data is stored.
2) If users need to get a complete value (32-bit value) from a 16 -bit multiplication, they have to use API22 MUL/MULP. Please refer to the explanation of API22 MUL/MULP for more information.
3) If the value gotten from the 32-bit multiplication can not be represented by a 32-bit signed value, and is greater than the maximum 32-bit positive value K2147483647 or less than the minimum 16bit negative value K-2147483648, the low 32-bit data is stored.
4) If users need to get a complete value (64-bit value) from a 32bit multiplication, they have to use ZL 22 DMUL/DMULP. Please refer to the explanation of API22 DMUL/DMULP for more information.

| $\begin{gathered} \text { ZL } \\ 115 \end{gathered}$ |  | DIV16 <br> DIV32 |  |  |  |  | S |  | S2 |  | D |  |  |  |  | 16-bit Division <br> 32-bit Division |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DIV: 7 steps DDIV: 13 step | $\begin{aligned} & \text { 16-bit } \\ & \text { 32-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S1: Dividend S2: Divisor D: Quotient and remainder.
- In 16-bit instruction, D occupies one device.
- In 32-bit instruction, D occupies 2 consecutive devices.
- This instruction divides $S 1$ and $S 2$ in BIN format and stores the result in D. Be careful with the positive/negative signs of $\mathrm{S} 1, \mathrm{~S} 2$ and D when doing 16 -bit and 32 -bit operations.
- This instruction will not be executed when the divisor is 0 .
- In 16-bit BIN division:

- When D serves as a bit device, it can designate K1 ~ K4 and construct a 16-bit result, occupying a 16-bit quotient.
- In 32-bit BIN division:


When D serves as a bit device, it can designate K1 ~ K8 and construct a 32-bit result, occupying a 32-bit quotient.

2, Program Example 1:

- If $M 0=$ On, the value in DO (K103) will be divided by the value in D10 (K5) and the quotient will be stored in D20. On/Off of the highest bit indicates the positive/negative status of the result value.
D0/D10=D20
$\Rightarrow \mathrm{K} 103 / \mathrm{K} 5=\mathrm{K} 20$. The remainder is K 3.
$\Rightarrow \mathrm{D} 20=\mathrm{K} 20$ (The remainder is left out.)

3. Program Example 2:

- If $M 0=O n$, the value in (D1, D0) $(K 81,000)$ will be divided by the value in (D11, D10) $(K 40,000)$ and the quotient will be stored in (D21, D20). On/Off of the highest bit indicates the positive/negative status of the result value.

(D1,D0)(D11,D10)=(D21,D20)
$\Rightarrow K 81,000 / \mathrm{K} 40,000=\mathrm{K} 2$. The remainder is $\mathrm{K} 1,000$.
$\Rightarrow \quad(\mathrm{D} 21, \mathrm{D} 20)=\mathrm{K} 2$ (The remainder is left out.)
Remarks:
- If users need to record a remainder by a 16-bit division, they have to use API23 DIV/DIVP. Please refer to the explanation of API23 DIV/DIVP for more information.
- If users need to record a remainder by a 32-bit division, they have to use API23 DDIV/DDIVP. Please refer to the explanation of API23 DDIV/DDIVP for more information.


1, Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $S$ : Source (angle) D: Result (radian)
- Radian $=$ degree $\times(\pi / 180)$


## 2, Program Example:

- When $\mathrm{XO}=\mathrm{On}$, designate the degree of binary floating point (D1, D0). Convert the angle into radian and store the result in binary floating point in (D11, D10)。


| $\begin{array}{\|c\|} \hline \text { ZL } \\ 117 \end{array}$ |  | DEG |  |  |  |  |  |  |  | D |  |  |  |  |  | Radian $\rightarrow$ Angle |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DDEG, DDEGP: 9 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  | KnM | KnS | T | C | D | E |  | F |
| S |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $\quad$ S: Source (radian) D: Result (angle)
- Degree $=$ radian $\times(180 / \pi)$


## 2, Program Example:

When $\mathrm{XO}=\mathrm{On}$, designate the angle of binary floating point (D1, DO). Convert the radian into angle and store the result in binary floating point in (D11, D10).


| $\begin{gathered} \text { ZL } \\ 118 \end{gathered}$ |  | EBCD |  |  |  |  | S |  | D |  |  |  |  |  |  | Float to Scientific Conversion |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DEBCD, DEBCDP: 9 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  |  |  | KnM | KnS | T | C |  | D | E | F |
| S |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- S:Source D:Result
- This instruction converts binary floating point value in the register designated by $S$ into decimal floating point value and stores it in the register designated by $D$.
- PLC conducts floating point operation in binary format. DEBCD instruction is exclusively for converting floating points from binary to decimal.

2. Program Example:

- When $\mathrm{XO}=\mathrm{On}$, the binary floating points in D1 and DO will be converted into decimal floating points and stored in D3 and D2.



1, Explanations:

- S:Source D: Result
- This instruction converts decimal floating point value in the register designated by $S$ into binary floating point value and stores it in the register designated by D.
- DEBIN instruction is exclusively for converting floating points from decimal to binary.
- Range of decimal floating point real numbers: -9.999 ~+9.999. Range of exponants: $-41 \sim+35$. Range of PLC decimal floating points: $\pm 1,175 \times 10^{-41} \sim \pm 3,402 \times 10^{+35}$

2, Program Example 1:

- When $\mathrm{X} 1=$ On, the decimal floating points in D1 and D0 will be converted into binary floating points and stored in D3 and D2.



## 3. Program Example 2:

- Use FLT instruction (API 149) to convert BIN integer into binary floating point before performing floating point operation. The value to be converted must be BIN integer and use DEBIN instruction to convert the floating point into a binary one.
- When $\mathrm{XO}=$ On, move K3, 140 to D0 and K-3 to D1 to generate decimal floating point $\left(3.14=3140 \times 10^{-3}\right)$ 。

```
*******)
K3140->D0] llll}[\begin{array}{l}{[D1]}\\{3140\times10}
K-3 -> D1 ] [D0]
(D1, D0) }->\mathrm{ (D3, D2)
3140\times10}\mp@subsup{0}{}{-3}->\mathrm{ Binary floating point
```


## 8.3 (ZL 120-129) Floating point arithmetic

| $\begin{gathered} \hline \text { ZL } \\ 120 \\ \hline \end{gathered}$ |  | EADD |  |  |  |  | S 1 S |  |  | 2 D |  |  |  |  |  | Floating Point Addition |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DEADD: 13 steps32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM |  |  |  | KnS | T | c |  | D | E | F |
| S1 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S1:Summand

S2: Addend
D: Sum

- $\quad S 1+S 2=D$. The floating point value in the register designated by $S 1$ and $S 2$ are added up and the result is stored in the register designated by D. The addition is conducted in binary floating point system.
- If S 1 or S 2 is an designated constant K or H , the instruction will convert the constant into a binary floating point value before the operation.
- $S 1$ and $S 2$ can designate the same register. In this case, if the "continuous execution" instruction is in use, during the period when the criteria contact in On, the register will be added once in every scan by pulse execution instruction DEADDP.

2, Program Example 1:

- When $\mathrm{XO}=\mathrm{On}$, binary floating point (D1, D0) + binary floating point (D3, D2) and the result is stored in (D1 1, D10).


3. Program Example 2:

- When X2 = On, binary floating point (D11, D10) + K1234 (automatically converted into binary floating point) and the result is stored in (D21, D20).

| X 2 | DEADD | D10 | K1234 | D20 |
| :---: | :--- | :--- | :--- | :--- |



1, Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S1:Minuend S2:Subtrahend D: Remainder
- $\mathrm{S} 1-\mathrm{S} 2=\mathrm{D}$. The floating point value in the register designated by S 2 is subtracted from the floating point value in the register assigned by $\$ 1$ and the result is stored in the register designated by D . The subtraction is conducted in binary floating point system.
- If S 1 or S 2 is an designated constant K or H , the instruction will convert the constant into a binary floating point value before the operation.
- $S 1$ and $S 2$ can designate the same register. In this case, if the "continuous execution" instruction is in use, during the period when the criteria contact in On, the register will be subtracted once in every scan by pulse execution instruction DESUBP.

2, Program Example 1:

- When X0 = On, binary floating point (D1, D0) - binary floating point (D3, D2) and the result is stored in (D11, D10).

| $\mathrm{X0}$ | DESUB | D0 | D2 | D10 |
| :---: | :--- | :--- | :--- | :--- |

3. Program Example 2:

- When X2 = On, K1234 (automatically converted into binary floating point) - binary floating point (D1, D0) and the result is stored in (D11, D10).


8 Application Instructions ZL100~ZL149

| $\begin{gathered} \text { ZL } \\ 122 \end{gathered}$ |  | EMUL |  |  |  |  |  |  | S2 |  | D |  |  |  |  | Floating Point Multiplication |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DEMUL: 13 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |
| S1 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1, Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S1:Multiplicand S2: Multiplicator D: Product
- $\quad$ S $1 \times S 2=\mathrm{D}$. The floating point value in the register assigned by S 1 is multiplied with the floating point value in the register designated by S2 and the result is stored in the register designated by D. The multiplication is conducted in binary floating point system.
- If S 1 or S 2 is an designated constant K or H , the instruction will convert the constant into a binary floating point value before the operation.
- S1 and S2 can designate the same register. In this case, if the "continuous execution" instruction is in use, during the period when the criteria contact in On, the register will be multiplied once in every scan by pulse execution instruction DEMULP.

2, Program Example :

- When $\mathrm{X} 1=$ On, binary floating point ( $\mathrm{D} 1, \mathrm{D} 0) \times$ binary floating point ( $\mathrm{D} 11, \mathrm{D} 10$ ) and the result is stored in (D21, D20).

| X1 | DEMUL | D0 | D10 | D20 |
| :---: | :--- | :--- | :--- | :--- |

- When X2 $=$ On, K 1234 (automatically converted into binary floating point) $\times$ binary floating point (D1, D0) and the result is stored in (D11, D10).

| X 2 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| H | DEMUL | K1234 | D0 | D10 |


| $\begin{gathered} \mathrm{ZL} \\ 123 \end{gathered}$ |  | EDIV |  |  |  |  |  |  | S2 |  | D |  |  |  |  | Floating Point Division |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DEDIV: 13 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T | C | D | E | F |  |
| S1 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1, Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S1: Dividend S2: Divisor D: Quotient and remainder
- S1 $\div$ S2 = D. The floating point value in the register designated by $S 1$ is divided by the floating point value in the register assigned by $S 2$ and the result is stored in the register designated by $D$. The division is conducted in binary floating point system.
- If $S 1$ or $S 2$ is an designated constant $K$ or $H$, the instruction will convert the constant into a binary floating point value before the operation.
- If $S 2=0$, operation error will occur, the instruction will not be executed.

2, Program Example 1:

- When $\mathrm{X} 1=$ On, binary floating point (D1, D0) $\div$ binary floating point ( $\mathrm{D} 11, \mathrm{D} 10$ ) and the quotient is stored in (D21, D20).

| X 1 | DEDIV | D0 | D10 | D20 |
| :---: | :---: | :---: | :---: | :---: |

3. Program Example 2:

- When X2 = On, binary floating point (D1, D0) $\div$ K1234 (automatically converted into binary floating point) and the result is stored in (D11, D10).



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $S$ : Device for operation source D: Device for operation result
- $\quad e=2.71828$ as the base and $S$ as exponent for EXP operation: EXP ${ }^{[D+1, D]}=[S+1, S]$.
- Both positive and negative values are valid for $S$. When designating $D$ registers, the data should be 32-bit and the operation should be performed in floating point system. Therefore, $S$ should be converted into a floating point value.
- The content in $D=e^{s} ; e=2.71828, S=$ designated source data


## 2, Program Example:

- When M0 = On, convert (D1, D0) into binary floating point and store it in register (D11, D10).
- When $\mathrm{Ml}=$ On, use (D11, D10) as the exponent for EXP operation and store the binary floating point result in register (D21, D20).
- When M2 = On, convert the binary floating point (D21, D20) into decimal floating point (D30 $\times$ 10[D31]) and store it in register (D31, D30).



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $S$ : Device for operation source D: Device for operation result
- This instruction performs natural logarithm "LN" operation by $\mathrm{S}: \mathrm{LN}[\mathrm{S}+1, \mathrm{~S}]=[\mathrm{D}+\mathrm{I}, \mathrm{D}]$.
- Only positive values are valid for $S$. When designating $D$ registers, the data should be 32 -bit and the operation should be performed in floating point system. Therefore, $S$ should be converted into a floating point value.
- $e^{D}=S$. The content in $D=\operatorname{lnS} ; S=$ designated source data.


## 2, Program Example:

- When $M 0=O n$, convert (D1, D0) into binary floating point and store it in register (D11, D10).
- When $M 1=$ On, use register (DII, D10) as the real number for $L N$ operation and store the binary floating point result in register (D21, D20).
- When M2 = On, convert the binary floating point (D21, D20) into decimal floating point (D30 $\times 10^{[D 31]}$ ) and store it in register (D31, D30)


| $\begin{gathered} \hline \text { ZL } \\ 126 \end{gathered}$ |  |  | LOG |  |  |  | S 1 |  |  | S2 |  | D |  |  |  | Logarithm of Binary Floating Point |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DLOG: 13 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- si: Device for base
S2: Device for operation source
D: Device for operation result
- This instruction performs "log" operation of the content in S1 and S2 and stores the result in D.
- Only positives are valid for the content in S 1 and S 2 . When designating D registers, the data should be 32bit and the operation should be performed in floating point system. Therefore, S 1 and S 2 should be converted into floating point values.
- $\quad \mathrm{S} 1^{\mathrm{D}}=\mathrm{S} 2, \mathrm{D}=$ ? $\rightarrow \mathrm{Log}_{5}{ }^{\mathrm{S} 2}=\mathrm{D}$

Example: Assume $\mathrm{S}_{1}=5, \mathrm{~S}_{2}=125, \mathrm{D}=\log _{5}{ }^{125}=$ ?

$$
\mathrm{S}_{1}^{\mathrm{D}}=\mathrm{S}_{2} \rightarrow 5^{\mathrm{D}}=125 \rightarrow \mathrm{D}=\log _{5} 125=3
$$

2, Program Example:

- When $M 0=O$ n, convert (D1, D0) and (D3, D2) into binary floating points and store them in the 32-bit registers (D11, D10) and (D13, D12).
- When M1 = On, perform log operation on the binary floting points in 32-bit registers (D11, D10) and (D13, D12) and store the result in the 32-bit register (D21, D20).
- When M2 = On, convert the binary floating point (D21, D20) into decimal floating point (D30 $\times{ }^{\left.10^{[D 31]}\right)}$ and store it in register (D31, D30).



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $s$ :Source device

D: Operation result

- Range of $\mathrm{S}: \geqslant 0$
- This instruction performs a square root operation on the content in the register designated by S and stores the result in the register designated by D. The square root operation is performed in floating point system.
- If S is an designated constant K or H , the instruction will convert the constant into a binary floating point value before the operation.
- $\quad S$ can only be a positive value. Performing any square root operation on a negative value will result in an "operation error" and this instruction will not be executed.

2. Program Example 1:

- When $\mathrm{MO}=$ On, calculate the square root of the binary floating point (DI, DO) and store the result in register (D11, D10)。


3. Program Example 2:

- When $\mathrm{M} 2=$ On, calculate the square root of $\mathrm{K} 1,234$ (automatically converted into binary floating point) and store the result in register (DII, DIO).

| X 2 |  |  |  |
| :--- | :--- | :--- | :--- |
| H | DESQR | K1234 | D10 |


| $\begin{gathered} \mathrm{ZL} \\ 128 \\ \hline \end{gathered}$ |  | POW |  |  |  |  |  |  | S 1 |  | D |  |  |  |  | Floating Point Power Operation |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DPOW: 13 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  | KnM | Kns | T | C | D | E |  | F |
| S1 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1, Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S1: Device for base. S2: Device for exponent. D: Device for operation result
- This instruction performs power multiplication of binary floating point S1 and S2 and stores the result in D.
$D=P O W[S 1+1, S 1] \wedge[S 2+1, S 2]$
- Only positives are valid for the content in S 1 . Both positives and negatives are valid for the content in S 2 . When designating D registers, the data should be 32-bit and the operation should be performed in floating point system. Therefore, S 1 and $\$ 2$ should be converted into floating point values.

Example: When $\mathrm{SI}^{\mathrm{S} 2}=\mathrm{D}, \mathrm{D}=$ ?
Assume $\mathrm{S} 1=5, \mathrm{~S} 2=3, \mathrm{D}=5^{3}=125$
2. Program Example:

- When MO = On, convert (D1, D0) and (D3, D2) into binary floating points and store them in the 32-bit registers (D11, D10) and (D13, D12).
- When $M 1=$ On, perform POW operation on the binary floting points in 32-bit registers (D11, D10) and (D13, D12) and store the result in the 32-bit register (D21, D20).
- When M2 $=$ On, convert the binary floating point (D21, D20) into decimal floating point (D30 $\times 10[$ [D31]) and store it in register (D31, D30).



1, Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S: Source device D: Converted result
- S occupies 2 consecutive devices. See the specifications of each model for their range of use.
- The binary floating point value of the register designated by $S$ is converted to BIN integer and stored in the register designated by D. The decimal of BIN integer is left out.
- 3. This instruction is the inverse operation of API 49 FLT instruction.

16- bit instruction: $-32,768 \sim 32,767$
32-bit instruction: -2,147,483,648~2,147,483,647

2, Program Example:

- When $\mathrm{XO}=\mathrm{On}$, the binary floating point (D1, D0) will be converted into BIN integer and the result will be stored in (D10). The decimal of BIN integer will be left out.
- When $\mathrm{X} 1=$ On, the binary floating point (D21, D20) will be converted into BIN integer and the result will be stored in (D31, D30). The decimal of BIN integer will be left out.



## 8.4 (ZL 130-139) Trigonometric operation

| $\begin{gathered} \hline \text { ZL } \\ 130 \end{gathered}$ |  | SIN |  |  |  |  |  |  |  | D |  |  |  |  |  | Sine |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DSIN 9 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM |  |  | KnS | T | C | D |  | E | F |
| S |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S:Source value D: SIN result
- $0^{\circ} \leqslant S<360^{\circ}$. See the specifications of each model for their range of use.
- The program will be in radian mode and the RAD value $=$ angle $\times \pi / 180$.
- The SIN value obtained by $S$ is calculated and stored in the register designated by D. The figure below offers the relation between radian and the result.


2, Program Example 1

- When $\mathrm{XO}=\mathrm{On}$, use the RAD value of binary floating point ( $\mathrm{D}, \mathrm{DO}$ ) and obtain its SIN value. The binary floating point result will be stored in (D1 1, D10).


3. Program Example 2

- Input terminals XO and XI select the angle. The angles are converted into RAD value for calculating the SIN value.



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $s$ : Source value
D: COS result
- the program will be in radian mode and the RAD value $=$ angle $\times \pi / 180$.



## 2, Program Example 1:

- When $X 0=$ On, use the RAD value of binary floating point ( $D 1, D 0$ ) and obtain its $C O S$ value. The binary floating point result will be stored in (D11, D10).



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S : Source value

D: TAN result

- The program will be in radian mode and the RAD value $=$ angle $\times \pi / 180$.
- The TAN value obtained by $S$ is calculated and stored in the register designated by D . The figure below offers the relation between radian and the result.



## 2. Program Example 1:

- When $\mathrm{XO}=\mathrm{On}$, use the RAD value of binary floating point (D1, DO) and obtain its TAN value. The binary floating point result will be stored in (D11, D10).



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $s$ : Source value (binary floating point)
D: ASIN result.
- ASIN value=sin - 1. The figure below offers the relation between the entered sin value and the result.

- The decimal floating point of the SIN value designated by $S$ should be within -1.0~+1.0.

2, Program Example:

- When $\mathrm{XO}=$ On, obtain the ASIN value of binary floating point (D1, DO) and store the binary floating point result in (D11, D10).


| $\begin{gathered} \text { ZL } \\ 134 \end{gathered}$ |  | ACOS |  |  |  |  |  | S |  |  | D |  |  |  |  |  | Arc Cosine |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  |  | DACOS: 9 steps <br> 32-bit |
|  | X | Y | M | S | K | H | H | KnX | KnY | KnM |  |  | KnS | T | C | D |  | E | F |
| S |  |  |  |  | * | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $S$ : Source value (binary floating point)
D: ACOS result
- ACOS value $=\cos ^{-1}$. The figure below offers the relation between the entered $\cos$ value and the result.

- The decimal floating point of the $\operatorname{COS}$ value designated by $S$ should be within -1.0~+1.0.

2. Program Example:

- When $\mathrm{XO}=\mathrm{On}$, obtain the ACOS value of binary floating point (DI, DO) and store the binary floating point result in (D11, D10).



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S : Source value (binary floating point)
D: ATAN value
- ATAN value=tan ${ }^{-1}$. The figure below offers the relation between the entered tan value and the result.


2, Program Example:

- When $\mathrm{XO}=$ On, obtain the ATAN value of binary floating point (DI, DO) and store the binary floating point result in (D1 1, D10).



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $s$ : Source value (binary floating point)

D: SINH value.

- $\operatorname{SINH}$ value $=\left(e^{s}-e^{-s}\right) / 2$. The result is stored in $D$.


## 2, Program Example:

- When XO = On, obtain the SINH value of binary floating point (DI, DO) and store the binary floating point result in (D1 1, D10).


| $\begin{gathered} \text { ZL } \\ 137 \end{gathered}$ |  | COSH |  |  |  |  |  |  |  | D |  |  |  |  |  | Hyperbolic Cosine |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DCOSH: 9 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM |  |  | KnS | T | C | D |  | E | F |
| S |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- $s$ : Source value (binary floating point)
D: COSH value
- $\operatorname{COSH}$ value $=\left(e^{s}+e^{-s}\right) / 2$. The result is stored in $D$.


## 2, Program Example:

- When $\mathrm{XO}=\mathrm{On}$, obtain the COSH value of binary floating point (DI, DO) and store the binary floating point result in (D1 1, D10).



1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S : Source value (binary floating point)
D: TANH result
- TANH value $=\left(e^{s}-e^{-s}\right) /\left(e^{s}+e^{-s}\right)$. The result is stored in $D$.


## 2, Program Example:

- When $\mathrm{XO}=$ On, obtain the TANH value of binary floating point (D1, DO) and store the binary floating point result in (D11, D10).



## 8.5 (ZL 140-149) Special function instructions

| $\begin{gathered} \mathrm{ZL} \\ 143 \end{gathered}$ |  | DELAY |  |  |  |  | S |  |  |  |  |  |  |  |  | Delay Instruction |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DELAY: 3 steps <br> 16-bit |  |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |  |
| S |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |  |

1. Explanations:

- To execute the pulse type, add the NP rising edge " $\uparrow$ " command before the command.
- S : delay time (unit: 100 ms ).
- Range of $\mathrm{S}: \mathrm{K1} \sim \mathrm{~K} 1,000$. See the specifications of each model for their range of use.
- After DELAY instruction is executed, the program after DELAY in every scan period will execute delay outputs according to the delay time designated by the user.


## 2, Program Example:

- If XO is turned from Off to On, the external interruption will be generated. DELAY in the interrupt subroutine will be execute for 2 ms before the next step $(\mathrm{XI}=\mathrm{On}$ and $\mathrm{YO}=\mathrm{On}$ ) is executed.


| $\begin{gathered} \hline \text { ZL } \\ 144 \end{gathered}$ |  | GPWM |  |  |  |  | S 1 |  |  |  | D |  |  |  |  | General PWM Output |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | GPWM: 7 steps <br> 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| D |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- s : Width of output pulse $\quad$ S2: Pulse output cycle $\quad$ D: Pulse output device
- $\$ 2$ occupies 3 consecutive devices.
- $\quad S 1 \leqslant S 2$. See the specifications of each model for their range of use.
- Range of $\mathrm{S} 1: \dagger=0 \sim 32,767 \mathrm{~ms}$.
- Range of $\$ 2: t=1 \sim 32,767 \mathrm{~ms}$.
- $\mathrm{S} 2+1$ and $\mathrm{S} 2+2$ are parameters for the system. Do not occupy them.
- Pulse output devices D: Y, M, S.
- When being executed, GPWM instruction designates S1 and S2 and that pulses output will be from device D.
- When $S 1 \leqslant 0$, there will be no pulse output. When $S 1 \geqslant S 2$, the pulse output device will keep being On.
- $\quad \mathrm{S} 1$ and S 2 can be modified when GPWM instruction is being executed.


## 2. Program Example:

- When $\mathrm{XO}=\mathrm{On}, \mathrm{D} 0=\mathrm{K} 1,000, \mathrm{D} 2=\mathrm{K} 2,000$, and Y 10 will output the pulse illustrated below. When $\mathrm{XO}=\mathrm{Off}, \mathrm{Y} 10$ output will be Off.



1. Explanations:

- S: Device for swapping 8 high/low byte.
- If $D$ is used in device $F$, only 16 -bit instruction is applicable.
- As 16 -bit instruction: the contents in the 8 high bytes and 8 low bytes are swapped.
- As 32-bit instruction: the 8 high bytes and 8 low bytes in the two registers swap with each other respectively.
- This instruction adopts pulse execution instructions (SWAPP, DSWAPP).


## 2, Program Example 1:

- When $\mathrm{XO}=\mathrm{On}$, the high 8 bytes and low 8 bytes in DO will swap with each other.



## 3. Program Example 2:

- When $\mathrm{XO}=\mathrm{On}$, the high 8 bytes and low 8 bytes in D 11 will swap with each other and the high 8 bytes and low 8 bytes in D10 will swap with each other.



## 9 Application instructions ZL150~ZL199

## 9.1 (ZL 150-154) Special function instructions

| $\begin{gathered} \text { ZL } \\ 151 \end{gathered}$ |  | PWD |  |  |  |  |  | $S$ D |  |  |  |  |  |  |  |  | Detection of Input Pulse Width |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  |  | PWD: 5 steps 16-bit |
|  | X | Y | M | S | K | H |  | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| n |  |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instructio n.
- S: Source device D: Destination device for storing the detected result
- Range of S: X10~X15
- D must be in the range of D0 to D999., it occupies two consecutive devices.
- PWD instruction is for detecting the interval between the input signals; the valid frequency range is $1 \sim 1 \mathrm{kHz}$. If M1169 = Off, the instruction will continuously detect the intervals between the rising edges of the input signals and the falling edges of the input signals (time unit: 100us). If M1169 = On, the instruction will continuously detect the intervals between rising edges of the input signals (time unit: lus). It cannot designate the same X10~X17 as DCNT and ZRN instructions.
- D occupies two consecutive devices. The longest detection time is $21,474.83647$ seconds, about 357.9139 minutes or 5.9652 hours.
- There is no limitation on the times of using this instruction. However, only one instruction can be executed at a time.

2. Program Example:

- When $\mathrm{XO}=$ On, record the time span of $\mathrm{X10}=$ On and store it in D1 and DO.

| $\begin{gathered} \text { ZL } \\ 154 \end{gathered}$ | RAND |  |  |  |  |  | S1 S2 <br> Word Devices |  |  |  |  |  |  |  |  | Random Number |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F | RAND, RANDP: 7 steps 16-bit DRAND: 13 steps |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * | * | * |  |

1, Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instructio n.
- S1: Lower bound of the random number S2: Upper bound of the random number D: The random number produced
- $\mathrm{S} 1 \leqq S 2 ; K 0 \leqq S 1, S 2 \leqq K 32,767$ If the user enters $\mathrm{S} 1>S 2$, the PLC determines that the operation is wrong and the instruction is not executed.
- Entering S1 > S2 will result in operation error. The instruction will not be executed at this time, M1067, M1068 $=$ On and D1067 records the error code OE1A (hex).

2. Program Example:

- When $\mathrm{X10}=\mathrm{On}$, RAND will produce the random number between the lower bound D 0 and upper bound D10 and store the result in D20.



## 9.2 (ZL 155-159) Positioning control

| $\begin{gathered} \hline \text { ZL } \\ 156 \end{gathered}$ |  | ZRN |  |  |  |  |  |  | S2 | S3 |  | D |  |  |  | Zero Return |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | ZRN: 9 steps DZRN: 17 step | $\begin{aligned} & \text { 16-bit } \\ & \text { 32-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| S3 | * | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- $\quad$ S1:Zero return speed $\quad$ S2: Creep speed
S3: Near p oint signal (DOG)
D: Pulse output device (please use transistor output module)
- When $S 1$ and $S 2$ are used in device $F$, only 16 -bit instruction is applicable.
- $S 1$ specifies the speed at the beginning of home return. The 16 -bit command can be specified in the range of $10 \sim 32,767 \mathrm{~Hz}$, and the 32 -bit command can be specified in the range of $10 \sim 200,000 \mathrm{~Hz}$. When the specified speed is less than $10 \mathrm{~Hz}, 10 \mathrm{~Hz}$ is regarded as the homing speed; when the specified speed is greater than $200 \mathrm{kHz}, 200 \mathrm{kHz}$ is the homing speed.
- $\quad$ S2 specifies the inching speed, after the DOG signal On, specifies the speed of the low-speed part, the specified range is $10 \sim 32,767 \mathrm{~Hz}$.
- $\quad$ S3 Designated DOG signal input (A contact input).
- D pulse output device.
- When performing ZL 158 DRVI relative positioning or ZL 159 DRVA absolute addressing, the PLC has automatically generated forward/reverse pulses to increase and decrease the current value registers as shown in the table below. Therefore, the mechanical position can be grasped at any time, but since the data will disappear when the power of the PLC is turned off, the home position data of the mechanical action must be input when the home position return is performed for the first time.

9 Application instructions ZL150~ZL199

| Group No | PUL | DIR | current <br> number of <br> output <br> pulses (32- <br> bit <br> integer) | Pulse complete | Pulse sending | Emergency <br> stop | Start <br> frequencyk <br> 10-K32767 <br> Default <br> K200 | Accel/De <br> cel <br> timeK10- <br> K10000 <br> defaultK1 <br> 00 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| CH0 (Y0,Y1) | YO | Y1 | D1648 | M1029 | M1344 | M1308 | D1340 | D1343 |
| CH1 (Y2,Y3) | Y2 | Y3 | D1664 | M1030 | M1345 | M1309 | D1352 | D1353 |
| CH2 (Y4,Y5) | Y4 | Y5 | D1680 | M1036 | M1346 | M1310 | D1379 | D1381 |
| CH3 (Y6,Y7) | Y6 | Y7 | D1696 | M1037 | M1347 | M1311 | D1380 | D1382 |
| CH4 (Y10,Y11) | Y10 | Y11 | D1712 | M1102 | M1348 | M1312 | D1400 | D1383 |
| CH5 (Y12,Y13) | Y12 | Y13 | D1728 | M1103 | M1349 | M1313 | D1401 | D1384 |
| CH6 (Y14,Y15) | Y14 | Y15 | D1744 | M1104 | M1350 | M1314 | D1402 | D1385 |
| CH7 (Y16,Y17) | Y16 | Y17 | D1760 | M1105 | M1351 | M1315 | D1403 | D1386 |
| CH8 (Y20,Y21) | Y20 | Y21 | D1776 | M1106 | M1352 |  | D1404 | D1387 |
| CH9 (Y22,Y23) | Y22 | Y23 | D1792 | M1107 | M1353 |  | D1405 | D1388 |
| CH10 (Y24,Y25) | Y24 | Y25 | D1808 | M1108 | M1354 |  | D1406 | D1389 |
| CH11 (Y26,Y27) | Y26 | Y27 | D1824 | M1109 | M1355 |  | D1407 | D1390 |
| CH12 (Y30,Y31) | Y30 | Y31 | D1840 | M1110 | M1356 |  | D1408 | D1391 |
| CH13 (Y32,Y33) | Y32 | Y33 | D1856 | M1111 | M1357 |  | D1409 | D1392 |
| CH14 (Y34,Y35) | Y34 | Y35 | D1872 | M1112 | M1358 |  | D1410 | D1393 |
| CH15 (Y36,Y37) | Y36 | Y37 | D1888 | M1113 | M1359 |  | D1411 | D1394 |
| CH16 (Y40,Y41) | Y40 | Y41 | D1904 | M1114 | M1360 |  | D1412 | D1395 |
| CH17 (Y42,Y43) | Y42 | Y43 | D1920 | M1115 | M1361 |  | D1413 | D1396 |
| CH18 (Y44,Y45) | Y44 | Y45 | D1472 | M1116 | M1362 |  | D1414 | D1397 |
| CH19 (Y46,Y47) | Y46 | Y47 | D1488 | M1117 | M1363 |  | D1415 | D1398 |
| CH20 (Y50,Y51) | Y50 | Y51 | D1504 | M1118 | M1364 |  | D1416 | D1399 |
| CH21 (Y52,Y53) | Y52 | Y53 | D1520 | M1119 | M1365 |  | D1417 | D1420 |
| CH22 (Y54,Y55) | Y54 | Y55 | D1536 | M1205 | M1366 |  | D1418 | D1421 |
| CH23 (Y56,Y57) | Y56 | Y57 | D1552 | M1206 | M1367 |  | D1419 | D1422 |
| Remarks |  |  | D1648: <br> Low word of the current number of output pulses from CH0. D1649: <br> High word of the current number of output pulses from CH0. | After CH0- <br> CH23 pulse <br> output is <br> completed, the <br> corresponding <br> flag bit is ON | Only when the pulse is being sent, the flag bit corresponding to CH0CH 23 is ON | Off->On: <br> The high- <br> speed pulse <br> output <br> pauses <br> immediately. <br> On->Off: <br> Continuing to <br> output the <br> pulses which <br> have not been <br> output |  |  |

2. Program Example:

- When M10=On, start the home return action from YO output pulse at 20 kHz frequency. When it touches the DOG signal $\mathrm{X} 2=\mathrm{On}$, it will run in the opposite direction at 1 kHz frequency of inching speed, output pulse from YO to X2 $=$ Off and stop.

- The zero return operation:

1) When ZRN instruction is executed, set the frequency of the first acceleration segment as the start frequency. The acceleration time of special $D$ is used for reference. $S 1$ will start to move when the acceleration reaches the zero return speed.
2) When the DOG signal goes from Off to On, the zero return speed will decelerate to $S 2$ in the acceleration/deceleration time.
3) When the DOG signal goes from On to Off, the pulse output will immediately stop, 0 will be written in the present value.
4) When the pulse output is completed, the completion flag is ON and the in operation flag is OFF.
5) The ZRN (DZRN) instruction cannot search for the position of the near-point signal (DOG), and the homing operation can only be performed in one direction. The content of the current value register of the pulse amount corresponding to each channel in the home return will change towards the decreasing direction.

6) When the conditions for the start of the return to origin command are met, $\mathrm{CHO}(\mathrm{CH} 1)$ will read the value set by D1343 (D1353) as the acceleration and deceleration time. After accelerating to the origin
return speed, wait for the DOG origin signal to enter and then decelerate from the origin return speed to inches Moving speed until the DOG origin signal is OFF and immediately stop outputting pulses.

1. Explanations:

- S: Pulse output frequency DI: Pulse output device (please use transistor output module)

D2: Output device for the signal of rotation direction

- See remarks for the setting range of S, D1 and D2.
- 3. S is the designated pulse output frequency. The 16 -bit instruction can designate its range $0 \sim+32,767 \mathrm{~Hz}$, $0 \sim-32,768 \mathrm{~Hz}$. the ranges designated by $32-$ bit instruction are $0 \sim+200,000 \mathrm{~Hz}$ and $0 \sim-200,000 \mathrm{~Hz}$. " $+/-$ " signs indicate forward/backward directions. During the pulse output, the frequency can be changed, but not the frequencies of different directions.
- D1 is the pulse output device:

Y0.Y2.Y4.Y6.Y10.Y12.Y14.Y16.Y18.Y20.Y22.Y24.Y26.Y30.Y32.Y34.Y36.Y40.Y42.Y44.Y46.

- The operation of D2 corresponds to the " + " or "- " of $S$. When $S$ is " + ", D2 will be On; when $S$ is "-". D2 will be Off.
- Anyway, when M1207 is OFF, the PLSV instruction does not set acceleration or deceleration, so it doesn't perform acceleration and deceleration at the beginning and stop.
- Anyway, when M1207 is ON, PLSV instruction takes acceleration and deceleration Settings, so the acceleration and deceleration actions start and stop are performed. The acceleration and deceleration time is set according to JC156 ZRN instruction.
- Anyway, when M1207 is OFF, PLSV instructions execute pulse output, and stop directly without decelerating if the driving condition changes to OFF.
- Virtual gateway When M1207 is ON, the PLSV instruction executes pulse output, decelerating and stopping if the driving condition changes to Off.

2, Program Example:

- When M10 = On, YO will output pulses at $20 \mathrm{kHz}, \mathrm{Yl}=$ On indicates forward pulses.


| Group No | PUL | DIR | current number of <br> output pulses (32-bit <br> integer) | Pulse <br> complete | Pulse sending | Accel/Decel <br> timeK10-K10000 <br> defaultK100 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathrm{CH}(\mathrm{YO}, \mathrm{Y} 1)$ | Y 0 | Y 1 | D 1648 | M 1029 | M 1344 | D 1343 |
| $\mathrm{CH} 1(\mathrm{Y} 2, \mathrm{Y} 3)$ | Y 2 | Y 3 | D 1664 | M 1030 | M 1345 | D 1353 |
| $\mathrm{CH} 2(\mathrm{Y} 4, \mathrm{Y} 5)$ | Y 4 | Y 5 | D 1680 | M 1036 | M 1346 | D 1381 |
| $\mathrm{CH} 3(\mathrm{Y} 6, \mathrm{Y} 7)$ | Y 6 | Y 7 | D 1696 | M 1037 | M 1347 | D 1382 |



1. Explanations:

- S1: Number of output pulses (relative designation)

S2: Pulse output frequency
DI: Pulse output device (please use transistor output module)
D2: Output device for the signal of rotation direction

- See remarks for the setting range of S1, S2, D1 and D2.
- $S 1$ is the number of output pulses (relative designation). The 16 -bit instruction can designate the range $32,768 \sim+32,767$. The range designated by 32 -bit instruction is $-2,147,483,648 \sim+2,147,483,647$. If the value in $S 1$ is 0 , that means no output and no action.
- $S 2$ is the designated pulse output frequency. The 16 -bit instruction can designate its range $10 \sim 32,767 \mathrm{~Hz}$. The range designated by 32 -bit instruction is $10 \sim 200,000 \mathrm{~Hz}$.
- The operation of D 2 corresponds to the "+" or "- " of S 1 . When S 1 is "-", D 2 will be Off ; when S 1 is "+ ", D2 will be On. D2 will not be Off immediately after the pulse output is over; it will be Off only when the drive contact of the instruction turns Off.
- Specify the number of pulse output S 1 will become the current value register of $\mathrm{CHO}(\mathrm{YO}, \mathrm{Y})$ pulse (D1648 high bit, D1649 low bit) 32-bit data, CH1 (Y2, Y3) pulse current value register (D1664 high bit, D1 665 low bit) 32 bit Data, and so on. In the reverse direction, the content of the current value register will decrease.
- When DRVI instruction is executing pulse output, you cannot change the content of all operands. The changes will be valid next time when DRVI instruction is enabled.
- When the driving condition of the DRVI command becomes Off, even if the $\mathrm{CHO}(\mathrm{CHI})$ pulse sending indicator M1344 (M1345) is On, the DRVI command cannot be driven again.
- DRVI and DDRVI commands output at 200kHz when the absolute value of the input frequency>200kHz,
and output at 10 Hz when the absolute value of the input frequency<10Hz.
- D1343 (D1353) is the acceleration and deceleration time setting of $\mathrm{CHO}(\mathrm{CH} 1)$ for the first stage of acceleration and the last stage of deceleration. The acceleration and deceleration time is $1 \sim 10,000 \mathrm{~ms}$. If it is higher than $10,000 \mathrm{~ms}$, the factory default value is 100 ms .
- D1340 (D1352) is the $\mathrm{CHO}(\mathrm{CH1})$ start/stop frequency setting. If the pulse output frequency specified by S 2 is less than or equal to the start/stop frequency, the start/stop frequency will be used as the pulse output frequency.
- Please refer to the table for host pulse output channels:

| Group No | PUL | DIR | current number of output pulses (32bit integer) | Pulse complete | Pulse sending | Emergency stop | Start <br> frequencyk <br> 10-K32767 <br> Default <br> K200 | Accel/De <br> cel <br> timeK10- <br> K10000 <br> defaultK1 <br> 00 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| CH0 (Y0,Y1) | YO | Y1 | D1648 | M1029 | M1344 | M1308 | D1340 | D1343 |
| CH1 (Y2,Y3) | Y2 | Y3 | D1664 | M1030 | M1345 | M1309 | D1352 | D1353 |
| CH2 (Y4,Y5) | Y4 | Y5 | D1680 | M1036 | M1346 | M1310 | D1379 | D1381 |
| CH3 (Y6,Y7) | Y6 | Y7 | D1696 | M1037 | M1347 | M1311 | D1380 | D1382 |
| CH4 (Y10,Y11) | Y10 | Y11 | D1712 | M1102 | M1348 | M1312 | D1400 | D1383 |
| CH5 (Y12,Y13) | Y12 | Y13 | D1728 | M1103 | M1349 | M1313 | D1401 | D1384 |
| CH6 (Y14,Y15) | Y14 | Y15 | D1744 | M1104 | M1350 | M1314 | D1402 | D1385 |
| CH7 (Y16,Y17) | Y16 | Y17 | D1760 | M1105 | M1351 | M1315 | D1403 | D1386 |
| CH8 (Y20,Y21) | Y20 | Y21 | D1776 | M1106 | M1352 |  | D1404 | D1387 |
| CH9 (Y22,Y23) | Y22 | Y23 | D1792 | M1107 | M1353 |  | D1405 | D1388 |
| CH10 (Y24,Y25) | Y24 | Y25 | D1808 | M1108 | M1354 |  | D1406 | D1389 |
| CH11 (Y26,Y27) | Y26 | Y27 | D1824 | M1109 | M1355 |  | D1407 | D1390 |
| CH12 (Y30,Y31) | Y30 | Y31 | D1840 | M1110 | M1356 |  | D1408 | D1391 |
| CH13 (Y32,Y33) | Y32 | Y33 | D1856 | M1111 | M1357 |  | D1409 | D1392 |
| CH14 (Y34,Y35) | Y34 | Y35 | D1872 | M1112 | M1358 |  | D1410 | D1393 |
| CH15 (Y36,Y37) | Y36 | Y37 | D1888 | M1113 | M1359 |  | D1411 | D1394 |
| CH16 (Y40, Y41) | Y40 | Y41 | D1904 | M1114 | M1360 |  | D1412 | D1395 |
| CH17 (Y42, Y43) | Y42 | Y43 | D1920 | M1115 | M1361 |  | D1413 | D1396 |
| CH18 (Y44,Y45) | Y44 | Y45 | D1472 | M1116 | M1362 |  | D1414 | D1397 |
| CH19 (Y46,Y47) | Y46 | Y47 | D1488 | M1117 | M1363 |  | D1415 | D1398 |
| CH20 (Y50,Y51) | Y50 | Y51 | D1504 | M1118 | M1364 |  | D1416 | D1399 |
| CH21 (Y52,Y53) | Y52 | Y53 | D1520 | M1119 | M1365 |  | D1417 | D1420 |
| CH22 (Y54,Y55) | Y54 | Y55 | D1536 | M1205 | M1366 |  | D1418 | D1421 |
| CH23 (Y56,Y57) | Y56 | Y57 | D1552 | M1206 | M1367 |  | D1419 | D1422 |



## 2. Program Example:

- When M10= On, YO will output 20,000 pulses (relative designation) at $2 \mathrm{kHz} . \mathrm{Yl}=$ On indicates the pulses are executed in forward direction.


Remarks:

1) Relative position control: Designating the traveling distance starting from the current position by " + / " signs; also known as a relative driving method.

2) Settings of relative positioning and the acceleration/deceleration speed:


| $\begin{gathered} \text { ZL } \\ 159 \end{gathered}$ |  | DRVA |  |  |  |  |  |  | S2 | D2 |  |  |  |  |  | Drive to Absolute |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DRVA: 9 steps DDRVA: 17 steps |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| D1 |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| D2 |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S1: Number of output pulses (absolute designation)

S2: Pulse output frequency
D1: Pulse output device (please use transistor output module) D2: Output device for the signal of rotation direction

- $\quad S 1$ is the number of output pulses (absolute designation). The 16 -bit instruction can designate the range $32,768 \sim+32,767$. The range designated by 32 -bit instruction is $-2,147,483,648 \sim+2,147,483,647$. If the absolute position and the current position in $S 1$ are the same, which means the relative output pulse is 0 . Then to execute this instruciton will NOT output any pulse but the special $M$ flag will be $O N$, indicating the output is complete.
- $\quad S 2$ is the designated pulse output frequency. The 16 -bit instruction can designate its range $10 \sim 32,767 \mathrm{~Hz}$. The range designated by 32-bit instruction is $10 \sim 200,000 \mathrm{~Hz}$.
- D2 The output device of the rotation direction signal. When S 1 is greater than the current relative position, D2: Off. When S1 is less than the current relative position, D2: On, D2 will not be Off immediately after the pulse output ends, you must wait for the command to execute when the contact switch is Off D2: Off.
- Specify the number of pulse output S 1 will become the current value register of $\mathrm{CHO}(\mathrm{YO}, \mathrm{Y} 1$ ) pulse (D1648 high bit, D1649 low bit) 32-bit data, CH1 (Y2, Y3) pulse current value register (D1664 high bit, D1665 low bit) 32 bit Data, and so on. In the reverse direction, the content of the current value register will decrease.
- When DRVA instruction is executing pulse output, you cannot change the content of all operands. The changes will be valid next time when DRVA instruction is enabled.
- When the driving condition of the DRVA command becomes Off, even if the $\mathrm{CHO}(\mathrm{CHI})$ pulse sending indicator M1344 (M1345) is On, the DRVA command cannot be driven again.
- DRVI and DDRVI commands output at 200 kHz when the absolute value of the input frequency>200kHz,
and output at 10 Hz when the absolute value of the input frequency<10Hz.
- D1343 (D1353) is the acceleration and deceleration time setting of $\mathrm{CHO}(\mathrm{CH} 1)$ for the first stage of acceleration and the last stage of deceleration. The acceleration and deceleration time is $1 \sim 10,000 \mathrm{~ms}$. If it is higher than $10,000 \mathrm{~ms}$, the factory default value is 100 ms .
- D1340 (D1352) is the $\mathrm{CHO}(\mathrm{CH1})$ start/stop frequency setting. If the pulse output frequency specified by S 2 is less than or equal to the start/stop frequency, the start/stop frequency will be used as the pulse output frequency.
- Please refer to the table for host pulse output channels:

| Group No | PUL | DIR | current number of output pulses (32bit integer) | Pulse complete | Pulse sending | Emergency stop | Start <br> frequencyk <br> 10-K32767 <br> Default <br> K200 | Accel/De <br> cel <br> timeK10- <br> K10000 <br> defaultK1 <br> 00 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| CH0 (Y0,Y1) | YO | Y1 | D1648 | M1029 | M1344 | M1308 | D1340 | D1343 |
| CH1 (Y2,Y3) | Y2 | Y3 | D1664 | M1030 | M1345 | M1309 | D1352 | D1353 |
| CH2 (Y4,Y5) | Y4 | Y5 | D1680 | M1036 | M1346 | M1310 | D1379 | D1381 |
| CH3 (Y6,Y7) | Y6 | Y7 | D1696 | M1037 | M1347 | M1311 | D1380 | D1382 |
| CH4 (Y10,Y11) | Y10 | Y11 | D1712 | M1102 | M1348 | M1312 | D1400 | D1383 |
| CH5 (Y12,Y13) | Y12 | Y13 | D1728 | M1103 | M1349 | M1313 | D1401 | D1384 |
| CH6 (Y14,Y15) | Y14 | Y15 | D1744 | M1104 | M1350 | M1314 | D1402 | D1385 |
| CH7 (Y16,Y17) | Y16 | Y17 | D1760 | M1105 | M1351 | M1315 | D1403 | D1386 |
| CH8 (Y20,Y21) | Y20 | Y21 | D1776 | M1106 | M1352 |  | D1404 | D1387 |
| CH9 (Y22,Y23) | Y22 | Y23 | D1792 | M1107 | M1353 |  | D1405 | D1388 |
| CH10 (Y24,Y25) | Y24 | Y25 | D1808 | M1108 | M1354 |  | D1406 | D1389 |
| CH11 (Y26,Y27) | Y26 | Y27 | D1824 | M1109 | M1355 |  | D1407 | D1390 |
| CH12 (Y30,Y31) | Y30 | Y31 | D1840 | M1110 | M1356 |  | D1408 | D1391 |
| CH13 (Y32,Y33) | Y32 | Y33 | D1856 | M1111 | M1357 |  | D1409 | D1392 |
| CH14 (Y34,Y35) | Y34 | Y35 | D1872 | M1112 | M1358 |  | D1410 | D1393 |
| CH15 (Y36,Y37) | Y36 | Y37 | D1888 | M1113 | M1359 |  | D1411 | D1394 |
| CH16 (Y40, Y41) | Y40 | Y41 | D1904 | M1114 | M1360 |  | D1412 | D1395 |
| CH17 (Y42, Y43) | Y42 | Y43 | D1920 | M1115 | M1361 |  | D1413 | D1396 |
| CH18 (Y44,Y45) | Y44 | Y45 | D1472 | M1116 | M1362 |  | D1414 | D1397 |
| CH19 (Y46,Y47) | Y46 | Y47 | D1488 | M1117 | M1363 |  | D1415 | D1398 |
| CH20 (Y50,Y51) | Y50 | Y51 | D1504 | M1118 | M1364 |  | D1416 | D1399 |
| CH21 (Y52,Y53) | Y52 | Y53 | D1520 | M1119 | M1365 |  | D1417 | D1420 |
| CH22 (Y54,Y55) | Y54 | Y55 | D1536 | M1205 | M1366 |  | D1418 | D1421 |
| CH23 (Y56,Y57) | Y56 | Y57 | D1552 | M1206 | M1367 |  | D1419 | D1422 |



## 2, Program Example:

- When M10= On, Y0 will output 20,000 pulses (absolute designation) at 2 kHz . $\mathrm{Y} 1=$ On indicates the pulses are executed in forward direction.


Remarks:

1) Absolute position control: Designating the traveling distance starting from the zero point (0); also known as a absolute driving method.


Target position
2) Settings of absolute positioning and the acceleration/deceleration speed:


## 9.3 (ZL 160-169) Perpetual calendar



1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- Range of S1: K0 ~ K23; range of S2 and S3: K0 ~K59.
- $S$ will occupy 3 consecutive devices; D will occupy 3 consecutive points.
- S1,S2 and S3 are compared with the present values of "hour" , "minute" and "second" starting from $S$. The comparison result is stored in $D$.
- $S$ is the "hour" of the current time (KO ~K23) in RTC; $S+1$ is the "minute" (KO ~K59) and $S+2$ is the "second" (K0~K59).
- $S$ is read by TRD instruction and the comparison is started by TCMP instruction. If $S$ exceeds the range, the program will regard this as an operation error and the instruction will not be executed.

2. Program Example:

- When $\times 10=$ On, the instruction will compare the current time in RTC (D20 ~ D22) with the set value 12:20:45 and display the result in M10 ~M12. When X10 goes from On to Off, the instruction will not be executed, but the On/Off stauts prior to M10~M12 will remain.
- Connect M10~M12 in series or in parallel to obtain the result of $\geqq, \leqq$, and $\neq$.

| TCMP | K12 | K20 | K45 | D20 | M10 |
| :--- | :--- | :--- | :--- | :--- | :---: |



1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1: Lower bound of the time for comparison

S2: Upper bound of the time for comparison
S: Current time of RTC
D: Comparison result

- $\quad$ S1, S2, and $S$ will occupy 3 consecutive devices.
- The content in S 1 must be less than the content in S 2 .
- D will occupy 3 consecutive points.
- $\quad \mathrm{S}$ is compared with S 1 and S 2 . The comparsion result is stored in D.
- $\mathrm{S} 1, \mathrm{~S} 1+1, \mathrm{~S} 1+2$ : The "hour", "minute" and "second" of the lower bound of the time for comparison.
- $\$ 2, S 2+1, S 2+2$ : The "hour" , "minute" and "second" ond" of the upper bound of the time for comparison.
- $S, S+1, S+2$ : The "hour" , "minute" and "second" of the current time of RTC.
- DO designated by S is read by TRD instruction and the comparison is started by TZCP instruction. If $\mathrm{S} 1, \mathrm{~S} 2$, and $S$ exceed their ranges, the program will regard this as an operation error and the instruction will not be executed.
- When $S<S 1$ and $S<S 2$, D will be On. When $S>S 1$ and $S>S 2, D+2$ will be On. In other occasions, $D+1$ will be On.


## 2. Program Example:

- When X10= On, TZCP instruction will be executed and one of M10~M12 will be On. When X10 $=$ Off, TZCP instruction will not be executed and the status of M10 ~ M12 prior to X10 = Off will remain unchanged.



1, Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1, S2, and D will occupy 3 consecutive devices.
- $\mathrm{S} 1+\mathrm{S} 2=\mathrm{D}$. The hour, minute, and second of the RTC designated in S 1 plus the hour, minute, and second designated in S 2 . The result is stored in the hour, minute, and second of the register designated in D .
- If $S 1$ and $S 2$ exceed their ranges, the program will regard this as an operation error and the instruction will not be executed.
- If the sum is larger than 24 hours, the value in D will be the result of "sum minuses 24 hours" .


## 2, Program Example:

- When $\times 10=$ On, TADD instruction will be executed and the hour, minute and second in RTC designated in D0 ~ D2 will plus the hour, minute and second in RTC designated in D10 ~ D12. The sum is stored in the hour, minute and second of the register designated in D20 ~ D22.


9 Application instructions ZL150~ZL199


1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1:Time minuend S2: Time subtrahend D: Time remainder
- $\mathrm{S} 1, \mathrm{~S} 2$, and $D$ will occupy 3 consecutive devices.
- $\mathrm{S} 1-\mathrm{S} 2=\mathrm{D}$. The hour, minute, and second of the RTC designated in S 1 minus the hour, minute, and second designated in S 2 . The result is stored in the hour, minute, and second of the register designated in D.
- If S 1 and S 2 exceed their ranges, the program will regard this as an operation error and the instruction will not be executed.
- If the remainder is a negative value, the value in D will be the result of "the negative value pluses 24 hours" .

2. Program Example:

- When $\times 10=$ On, TADD instruction will be executed and the hour, minute and second in RTC designated in D0, ~ D2 will minus the hour, minute and second in RTC designated in D10~D12. The remainder is stored in the hour, minute and second of the register designated in D20 ~ D22.




## 1, Instructions

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- A device that stores a perpetual calendar when the time is read out.
- According to a clock, seven data sets -- year, week, month, day, hour, minute, second -- are stored in D1319 to D1313, according to the TRD instruction, which lets programmers read the time directly into a specified set of seven registers.
- Anyway, D1319 reads only the right two bits of THE AD year, according to the supplementary instructions for reading all four bits.

2, Sample application

- According to the system, when $\mathrm{XO}=\mathrm{On}$, the clock reads the time of the calendar into the specified REGISTERS DO~D6.
- Buy a ticket for D1318, using 1 for Monday, 2 for Tuesday, and so on, and 7 for Sunday.


## Description:

1) Mark and special register of perpetual calendar clock:

| number | The name says | Action function |
| :--- | :--- | :--- |
| M1016 | Perpetual calendar <br> AD year display | Off when D1319 shows AD 2 to the right <br> On D1319 shows the year AD 2 digits to the right plus 2000 |
| M1017 | Plus or minus 30 <br> seconds | Off $\rightarrow$ On for correction when triggered. <br> ( $\sim \sim 29$ seconds return 0, 30~59 seconds, minute plus 1, second return <br> 0). |
| M1076 | The calendar is faulty | Set value ON when out of set range (this check will only be done <br> when starting up) |


| number | The name says | Action function |
| :--- | :--- | :--- |
| D1313 | second | $0 \sim 59$ |
| D1314 | points | $0 \sim 59$ |
| D1315 | when | $0 \sim 23$ |
| D1316 | day | $1 \sim 31$ |
| D1317 | month | $1 \sim 12$ |
| D1318 | week | $1 \sim 7$ |
| D1319 | years | $0 \sim 99$ (2nd from the right of AD) |

- A method of correcting a perpetual calendar clock
- built-in perpetual calendar clock, its correction method is correction time special instruction, please refer to TWR instruction (ZL 167) for details.
- Lent The YEAR displays a 4-digit number:

1) The year normally displays only 2 digits (for example: 2003 displays only 03). If you want to display 4 digits, please type the following program at the beginning of the program.
2) The AD year display is changed from 2 digits to 4 digits, showing the year of THE right 2 digits plus 2000 .
3) If you want to write the new setting time in the mode of 4-digit display in THE YEAR of AD, only 2-digit can be written, and the valid value of this 2-digit is "0~99", which reflects the year of AD is "2000~2099", the relationship between them is as follows. Example: 00= year $200003=$ year $200350=$ year $205099=$ year 2099.

| $\begin{gathered} \hline \text { JC } \\ 167 \end{gathered}$ |  | TWR |  |  |  |  | S |  |  |  |  |  |  |  |  | Write calendar data |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | A 16-bit instructions <br> TRD continuous execution <br> 32-bit instruction <br> There is no |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| D |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |
| - Note that the D operand uses seven devices in a row Please refer to the functional specification table for the use range of each device |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

## 1, Instructions

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- Anyway, S : The device that stores new values to be written to the calendar.
- To adjust a calendar clock, you use this command to write the correct current time into the built-in calendar clock, anyway.
- According to the scheme, when the command is executed, the new set time is written into the PLC's internal calendar clock immediately, so when running the command, pay attention to whether the new set time is written to the current time at the time of writing.
- Anyway, if the value of $S$ content is out of the range, the operation is regarded as an error and the command is not executed.

2, Sample application

- Buy a way to write the correct current time into an implicit calendar clock when $\mathrm{XO}=\mathrm{On}$.


9 Application instructions ZL150~ZL199

|  | General D | project | content | $\rightarrow$ | Te D | project |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| New setting time | D20 | years | 00~99 |  | D1319 | years | Perpetual calendar clock |
|  | D21 | week | 1~7 |  | D1318 | week |  |
|  | D22 | month | 1~12 | $\rightarrow$ | D1317 | month |  |
|  | D23 | day | 1~31 | $\rightarrow$ | D1316 | day |  |
|  | D24 | when | 0~23 | $\rightarrow$ | D1315 | when |  |
|  | D25 | points | 0~59 | $\rightarrow$ | D1314 | points |  |
|  | D26 | second | 0~59 | $\rightarrow$ | D1313 | second |  |

## 3, Example program 2

- Perpetual calendar current time setting, adjust the current time to 15:27:30 on Tuesday, August 19, 2003.
- The content of D0~D6 sets the time for the new perpetual calendar.
- X10=On can replace the current time of the perpetual calendar clock as the set value.

Every time X 11 is On, the perpetual calendar clock will perform a correction action of $\pm 30$ seconds. The so-called correction is that when the second hand of the perpetual calendar clock is between 1 and 29 , it will be automatically classified as "0" seconds and the minute hand will remain unchanged. is automatically reclassified to "0" seconds and the minute hand adds 1 minute.


| $\begin{gathered} \text { ZL } \\ 168 \end{gathered}$ |  | MVM |  |  |  |  |  |  |  | S2 D |  |  |  |  |  | Move the Designated Bit |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | MVM, MVMP: 7 steps <br> DMVM,DMVMP: 13 steps | 16-bit 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  | KnM | KnS | T | C | D | E |  |  | F |
| S1 |  |  |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| D |  |  |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |

1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- $\quad$ s1:Source device 1 S2: Bits to be masked (OFF)

D: Source device 2 / Operation results $\quad[\mathrm{D}=(\mathrm{S} 1 \& 52) \mid(\mathrm{D} \& \sim$ S2)]

- The instruction conducts logical AND operation between S1 and S2 first, logical AND operation between $D$ and $\sim \$ 2$ secondly, and combines the $1^{15 t}$ and $2^{\text {nd }}$ results in $D$ by logical OR operation.
- Rule of Logical AND operation: 0 AND $1=0,1$ AND $0=0,0$ AND $0=0,1$ AND $1=1$.
- Rule of Logical OR operation: 0 OR $1=1,1 O R 0=1,0 O R 0=0,1 O R 1=1$.

2, Program Example 1:

- When XO = ON, MVM instruction conducts logical AND operation between 16-bit register DO and H' FFOO first, logical AND operation between D4 and H' OOFF secondly, and combines the 1st and 2nd results in D4 by logical OR operation.

$$
\begin{array}{|c|c|c|c|c|}
\hline \mathrm{X0} \\
\hline \mathrm{H} & \begin{array}{|l|l|}
\hline \text { MVM } & \text { D0 } \\
\hline
\end{array} & \text { HFF00 } & \text { D4 } \\
\hline
\end{array}
$$


3. Program Example 2:

- Simplify instructions:

| X0 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| -1- | WAND | HFFO0 | D110 | D100 |
|  | WAND | H00FF | D120 | D0 |
|  | WOR | D0 | D100 | D100 |


| X0 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
|  | MVM | D110 | HFF00 | D100 |


| $\begin{gathered} \hline \text { ZL } \\ 169 \end{gathered}$ | D | HOUR |  |  |  |  |  |  | D1 D2 |  |  |  |  |  |  |  | Hour Meter |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  |  | HOUR: 7 steps <br> DHOUR: 13 steps | $\begin{aligned} & \text { 16-bit } \\ & 32 \text {-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C |  |  | E | F |  |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * |  |  |  | * |  |  |
| D1 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| D2 |  | * | * | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- $S$ : Period of time when D2 is On (in hour)

D1: Current value being measured (in hour)
D2: Output device

- If $S$ is used in device $F$, only 16 -bit instruction is applicable.
- D1 will occupy 2 consecutive points. D1 +1 uses 16 -bit register in 16 -bit or 32 -bit instruction.
- Range of S: K1~K32,767 (unit: hour); range of DI: K0~K32,767 (unit: hour). DI + 1 refers to the current time that is less than an hour (range: K0 ~K3,599; unit: second).
- This instruction times the time and when the time reaches the set time (in hour), D2 will be On. This function allows the user to time the operation of the machine or conduct maintenance works.
- After D2 is On, the timer will resume the timing.
- In the 16-bit instruction, when the current time measured reaches the maximum 32,767 hours $/ 3,599$ seconds, the timing will stop. To restart the timing, Dl and $\mathrm{D1}+1$ have to be reset to " 0 " .
- n the 32-bit instruction, when the current time measured reaches the maximum 2,147,483,647 hours $/ 3,599$ seconds, the timing will stop. To restart the timing, D1 ~ D1 + 2 have to be reset to "0" .


## 2. Program Example 1:

- In 16 -bit instruction, when $\mathrm{XO}=\mathrm{On}, \mathrm{Y} 10$ will be On and the timing will start. When the timing reaches 100 hours, YO will be On and DO will record the current time measured (in hour) and DI will record the current time that is less than an hour ( $0 \sim 3,599$; unit: second).


3. Program Example 2:

- In 32-bit instruction, when $X 0=O n, Y 10$ will be On and the timing will start. When the timing reaches 40,000 hours, YO will be On. D1 and D0 will record the current time measured (in hour) and D2 will record the current time that is less than an hour ( $0 \sim 3,599$; unit: second).



## 9.4 (ZL 170-179) Gray code conversion/floating point arithmetic



1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S: Source device for BIN value D: Device for storing Gray code
- If $S$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- This instruction converts the BIN value in the device designated in $S$ into Gray code and stores the value in D.
- See the ranges of $S$ as indicated below. If $S$ exceeds the ranges, the program will regard it as an operation error and the instruction will not be executed. M1067 and M 1068 will be On and D 1067 will record the error code OE1A (hex).

In 16-bit instruction: 0~32,767
In 32-bit instruction: $0 \sim 2,147,483,647$
2. Program Example:

- When $\mathrm{XO}=\mathrm{On}$, the instruction will convert constant $\mathrm{K} 6,513$ into Gray code and store the result in K4Y20.



1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- $\quad$ : Source device for Gray code D: Device for storing BIN value
- If $S$ and $D$ are used in device $F$, only 16 -bit instruction is applicable.
- This instruction converts the Gray code in the device designated in S into BIN value and stores the value in
D.
- This instruction converts the content (in Gray code) in the absolute position encoder connected at the PLC input terminal into BIN value and store the result in the designated register.
- See the ranges of $S$ as indicated below. If $S$ exceeds the ranges, the program will regard it as an operation error and the instruction will not be executed.

In 16-bit instruction: 0~32,767
In 32-bit instruction: $0 \sim 2,147,483,647$
2, Program Example:

- When X20 = On, the Gray code in the absolute position encoder connected at X0 ~ X17 will be converted into BIN value and stored in D10.


K4X0



| $\begin{array}{\|c\|} \hline \mathrm{ZL} \\ 172 \\ \hline \end{array}$ |  | ADDR |  |  |  |  |  | S 1 S2 |  |  |  | D |  |  |  |  | Floating Point Addition |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  |  | DADDR 13 steps32-bit |
|  | X | Y | M | S | K |  | H | KnX | KnY | KnM | KnS |  |  |  | T | C |  | D | E | F |
| S1 |  |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1, Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1: Floating point summand
S2: Floating point addend
D: Sum
- $\quad$ S1 and 52 can be floating point values (FX.XX).
- In DADDR instruction, floating point values (e.g. F1.2) can be entered directly into S1 and S2 or stored in register $D$ for operation. When the instruction is being executed, operand $D$ will store the operation result.
- When S1 and S2 stores the floating point values in register D, their functions are the same as API 120 EADD.
- S1 and S2 can designate the same register. In this case, if the "continuous execution" type instruction is in use and during the On period of the drive contact, the register will be added once in every scan by a "pulse execution" type instruction (DADDRP).

2, Program Example 1:

- When $\mathrm{XO}=$ On, the floating point F1. 20000004768372 will plus F 2.20000004768372 and the result F3. 40000009536743 will be stored in the data registers (D10, D11).

| XO | DADDR | F1.20000004768372 | F2.20000004768372 | D10 |
| :--- | :--- | :--- | :--- | :--- |

## 3. Program Example 2:

- When $\mathrm{XO}=\mathrm{On}$, the floating point value ( $\mathrm{D} 1, \mathrm{D} 0$ ) + floating point value ( $\mathrm{D} 3, \mathrm{D} 2$ ) and the result will be stored in the registers designated in (D11, D10).

| X0 | DADDR | D0 | D2 | D10 |
| :---: | :---: | :---: | :---: | :---: |



1, Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1:Floating point minuend S2: Floating point subtrahend D: Remainder
- $S 1$ and $S 2$ can be floating point values (FX.XX).
- In DSUBR instruction, floating point values (e.g. F1.2) can be entered directly into S1 and S2 or stored in register $D$ for operation. When the instruction is being executed, operand $D$ will store the operation result.
- When S1 and S2 stores the floating point values in register D, their functions are the same as API 121 ESUB.
- $S 1$ and $S 2$ can designate the same register. In this case, if the "continuous execution" type instruction is in use and during the On period of the drive contact, the register will be subtracted once in every scan by a "pulse execution" type instruction (DSUBRP).

2, Program Example 1:

- When $\mathrm{XO}=$ On, the floating point F1. 20000004768372 will minus F 2.20000004768372 and the result $\mathrm{F}-1$ will be stored in the data registers (D10, D11).

| X0 | DSUBR | F1.20000004768372 | F2.20000004768372 | D10 |
| :--- | :--- | :--- | :--- | :--- |

- When $X 0=$ On, the floating point value (D1, D0) - floating point value (D3, D2) and the result will be stored in the registers designated in (D11, D10).

| X0 | DSUBR | D0 | D2 | D10 |
| :--- | :--- | :--- | :--- | :--- |



1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- s : Floating point multiplicand
S2: Floating point multiplicator
D: Product
- $S 1$ and $S 2$ can be floating point values (FX.XX).
- In DMULR instruction, floating point values (e.g. F1.2) can be entered directly into S1 and S2 or stored in register $D$ for operation. When the instruction is being executed, operand $D$ will store the operation result.
- When $S 1$ and $S 2$ stores the floating point values in register $D$, their functions are the same as API 122 EMUL.
- $S 1$ and $S 2$ can designate the same register. In this case, if the "continuous execution" type instruction is in use and during the On period of the drive contact, the register will be multiplied once in every scan by a "pulse execution" type instruction (DMULRP).

2, Program Example 1:

- When $X 0=$ On, the floating point F1. 20000004768372 will multiply $F 2.20000004768372$ and the result F2.64000010490417 will be stored in the data registers (D10, D11).

| XO |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| DMULR | F1.20000004768372 | F2.20000004768372 | D10 |

3. Program Example 2:

- When $\mathrm{XI}=\mathrm{On}$, the floating point value (D1, D0) $\times$ floating point value ( $\mathrm{D} 11, \mathrm{D} 10$ ) and the result will be stored in the registers designated in (D21, D20).

| X1 | DMULR | D0 | D10 | D20 |
| :---: | :--- | :--- | :--- | :--- |



1, Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1: Floating point dividend
S2: Floating point divisor
D: Quotient
- $S 1$ and $S 2$ can be floating point values.
- In DDIVR instruction, floating point values (e.g. F1.2) can be entered directly into S1 and S2 or stored in register $D$ for operation. When the instruction is being executed, operand $D$ will store the operation result.
- When S 1 and S 2 stores the floating point values in register D, their functions are the same as API 123 EDIV.
- If $S 2$ is " 0 " , the program will regard it as an operation error and the instruction will not be executed.


## 2, Program Example 1:

- When $\mathrm{XO}=$ On, the floating point F1.20000004768372 will be divided by F2.20000004768372 and the result F0.545454561710358 will be stored in the data registers (D10, D11).

| XO | DDIVR | F1.20000004768372 | F2.20000004768372 | D10 |
| :--- | :--- | :--- | :--- | :--- |

- When $\mathrm{XI}=$ On, the floating point value (D1, D0) $\div$ floating point value ( $\mathrm{D} 11, \mathrm{D} 10$ ) and the quotient will be stored in the registers designated in (D21, D20).

| X 1 | DIVR | D0 | D10 | D20 |
| :--- | :--- | :--- | :--- | :--- |


| $\begin{gathered} \text { ZL } \\ 176 \end{gathered}$ | MMOV |  |  |  |  |  | $S \quad \mathrm{D}$ |  |  |  |  |  |  |  |  | Magnifying Transfer with <br> Sign Extension |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | MMOV: 5 steps <br> 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |

1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S : Data source (16-bit)

D: Data destination (32-bit)

- MMOV instruction sends the data in the 16 -bit $S$ device to the 32 -bit $D$ device. The designated sign bit will be copied and stored in the destination device.


## 2, Program Example 1:

- When $\mathrm{X} 23=$ On, the data in D4 will be sent to D6 and D7.



1, Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- $s$ : Source device
n : Data length to be summed up
D: Device for storing the result
- WSUM instruction sums up $n$ devices starting from $S$ and store the result in $D$.
- If the specified source devices $S$ are out of valid range, only the devices in valid range will be processed.
- Valid range for $\mathrm{n}: 1 \sim 64$. If the specified $n$ value is out of the available range (1~64), PLC will take the upper (64) or lower (1) bound value as the set value.
- D used in the 16 -bit/32-bit instruction is a 32 -bit register.

2, Program example 1:

- When $\times 10=O N, 3$ consecutive devices $(\mathrm{n}=3)$ from D0 will be summed up and the result will be stored in (D11, D10).


3. Program example 2 :

- When $\mathrm{X10}=\mathrm{ON}, 3$ consecutive devices $(\mathrm{n}=3)$ from ( $\mathrm{D} 1, \mathrm{D} 0$ ) will be summed up and the result will be stored in (D1 1, D10).



## 9.5 (ZL 180-190) Matrix processing

| $\begin{gathered} \text { ZL } \\ 180 \end{gathered}$ | MAND |  |  |  |  |  |  |  | S2 D |  |  | n |  |  |  | Matrix 'AND' | Operation |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | MAND: 9 steps <br> 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| S2 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |

1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1:Matrix source device 1 S2: Matrix source device 2

D: Operation result $n$ : Array length

- Range of $\mathrm{n}: \mathrm{Kl} \sim \mathrm{K} 256$.
- S1, and S2 designate KnX, KnY, KnM and KnS; D designates KnYm KnM and KnS.
- esignate $\mathrm{n} \leqq 4$.
- The two matrix sources $\$ 1$ and $\$ 2$ perform matrix 'AND' operation according to the array length n . The result is stored in D.
- Operation rule of matix 'AND' : The result will be 1 if both two bits are 1 ; otherwise the result will be 0 .


## 2. Program Example:

- When $\mathrm{XO}=$ On, the 3 arrays of 16 -bit registers DO $\sim$ D2 and the 3 arrays of 16 -bit registers D10~D12 will perform a matrix 'AND' operation. The result will be stored in the 3 arrays of 16 -bitd registers D20 ~ D22.

| X0 | MAND | D0 | D10 | D20 | K3 |
| :---: | :---: | :---: | :---: | :---: | :---: |



Remarks:

1) A matix consists of more than 1 consecutive 16 -bit registers. The number of registers in the matrix is the length of the array ( n ). A matrix contains $16 \times \mathrm{n}$ bits (points) and there is only 1 bit (point) offered for an operand at a time.
2) The matrix instruction gathers a series of $16 \times n$ bits ( $b 0 \sim b 16 n-1$ ) and designates a single point for operation. The point will not be seen as a value.
3) The matrix instruction processes the moving, copying, comparing and searching of one-to-many or many-to-many matrix status, which is a very handy and important application instruction.
4) The matrix operation will need a 16 -bit register to designate a point among the 16 n points in the matrix for the operation. The register is the Pointer (Pr) of the matrix, designated by the user in the instruction. The vaild range of $\operatorname{Pr}$ is $0 \sim 16 \mathrm{n}-1$, corresponding to b0 ~b16n-1 in the matrix.
5) There are left displacement, right displacement and rotation in a matrix operation. The bit number decreases from left to right (see the figure below).

6) The matrix width $(C)$ is fixed at 16 bits.
7) Pr: matrix pointer. E.g. if $\operatorname{Pr}$ is 15 , the designated point will be bl5.
8) Array length $(R)$ is $n: n=1 \sim 256$.

Example: The matrix is composed of $D 0, n=3 ; D 0=H A A A A, D 1=H 5555, D 2=$ HAAFF

| $\mathrm{C}_{15} \mathrm{C}_{14} \mathrm{C}_{13} \mathrm{C}_{12} \mathrm{C}_{11} \mathrm{C}_{10} \mathrm{C}_{8} \mathrm{C}_{8} \mathrm{C}_{7} \mathrm{C}_{6} \mathrm{C}_{5} \mathrm{C}_{4} \mathrm{C}_{3} \mathrm{C}_{2} \mathrm{C}_{1} \mathrm{C}_{0}$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{R}_{0}$ | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | , | 1 | ${ }^{1}$ | 1 | 0 |  |
| $\mathrm{R}_{1}$ | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | D1 |
| $\mathrm{R}_{2}$ | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | D2 |

Example: The matrix is composed of $\mathrm{K} 2 \mathrm{XO}, \mathrm{n}=3 ; \mathrm{K} 2 \mathrm{XO}=\mathrm{H} 37, \mathrm{~K} 2 \mathrm{X} 10=\mathrm{H} 68, \mathrm{~K} 2 \mathrm{X} 20=\mathrm{H} 45$

| $\mathrm{C}_{15} \mathrm{C}_{14} \mathrm{C}_{13} \mathrm{C}_{12} \mathrm{C}_{11} \mathrm{C}_{10} \mathrm{C}_{8} \mathrm{C}_{8}$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{R}_{0}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 1 |  | $\mathrm{X}_{0} \sim \mathrm{X}_{7}$ |
| $\mathrm{R}_{1}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 0 | $\mathrm{X}_{10} \sim \mathrm{X}_{17}$ |
| $\mathrm{R}_{2}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 1 | $\mathrm{X}_{20} \sim \mathrm{X}_{27}$ |


| $\begin{gathered} \mathrm{ZL} \\ 181 \end{gathered}$ |  | MOR |  |  |  |  | S1 S2 D |  |  |  |  | n |  |  |  | Matrix 'OR' | Operation |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | MOR: 9 steps <br> 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| S2 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |

1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1:Matrix source device 1 S2: Matrix source device 2. D: Operation result $n$ : Array length
- Range of $\mathrm{n}: \mathrm{Kl} \sim \mathrm{K} 256$.
- S1, and S2 designate KnX, KnY, KnM and KnS; D designates KnYm KnM and Kns.
- esignate $\mathrm{n} \leqq 4$.
- The two matrix sources S1 and S2 perform matrix 'OR' operation according to the array length n . The result is stored in $D$.
- Operation rule of matrix 'OR' : The result will be 1 if either of the two bits is 1 . The result is 0 only when both two bits are 0 .

2, Program Example:

- When $\mathrm{XO}=$ On, the 3 arrays of 16 -bit registers D0 $\sim$ D2 and the 3 arrays of 16 -bit registers D10~D12 will perform a matrix 'OR' operation. The result will be stored in the 3 arrays of 16 -bit registers D20 ~ D22.

| X0 | MOR | D0 | D10 | D20 | K3 |
| :--- | :--- | :--- | :--- | :--- | :--- |



| $\begin{gathered} \text { ZL } \\ 182 \end{gathered}$ |  | MXOR |  |  |  |  | S 1 S2 D |  |  |  |  |  |  |  |  | Matrix 'XOR' | Operation |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | MXOR: 9 steps <br> 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| S2 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |

1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1: Matrix source device 1 S2: Matrix source device $2 \quad$ D: Operation result $n$ : Array length
- Range of $\mathrm{n}: \mathrm{Kl} \sim \mathrm{K} 256$.
- S1, and S2 designate KnX, KnY, KnM and KnS; D designates KnYm KnM and KnS .
- The two matrix sources S1 and S2 perform matrix 'XOR' operation according to the array length n . The result is stored in D.
- Operation rule of matrix 'XOR' : The result will be 1 if the two bits are different. The result will be 0 if the two bits are the same.

2. Program Example:

- When $\mathrm{XO}=\mathrm{On}$, the 3 arrays of 16 -bit registers D0 $\sim$ D2 and the 3 arrays of 16 -bit registers D10~D12 will perform a matrix 'XOR' operation. The result will be stored in the 3 arrays of 16-bit registers D20 ~ D22.


| $\begin{gathered} \hline \text { ZL } \\ 183 \end{gathered}$ | MXNR |  |  |  |  |  | S 1 S2 D |  |  |  |  |  |  |  |  | Matrix | Operation |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | MXNR: 9 steps <br> 16-bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| S2 |  |  |  |  |  |  | * | * | * | * | * | * | * |  |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |  |

## 1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S1:Matrix source device 1
S2: Matrix source device 2
D: Operation result
n : Array length
- Range of $\mathrm{n}: \mathrm{Kl} \sim \mathrm{K} 256$.
- S1, and S2 designate KnX, KnY, KnM and KnS; D designates KnYm KnM and Kns.
- The two matrix sources S1 and S2 perform matrix 'XNR' operation according to the array length n . The result is stored in D.
- Operation rule of matrix 'XNR' : The result will be 1 if the two bits are the same. The result will be 0 if the two bits are different.

2. Program Example:

- When $\mathrm{XO}=\mathrm{On}$, the 3 arrays of 16 -bit registers D0 $\sim$ D2 and the 3 arrays of 16 -bit registers D10~D12 will perform a matrix 'XNR' operation. The result will be stored in the 3 arrays of 16-bit registers D20~D22.

| M0 | MXNR | D0 | D10 | D20 |
| :--- | :--- | :--- | :--- | :--- |


| b15 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Before Execution | (S1) |  | 0 |  | 01 | 10 | 11 | 10 | 1 | 0 |  | 0 | 1 |  | 1 |  | 1 |
|  |  |  | 0 | 1 | 01 | 10 | 1 | 10 | 1 | 0 |  | 0 | 1 | 0 | 1 |  | 1 |
|  |  |  | 0 | 1 | 01 | 10 | 1 | 10 | 1 | 0 |  | 0 | 1 | 0 | 1 |  | 1 |
|  |  |  | MXNR |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | (S2) |  |  | 0 | 00 | 01 | 11 | 11 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |  | 1 |
|  |  |  | 0 | 0 | 0 | 01 | 11 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |  | 1 |
|  |  |  | 0 |  |  |  |  |  |  | I |  | 1 | 0 |  | 1 |  | 1 |
| After Execution |  |  | $F$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | (D) |  | 1 | 0 |  | 0 | 01 | 10 |  | 0 |  | 0 | 0 |  | 1 |  | 1 |
|  |  |  | 1 | 0 | 10 | 00 | 01 | 10 |  | 0 |  | 0 | 0 |  | 1 |  | 1 |
|  |  |  | 1 | 0 |  | 0 | 01 | 10 | 1 | 0 |  | 0 | 0 |  | 1 |  | 1 |



1, Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction.
- S: Matrix source device
D: Operation result
n : Array length
- Range of $\mathrm{n}: \mathrm{Kl} \sim \mathrm{K} 256$
- $S$ designates $\mathrm{KnX}, \mathrm{KnY}, \mathrm{KnM}$ and KnS ; D designates $\mathrm{KnY}, \mathrm{KnM}$ and KnS .
- $S$ performs an inverse matrix operation according to the array length $n$. The result is stored in $D$.

2, Program Example:

- When $\mathrm{XO}=\mathrm{On}$, the 3 arrays of 16 -bit registers D0 ~ D2 perform a matrix inverse operation. The result will be stored in the 3 arrays of 16 -bit registers D20 ~ D22.

| X0 | MINV | D0 | D20 | K3 |
| :--- | :--- | :--- | :--- | :--- |



## 9.6 (ZL 191-199) Positioning instructions

| $\begin{gathered} \text { ZL } \\ 192 \end{gathered}$ | D | PPMA |  |  |  |  | S1 S2 <br> Word Devices |  |  |  |  |  |  |  |  | 2-Axis Absolute Point to <br> Point Motion |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  |  |  |  |  |  |  |  |  |  |  |  | DPPMA: 17 steps32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S1: Number of output pulses of $X$ axis $\quad$ S2: Number of output pulses of $Y$ axis
S: Max. point to point output frequency
D: Pulse output device
- In terms of pulse output methods, this instructin only supports "pulse + direction" mode.
- S 1 and S 2 are the designated (absolute designation) number of output pulses in X axis ( YO or Y 4 ) and Y axis ( Y 2 or Y 6 ). The range of the number is $-2,147,483,648 \sim+2,147,483,647$ (+/- represents the forward/backward direction). When in forward direction, the pulse present value registers CHO (D1649 high word, D1648 low word), CH1 (D1665high word, D1664 low word), CH2 (D1681high word, D1680 low word) and CH3 (D1697 high word, D1696 low word) will increase. When in backward direction, the present value will decrease.
- D can designate Y0 and Y4.

When YO is designated:
YO refers to 1 st group X-axis pulse output device.
Y 1 refers to 1 st group X -axis direction signal.
Y2 refers to 1st group Y-axis pulse output device.
Y3 refers to 1st group Y-axis direction signal.
Y4 refers to 2 nd group $X$-axis pulse output device.
Y 5 refers to 2 nd group X -axis direction signal.
Y6 refers to 2 nd group $Y$-axis pulse output device.
Y7 refers to 2 nd group $Y$-axis direction signal.

- When direction signal outputs, Off will not occur immediately after the pulse output is over. Direction signal will turn Off when the drive contact is Off.
- Refer to DDRVI and DDRVA instructions for special $M$ and $D$ corresponding to each channel.
- The time shall be longer than 10 ms . If the time is shorter than 10 ms or longer than $10,000 \mathrm{~ms}$, the output will be operated at 10 ms . Default setting $=100 \mathrm{~ms}$.
- If the maximum output frequency setting is less than 10 Hz , the output will be operated at 10 Hz . If the setting is more than 200 kHz , the output will be operated at 200 kHz .
- When the 2-axis synchronous motion instruction is enabled, the start frequency and acceleration/deceleration time in $Y$ axis will be same as the settings in $X$ axis.
- The number of output pulses for the 2-axis motion shall not be the values within $1 \sim 59$; otherwise the line drawn will not be straight enough.
- There is no limitation on the number of times using the instruction. However, assume CH 1 or CH 2 output is in use, the 1 st group $X / Y$ axis will not be able to output. If CH 3 or CH 4 output is in use, the 2 nd group $\mathrm{X} / \mathrm{Y}$ axis will not be able to output.

2. Program Example: Draw a rhombus as the figure below


- Steps:

1) Set the four coordinate ( $-27,000,-27,000$ ), ( $0,-55,000$ ), ( $27,000,-27,000$ ), ( 0,0 ) (as the figure above).

Place them in the 32-bit (D200, D202), (D204, D206), (D208, D210), (D212, D214).
2) Write program codes as follows.
3) PLC RUN. Set MO as On and start the 2-axis line drawing.


- Motion explanation:

When PLC RUN and M0 $=$ On, PLC will start the first point-to-point motion by 100kHz. D0 will plus 1 whenever a point-to-point motion is completed and the second point-to-point motion will start to execute automatically. The same motion will keep executing until the fourth point-to-point motion is completed.

| $\begin{gathered} \text { ZL } \\ 194 \end{gathered}$ | D | CIMA |  |  |  |  | $S 1 \quad S 2 \quad S$ <br> Word Devices |  |  |  |  |  |  |  |  | 2-Axis Absolute Position Arc Interpolation |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | D | vic |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  | DCIMA: 17 steps |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  | 32-bit |
| S |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

## Mode 0:

- S1: Number of output pulses of $X$ axis $\quad$ S2: Number of output pulses of $Y$ axis
S: Parameter setting
D: Pulse output device
- In terms of pulse output methods, this instructin only supports "pulse + direction" mode.
- $\quad$ S1 and $S 2$ are the designated (absolute designation) number of output pulses in $X$ axis ( $Y 0$ or $Y 4$ ) and $Y$ axis ( Y 2 or Y 6 ). The range of the number is $-2,147,483,648 \sim+2,147,483,647$. The pulse present value register will increase when in the positive direction. In the opposite direction, it will decrease.
- $S$ (direction and resolution setting): Set K0 to output 10 segments clockwise (normal resolution), and set K2 to output 20 segments clockwise (higher resolution), you can draw a $90^{\circ}$ arc as shown in Figure (1), (2); set K 1 to output 10 counterclockwise segments (normal resolution), and set K3 to output 20 counterclockwise segments (higher resolution), you can draw a $90^{\circ}$ arc as shown in (3), (4) shown. $S$ is $K 0$ or $K 1$, which means working in mode 0 , and $S$ is $K 2$ or $K 3$, which means working in mode 1.
- $S+1 \sim S+2$ (walking frequency setting): The general setting range is $10 \mathrm{hz} \sim \mathrm{K} 200000 \mathrm{hz}$.




- D can designate Y0 and Y4.

When YO is designated:
YO refers to 1st group X-axis pulse output device.
Yl refers to 1 st group X -axis direction signal.
Y2 refers to 1st group Y-axis pulse output device.
Y3 refers to 1st group Y-axis direction signal.
When $Y 4$ is designated:
Y4 refers to 2nd group X-axis pulse output device.
Y5 refers to 2nd group X-axis direction signal.
Y6 refers to 2nd group Y-axis pulse output device.
Y7 refers to 2nd group Y-axis direction signal.
When direction signal outputs, Off will not occur immediately after the pulse output is over. Direction signal will turn Off when the drive contact is Off.

- Draw four $90^{\circ}$ arcs.
- When the direction signal is On, the direction is positive. When the direction signal is Off, the direction is negative. When $S$ is set as $K 0, K 2$, the arcs will be clockwise (see figure 5 ). When $S$ is set as $K 1, K 3$, the arcs will be counterclockwise (see figure 6).


Figure 5


Figure 6

- When the 2-axis motion is being executed in 10 segments (of average resolution), the operation time of the instruction when the instruction is first enabled is approximately 5 ms . The number of output pulses cannot be less than 100 and more than $1,000,000$; otherwise, the instruction cannot be enabled.
- When the 2-axis motion is being executed in 20 segments (of high resolution), the operation time of the instruction when the instruction is first enabled is approximately 10 ms . The number of output pulses cannot be less than 1,000 and more than 10,000,000; otherwise, the instruction cannot be enabled.
- If you wish the number of pulses in 10 -segment or 20 -segment motion to be off the range, you may adjust the gear ratio of the servo for obtaining your desired number.
- Every time when the instruction is executed, only one $90^{\circ}$ arc can be drawn. It is not necessary that the arc has to be a precise arc, i.e. the numbers of output pulses in $X$ and $Y$ axes can be different.
- There are no settings of start frequency and acceleration/deceleration time.
- There is no limitation on the number of times using the instruction.
- The settings of motion time in the high 16 bits of $S$ can be slower than the the fastest suggested time but shall not be faster than the fastest suggested time.
- The fastest suggested time for the arc interpolation:


## Mode 1:

- S1~S1+1 represent the center of the X -axis. S2~S2+1 represent the center of the Y -axis. S: parameter setting. D: Pulse output device.
- This command pulse output mode only supports "pulse + direction" mode.
- S (direction and resolution setting): Set KO to output 10 segments clockwise (normal resolution), and set K2 to output 20 segments clockwise (higher resolution), you can draw a $90^{\circ}$ arc as shown in Figure (1) , (2); set K1 to output 10 counterclockwise segments (normal resolution), and set K3 to output 20 counterclockwise segments (higher resolution), you can draw a $90^{\circ}$ arc as shown in (3), ( 4) shown.
- $S$ is $K 0$ or $K 1$, which means working in mode 0 , and $S$ is $K 2$ or $K 3$, which means working in mode 1 .
- $S+1 \sim S+2$ (walking frequency setting): The general setting range is $10 \mathrm{hz} \sim \mathrm{K} 200000 \mathrm{hz}$.
- $S+3 \sim S+4$ indicates that the length of the arc to be executed is in degrees, and the format is a floatingpoint number. For example, F150.23, expressed as 150.23 degrees.
- $S+5 \sim S+6$ represents the length of the arc that has been run (read-only, the format is a floating point number, such as F125.23, which means that it has run 125.23 degrees).
2, Program Example 1: Draw an ellipse as shown below, with mode 0.

- Steps:

1) Set the four coordinates $(0,0),(16000,22000),(32000,0),(16000,-22000)$ (as the figure above). Place them in the 32-bit (D200, D202), (D204, D206), (D208, D210), (D212, D214).
2) Select "draw clockwise arc" and "average resolution" ( $S=D 100=K 0$ ).
3) Select DCIMA instruction for drawing arc and write program codes as follows.
4) PLC RUN. Set MO as On and start the drawing of the ellipse.


- Motion explanation:

When PLC RUN and MO = On, PLC will start the drawing of the first segment of the arc. DO will plus

1 whenever a segment of arc is completed and the second segment of the arc will start to execute automatically. The same motion will keep executing until the fourth segment of arc is completed。
3. Program Example 2: Draw a tilted ellipse as the figure below


- Steps:

1) Find the max. and min. coordinates on $X$ and $Y$ axes (0, 0), (26000, 26000), (34000, 18000), (8000, 8000) (as the figure above). Place them respectively in the 32-bit (D200, D202), (D204, D206), (D208, D210) and (D212, D214).
2) Select "draw clockwise arc" and "average resolution" ( $\mathrm{S}=\mathrm{D} 100=\mathrm{KO}$ ).
3) Select DCIMA instruction for drawing arc and write program codes as follows.
4) PLC RUN. Set MO as On and start the drawing of the ellipse.

| $\mathrm{H}=$ | D0 | K1 | HDCIMA | D200 | D202 | D100 | Y0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| - $=$ | D0 | K2 | HDCIMA | D204 | D206 | D100 | Y0 |
| H= | D0 | K3 | HDCIMA | D208 | D210 | D100 | Y0 |
| - $=$ | D0 | K4 | - DCIMA | D212 | D214 | D100 | Y0 |



- Motion explanation:

When PLC RUN and MO = On, PLC will start the drawing of the first segment of the arc. DO will plus 1 whenever a segment of arc is completed and the second segment of the arc will start to execute automatically. The same motion will keep executing until the fourth segment of arc is completed.

| $\begin{gathered} \text { ZL } \\ 195 \end{gathered}$ | D | PTPO |  |  |  |  | S1 S2 <br> Word Devices |  |  |  |  |  |  |  |  | Single-Axis Pulse Output by <br> Table |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  |  |  |  |  |  |  |  |  |  |  |  | DPTPO: 13 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Explanations:

- S1:Source start device $\quad$ S2: Number of segments D: Pulse output device
- According to the value of $S 2+0$, every segment consecutively occupy four register $D .(S 1+0)$ refers to output frequency. ( $\$ 1+2$ ) refers to the number of output pulses.
- When the output frequency of $S 1$ is less than 1 , PLC will automatically modify it as 1 . When the value is larger than $200,000 \mathrm{kHz}$, PLC will automatically modify it as $200,000 \mathrm{kHz}$.
- $\quad$ S2 +0 : number of segments (range: $1 \sim 60$ ). $\$ 2+1$ : number of segments being executed. Whenever the program scans to this instruction, the instruction will automatically update the segment No. that is currently being executed. D can only designate output devices YO, Y2, Y4 and Y6 and can only perform pulse output control. For the pin for direction control, the user has to compile other programs to control.
- This instruction does not offer acceleration and deceleration functions. Therefore, when the instruction is disabled, the output pulses will stop immediately.
- In every program scan, each channel can only be executed by one instruction. However, there is no limitation on the number of times using this instruction.
- When the instruction is being executed, the user is not allowed to update the frequency or number of the segments. Changes made will not be able to make changes in the actual output.

2, Program Example:

- When $\mathrm{XO}=\mathrm{On}$, the output will be operated according to the set frequency and number of pulses in every segment.
- Format of the table:

| S2=D300, number of segments <br> $(D 300=K 60)$ | $S 1=D 0$, frequency $(S 1+0)$ | $S 1=D 0$, number of output pulses <br> $(S 1+2)$ |
| :--- | :--- | :--- |


| K1 $1^{\text {st }}$ segment $)$ | D1, D0 | D3, D2 |
| :--- | :--- | :--- |
| K2(2 ${ }^{\text {nd }}$ segment $)$ | D5, D4 | D7, D6 |
| $:$ | $:$ | $:$ |
| $:$ | $:$ | $:$ |
| K60(60 |  |  |

- Monitor the segment No. that is currently being executed in register D301.

- The pulse output curve:

- among them: $\dagger 1=(D 3, D 2) \div(D 1, D 0) ; \dagger 2=(D 7, D 6) \div(D 5, D 4) ; \dagger 60=(D 239, D 238) \div(D 237, D 236)$

| $\begin{gathered} \text { ZL } \\ 197 \end{gathered}$ |  | CLLM |  |  |  |  |  |  | S2 | S |  |  |  |  |  | Close Loop Position Control |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DCLLM: 17 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 | * |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1, Explanations:

- S1:Feedback source device

S2: Target number of feedbacks
S3: Target frequency of output
D: Pulse output device
The corresponding interruption of S1:

| Source device | X0 | X1 | X2 | X3 | C241-C254 |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Corresponding <br> outout | Y0 | Y2 | Y4 | Y6 | Y0 | Y2 | Y4 | Y6 |
| Interruption <br> No. | $100 \square$ | $110 \square$ | $120 \square$ | $130 \square$ | 1010 | 1020 | 1030 | 1040 |

note: $\square=1$ : rising-edige trigger, $\square=0$ falling-edge trigger

- 1) When S 1 designates X as the input points and the pulse output reaches the set target number of feedbacks in $\$ 2$, the output will continue to operate by the frequency of the last segment until the interruption of $X$ input points occurs.

2) When $\$ 1$ designates a high speed counter and the pulse output reaches the set target number of feedbacks in $\$ 2$, the output will continue to operate by the frequency of the last segment until the feedback pulses reaches the target number.
3) The range of S2: $-2,147,483,648 \sim+2,147,483,647$ (+/-represents the forward/backward direction). When in forward direction, the pulse present value registers CHO (D1649 high word, D1648 low word), CH1 (D1665 high word, D1664 low word), CH2 (D1681 high word, D1 680 low word) and CH3 (D1697 high word, D1696 low word) will increase. When in backward direction, the present value will decrease..

- If S 3 is lower than 10 Hz , the output will operate at 10 Hz ; if S 3 is higher than 200 kHz , the output will operate at 200 kHz .
- D can only designate $Y 0, Y 2, Y 4$ and $Y 6$ and the direction signals repectively are $Y 1, Y 3, Y 5$ and $Y 7$. When there is a direction signal output, the direction signal will not be Off immediately after the pulse output is completed. The direction signal will be Off only when the drive contact is Off.
- D1340, D1352, D1379 and D1380 are the settings of start/end frequencies of $\mathrm{CHO} \sim \mathrm{CH} 3$. The minimun frequency is 10 Hz and default is 200 Hz .
- D1343, D1353, D1381 and D1382 are the settings of the time of the first segment and the last deceleration segment of $\mathrm{CHO} \sim \mathrm{CH} 3$. The acceleration/deceleration time cannot be shorter than 10 ms . The outptu will be operated in 10 ms if the time set is shorter than 10 ms or longer than $10,000 \mathrm{~ms}$. The dafault setting is 100ms.
- D1131, D1132, D1478 and D1479 are the output/input ratio of the close loop control in CHO ~ CH3. K1 refers to 1 output pulse out of the 100 target feedback input pulses; K200 refers to 200 output pulses out of the 100 target feedback input pulses. D1131, D1132, D1478 and D1479 are the numerators of the ratio (range: K1 ~K10,000) and the denominator is fixed as K100 (the user does not have to enter a denominator).

| Group No | PUL | DIR | current <br> number <br> of output <br> pulses <br> (32-bit <br> integer) | Pulse <br> complete flag | Pulse sending | Emergency <br> stop <br> without <br> slowing <br> down | Start <br> frequencyk10- <br> K32767 <br> defaultK200 | Accel/Decel <br> timeK10- <br> K10000 <br> defaultK100 | deceleration time |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| CH0 (YO,Y1) | YO | Y1 | D1648 | M1029 | M1344 | M1308 | D1340 | D1343 | D1936 |
| CH1 (Y2,Y3) | Y2 | Y3 | D1664 | M1030 | M1345 | M1309 | D1352 | D1353 | D1937 |
| CH2 (Y4,Y5) | Y4 | Y5 | D1680 | M1036 | M1346 | M1310 | D1379 | D1381 | D1938 |
| CH3 (Y6,Y7) | Y6 | Y7 | D1696 | M1037 | M1347 | M1311 | D1380 | D1382 | D1939 |

## 2, Close Loop Explanations:

- Function: Immediately stop the high-speed pulse output according to the number of feedback pulses or external interruption signals.
- The execution:


Number of output pulses $=$
target number of feedbacks $\times$ percentage value/100

- How to adjust the time for the completion of the positioning:

1) The time for the completion of the positioning refers to the time for "acceleration + high speed + deceleration + idling" (see the figure above). For example, you can increase or decrease the entire number of output pulses by making adjustment on the percentage value and further increase or decrease the time required for the positioning.
2) Among the four segments of time, only the idling time cannot be adjusted directly by the user. However, you can determine if the execution result is good or bad by the length of the idling time. In theory, a bit of idling left is the best result for a positioning.
3) Owing to the close loop operation, the length of idling time will not be the same in every execution. Therefore, when the content in the special D for displaying the actial number of output pulses is smaller or larger than the calculated number of output pulses (taget number of feedbacks $x$ percentage value/100), you can improve the situation by adjusting the percentage value, acceleration/decelartion time or target frequency.

## 3. Program Example:

- Assume we adopt X0 as the external interruption, together with 1001 (rising-edge trigger) interruption program; target number of feedbacks $=50,000$; target frequency $=10 \mathrm{kHz} ; \mathrm{Y0}, \mathrm{Y} 1(\mathrm{CHO})$ as output pulses; start/end frequency $($ D1340 $)=200 \mathrm{~Hz}$; acceleration time $($ D1343 $)=300 \mathrm{~ms}$; deceleration time $($ D1936 $)=$ 600ms; percentage value (D1131) = 100; current number of output pulses (D1648, D1649) $=0$.

- Assume the first execution result as:

- Observe the result of the first execution:

1) The actual output number 49,200 - estimated output number 50,000 $=-800$ (a negative value). $A$ negative value indicates that the entire execution finishes earlier and has not completed yet.
2) Try to shorten the acceleration time (D1343) into 250 ms and deceleration time (D1936) into 550ms.

- Obtain the result of the second execution:

- Observe the result of the second execution:

1) The actual output number 50,020 - estimated output number 50,000 $=20$
2) $20 \times(1 / 200 \mathrm{~Hz})=100 \mathrm{~ms}$ (idling time)
3) 100 ms is an appropriate value. Therefore, set the acceleration time as 250 ms and deceleration time as 550 ms to complete the design.

4, Program Example 2:

- Assume the feedback of the encoder is an A/B phase input and we adopt C251 timing (we suggust you clear it to 0 before the execution); target number of feedbacks $=50,000$; target output frequency $=100 \mathrm{kHz}$; Y0, Y1 (CH0) as output pulses; start/end frequency (D1340) $=200 \mathrm{~Hz}$; acceleration time (D1343) $=300 \mathrm{~ms}$; deceleration time $($ D1936 $)=600 \mathrm{~ms}$; precentage value $(\mathrm{D} 1131)=100$; current number of output pulses $(D 1648, D 1649)=0$.

|  |  |  |  |
| :---: | :---: | :---: | :---: |
|  | MOV | K100 | D1131 |
| $\begin{aligned} & \text { ON only } \\ & \text { for } 1 \\ & \text { scan } \\ & \text { oftor } \end{aligned}$ | MOV | K100 | D1244 |
|  | MOV | K100 | D1340 |
|  |  |  | $\begin{aligned} & \text { The 1st } \\ & \text { step } \\ & \text { start } \\ & \text { fromion } \end{aligned}$ |
|  | MOV | K100 | D1343 |
|  |  |  | Accelera /Deceler time of © H n |
|  | MOV | K0 | D1696 |
|  |  |  |  |
|  | DMOV | K0 | D1648 |
|  |  |  |  |
|  | DMOV | K0 | C251 |




- Assume the first execution result as:

- Observe the result of the first execution:

1) The actual output number 50,600 - estimated output number $50,000=600$
2) $600 \times(1 / 200 \mathrm{~Hz})=3 \mathrm{~s}$ (idling time)
3) 3 seconds are too long. Therefore, increase the percentage value (D1131) to K101.

- Obatin the result of the second execution:

- Observe the result of the second execution:

1) a) The actual output number 50,560 - estimated output number $50,500=60$
2) b) $60 \times(1 / 200 \mathrm{~Hz})=300 \mathrm{~ms}$ (idling time)
3) 300 ms is an appropriate value. Therefore, set the percentage value (D1131) as K101 to complete the design.

| $\begin{gathered} \text { ZL } \\ 198 \end{gathered}$ |  | VSPO |  |  |  |  |  | S 1 | S2 | S3 |  | D |  |  |  | Variable speed pulse output |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | DVSPO: 17 steps <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | Kns | T |  |  | C | D |  | E | F |
| S1 |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| S2 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| S3 |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Operands:

- $\quad$ s1: Target frequency of output $\quad$ S2: Target number of pulses
S3: Gap time and gap frequency
D: Pulse output device (Y0, Y2, Y4, and Y6.)
- Max frequency for $\mathrm{S1}: 200 \mathrm{kHz}$. Target frequency can be modified during the execution of instruction. When S1 is modified, VSPO will ramp Up/down to the target frequency according to the ramp-up gap time and gap frequency set in S3.
- $\quad$ S2 target number of pulses is valid only when the instruction is executed first time. $\$ 2$ can NOT be modified during the execution of instruction. S 2 can be a negative value. When target number of pulses are specified with 0, PLC will perform continuous output and the special $D$ shows the current value that is counting and going in the forward direction but that does NOT include any control over the output point direction.
- The gap frequency in $\mathrm{S} 3+0$ is in the range of 6 Hz to 32767 Hz , and the gap time in $\mathrm{S} 3+1$ is in the range of 1 ms to 80 ms . If a setting value exceeds the available range, the PLC will take the maximum or the minimum value.
- D pulse output device supports Y0, Y2, Y4 and Y6. Y1, Y3, Y5 and Y7 are corresponding output direction. The forward direction is On
- Parameters set in $S 3$ can only be modified while modifying the value in $S 1$. When target frequency is set as $0, \mathrm{PLC}$ will ramp down to stop according to parameters set in $S 3$. If target frequency other than 0 is specified again, pulse output will ramp up to target frequency and operates untill target number of pulses are completed
- Function Explanations::

1) Pulse output diagram

a. Definitions:
$\dagger l \rightarrow$ target frequency of $1^{\text {tt }}$ shift
$\dagger 2 \rightarrow$ target frequency of $2^{\text {nd }}$ shift
†3 $\rightarrow$ target frequency of $3^{\text {rd }}$ shift
gl $\rightarrow$ ramp-up time of $1^{\text {st }}$ shift
g2 $\rightarrow$ ramp-up time of $2^{\text {nd }}$ shift
g3 $\rightarrow$ ramp-down time of $3^{\text {rd }}$ shift
S2 $\rightarrow$ total output pulses
b. Explanations on each shift

1, $1^{\text {st }}$ shift:
Assume $\dagger 1=6 \mathrm{kHz}$, gap freqency $=1 \mathrm{kHz}$, gap time $=10 \mathrm{~ms}$ Ramp-up steps of 1 st shift:


2, $2^{\text {nd }}$ shift:
Assume $\dagger 2=11 \mathrm{kHz}$, internal frequency $=2 \mathrm{kHz}$, gap time $=20 \mathrm{~ms}$ Ramp-up steps of 2 nd shift:


3, 3rd shift:
Assume $\dagger 3=3 \mathrm{kHz}$, gap frequency $=2 \mathrm{kHz}$, gap time $=20 \mathrm{~ms}$ Ramp-down steps of 3 rd shift:


Points to note:

1. Associated flags:

M1029: CH0 pulse output execution is completed
M1030: CH 1 pulse output execution is completed
M1036: CH 2 pulse output execution is completed
M1037: CH3 pulse output execution is completed

## 10 Application instructions ZL200~ZL313

## 10.1 (ZL 202-203) Special function instructions



1, Explanations:

- To perform pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instruction
- S1:Source value S2: Slope S3: Offset D: Destination device.
- Operation equation in the instruction: $D=(S 1 \times S 2) \div 1,000+S 3$.

Users have to obtain S2 and S3 (decimals are rounded up into 16-bit integers) by using the slope and offset equations below.

Slope equation: $S 2=[($ max. destination value $-\min$. destination value $) \div(\max$. source value $-\min$. source value)] $\times 1,000$.

Offset equation: $\mathrm{S} 3=\mathrm{min}$. destination value -min . source value $\times \mathrm{S} 2 \div 1,000$
4


2, Program Example 1:

- Assume $S 1=500, S 2=168, S 3=-4$. When $X 0=$ On, SCAL instruction will be executed and obtain the proportional value at D0.
- Equation: $\mathrm{DO}=(500 \times 168) \div 1,000+(-4)=80$.



3. Program Example 2:

- Assume S1 $=500, S 2=-168, S 3=534$. When X10 $=$ On, SCAL instruction will be executed and obtain the proportional value at D10.
- Equation: $\mathrm{DO}=(500 \times-168) \div 1,000+534=450$


Remarks:

- This instruction is applicable for known slope and offset. If slope and offset are unknown, use SCLP instruction for the calculation.
- $\quad S 2$ has to be within the range $-32,768 \sim 32,767$. If $S 2$ falls without the range, use SCLP instruction for the calculation.
- When using the slope equation, please be aware that the max. source value must > min. source value, but it is not necessary that max. destination value > min. destination value.
- If the value of $D>32,767, D=32,767$; if the value of $D<-32,768, D=-32,768$.


1. Explanations:

- SI:Source value.

S2: Parameter.
D: Destination device

- Settings of S 2 for 16 -bit instruction:

S2 occupies 4 consecutive devices in 16 -bit instruction:

| Device No. | Parameter | Range |  |
| :---: | :---: | :---: | :---: |
| S2: | Maximum source value | Integer | Floating point |
| S2 +1 | Minimum source value | $\begin{gathered} -2,147,483,648 ~ \\ 2,147,483,647 \end{gathered}$ | Range of 32-bit <br> floating point |
| S2 +2 | Maximum destination value |  |  |
| S2 +3 | Minimum destination value |  |  |

- Settings of S 2 for 32-bit instruction:
- S2 occupies 8 consecutive devices in 32-bit instruction.

| Device No. | Parameter | Range |  |
| :---: | :---: | :---: | :---: |
|  |  | Integer | Floating point |
| S2, S2 + 1 | Maximum source value | $\begin{gathered} -2,147,483,648 ~ \\ 2,147,483,647 \end{gathered}$ | Range of 32-bit floating point |
| S2+2, 3 | Minimum source value |  |  |
| S2 + 4, 5 | Maximum destination value |  |  |
| S2 + 6, 7 | Minimum destination value |  |  |

- Operation equation in the instruction: $\mathrm{D}=[(\mathrm{S} 1-\mathrm{min}$. source value) $\times$ (max. destination value -min . destination value)] $\div$ (max. source value - min. source value) + min. destination value.
- The operational relation between source value and destination value is as stated below:
$y=k x+b$
$y=$ Destination value (D)
$\mathrm{k}=$ Slope $=($ max. destination value - min. destination value $) \div($ max. source value - min. source value $)$ $b=$ Offset $=$ Min. destination value - Min. source value $\times$ slope

Bring all the parameters into equation $y=k x+b$ and obtain the equation in the instruction:
$y=k x+b=D=k \quad S 1+b=$ slope $\times S 1+$ offset $=$ slope $\times S 1+$ min. destination value - min. source value $\times$ slope $=$ slope $\times(S 1-\min$. source value $)+$ min. destination value $=(S 1-m i n$. source value $) \times($ max. destination value - min. destination value) $\div$ (max. source value - min. source value) + min. destination value.

- If $S 1>$ max. source value, $S 1=$ max. source value

If S 1 < min. source value, $\mathrm{Sl}=\mathrm{min}$. source value
When all the input values and parameters are set, the output curve is shown as the figure:

2. Program Example 1:

- Assume $\mathrm{S} 1=500$, max. source value $\mathrm{DO}=3,000$, min. source value $\mathrm{D} 1=200$, max. destination value $\mathrm{D} 2=$ 500 , and min. destination value D3 $=30$. When $X 0=O n, S C L P$ instruction will be executed and obtain the proportional value at D10.
- Equation: D10 $=[(500-200) \times(500-30)] \div(3,000-200)+30=80.35$. Round off the result into an integer D10 $=80$.


3. Program Example 2:

- Assume $\mathrm{S} 1=500$, max. source value $\mathrm{DO}=3,000$, min. source value $\mathrm{D} 1=200$, max. destination value $\mathrm{D} 2=$ 30 , and min. destination value D3 $=500$. When $\mathrm{XO}=\mathrm{On}, \mathrm{SCLP}$ instruction will be executed and obtain the proportional value at DIO.
- Equation: $\mathrm{D} 10=[(500-200) \times(30-500)] \div(3,000-200)+500=449.64$. Round off the result into an integer D10 $=450$.




## 4. Program Example 3:

- Assume the source of S1 D100 = F500, max. source value D0 = F3000, min. source value D2 = F200, max. destination value D4 $=$ F500, and min. destination value D6 $=F 30$. When $X 0=O n$, set up M1162, adopt floating point operation and execute DSCLP instruction. The proportional value will be obtained at D10.
- Equation: $\mathrm{D} 10=[(F 500-F 200) \times(F 500-F 30)] \div(F 3000-F 200)+F 30=F 80.35$. Round off the result into an integer D10 = F80.


Remarks:

- Range of $S 1$ for 16 -bit instruction: max. source value $\geqslant S 1 \geqslant \min$. source value; $-32,768 \sim 32,767$. If the value falls without the bounds, the bound value will be used for calculation.
- Range of integer S1 for 32-bit instruction: max. source value $\geqslant S 1 \geqslant \min$. source value; -2,147,483,648~ $2,147,483,647$. If the value falls without the bounds, the bound value will be used for calculation.
- Range of floating point $S 1$ for 32-bit instruction: max. source value $\geqslant S 1 \geqslant$ min. source value; following the range of 32- bit floating point. If the value falls without the bounds, the bound value will be used for calculation.
- Please be aware that the max. source value must > min. source value, but it is not necessary that max. destination value $>$ min. destination value.

| $\begin{gathered} \text { ZL } \\ 205 \end{gathered}$ | D | CMPT |  |  |  | P | S 1 S |  |  | 2 n |  | D |  |  |  | Compare table |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | CMPT: 9 steps <br> DCMPT: 17 steps <br> DCMPTP: 17 steps |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |
| n |  |  |  |  | * | * |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |

1. Explanations:

- To perform a pulse type, queve by adding an NP rising edge " $\uparrow$ " instruction to the front of the instr uction.
- Sl : Source device 1

S2: Source device 2
n : Data length/function
D: Destination device

- $\quad$ S1 and $\$ 2$ can be T/C/D devices, for C devices only 16-bit devices are applicable (C0~C199).
- The high 16 -bit value in the operand n used in the 32 -bit instruction is an invalid value.
- The low 8 -bit value in the operand n indicates the data length. The operand n used in the 16 -bit instruction should be within the range between 1 and 16. The operand n used in the 32 -bit instruction should be within the range between 1 and 32. PLC will take the upper/lower bound value if set value exceeds the available range.
- The high 8 -bit value in the operand n indicates the comparison condition.

| Value | K 0 | K 1 | K 2 | K 3 | K 4 |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Comparison condition | $\mathrm{S} 1=\mathrm{S} 2$ | $\mathrm{~S} 1<\mathrm{S} 2$ | $\mathrm{~S} 1<=\mathrm{S} 2$ | $\mathrm{~S} 1>\mathrm{S} 2$ | $\mathrm{~S} 1>=\mathrm{S} 2$ |

- If n used in the 16 -bit instruction is set to $\mathrm{HO108}$, it means that 8 pieces of data are compared to 8 pieces of data, and the "larger than" comparison is performed. If $n$ used in the 32 -bit instruction is set to H00000320, it means that 32 pieces data are compared to 32 pieces of data, and the "less than" comparison is performed.
- If the setting value for the comparison condition exceeds the range, or the firmware version does not support the comparison condition, the default "equal to" comparison is performed.
- The comparison values used in the 16 -bit instruction are signed values. The comparison values used in the 32- bit instruction are 32-bit values (M1162=Off), or floating-point values (M1162=On).
- Data written in operand $D$ will all be stored in 16-bit format or in 32-bit format. When data length is less than 16 or 32 , the null bits are fixed as 0, e.g. if $\mathrm{n}=\mathrm{K8}$, bit $0 \sim 7$ will be set according to compare results, and bit 8~15 will all be 0 .
- If the comparison result meets the condition, the corresponding bit is set to 1 . Otherwise, it is set to 0 .


## 2, Program example:

- When $\mathrm{MO}=\mathrm{ON}$, compare the 16-bit value in D0~D7 with D20~D27 and store the results in D100.

| M0 |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |

- Content in D0~D7:

| number | D0 | D1 | D2 | D3 | D4 | D5 | D6 | D7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| value | K10 | K20 | K30 | K40 | K50 | K60 | K70 | K80 |

- Content in D20 ${ }^{\sim}$ D27:

| number | D20 | D21 | D22 | D23 | D24 | D25 | D26 | D27 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| value | K12 | K20 | K33 | K44 | K50 | K66 | K70 | K88 |

- After the comparison of CMPT instruction, the associated bit will be 1 if two devices have the same value, and other bits will all be 0 . Therefore the results in D100 will be as below:

| D100 | BitO | Bit 1 | Bit02 | Bit03 | Bit04 | Bit05 | Bit06 | Bit07 | Bit8~15 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | $0 \cdots 0$ |
|  | H0052 (K82) |  |  |  |  |  |  |  |  |


| $\begin{gathered} \text { ZL } \\ 207 \end{gathered}$ |  | CSFO |  |  |  |  | S Sl S |  |  |  |  |  |  |  |  | Catch speed and proportional output |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | CSFO: 7 steps <br> 16-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S | * |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| S1 |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |
| D |  |  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |

1. Explanations:

- To perform a pulse type, queue by adding an NP rising edge " $\uparrow$ " instruction to the front of the instructio n.
- $s$ : Source device of signal input (Only XO and XI are available)

S1: Sample time setting and the input speed information
D: Output proportion setting and output speed information

- When $S$ specifies X0, PLC only uses XO input point and its associated high speed pulse output: Y0, in this case Y 1 is normal output point. When S specifies $\mathrm{XI}, \mathrm{PLC}$ uses XO (A phase) and XI ( $B$ phase) input points and their associated output: Y0 (Pulse) / Y1 (Dir).
- If $S$ specifies XI with 2-phase 2 inputs, the counting mode is fixed as quadruple frequency.
- During pulse output process of YO, special registers (D1649, D1648) storing the current number of output pulses will be updated when program scan proceeds to this instruction.
- $\quad$ S1 occupies consecutive 416 -bit registers. $S 1+0$ specifies the sampling times, i.e. when $\mathrm{S} 1+0$ specifies K 1 , PLC catches the speed every time when 1 pulse is outputted. Valid range for $\$ 1+0$ in 1 -phase 1 -input mode: K1~K100, and 2-phase 2-input mode: K2~K100. If the specified value exceeds the valid range, PLC will take the lower/upper bound value as the set value. Sample time can be changed during PLC operation, however the modified value will take effect until program scan proceeds to this instruction. S1+1 indicates the latest speed sampled by PLC (Read-only). Unit: 1 Hz . Valid range: $\pm 10 \mathrm{kHz}$. S1+2 and S1+3 indicate the accumulated number of pulses in 32-bit data (Read-only).
- D occupies 3 consecutive 16-bit registers. D +0 specifies the output proportion value. Valid range: K1 ( $1 \%$ ) ~K10000 ( $10000 \%$ ). If the specified value exceeds the valid range, PLC will take the lower/upper bound value as the set value. Output proportion can be changed during PLC operation, however the modified value will take effect until program scan proceeds to this instruction. $\mathrm{D}+2$ and $\mathrm{D}+1$ indicates the
output speed in 32-bit data. Unit: 1Hz. Valid range: $\pm 200 \mathrm{kHz}$. When selecting the MPG mode, it takes one more 16 -bit register. $\mathrm{D}+3$ indicates the pulse output channel, ranging from KO to K 3 , indicating the output channels $\mathrm{CHO} \sim \mathrm{CH} 3$. When selecting the general mode, the pulse output channel is fixed to CHO . Note: if you need to change the mode from the MPG mode to the general mode or vise versa, you need to close the instruction and re-execute the instruction to ensure the channel switching can be normally done。
- The pulse output channel selecting: when S input point uses XO as the source, the corresponding pulse output points are $\mathrm{Y} 0, \mathrm{Y} 2, \mathrm{Y} 4, \mathrm{Y} 6$ and the general pulse output points are $\mathrm{Y} 1, \mathrm{Y} 3, \mathrm{Y} 5, \mathrm{Y} 7$. When S input point uses $\mathrm{X1}$ as the source, the corresponding output points are Y0(Pulse) / Y1(Dir) or Y2(Pulse) / Y3(Dir) or Y4(Pulse) /Y5(Dir) or Y6(Pulse) / Y7(Dir) high speed output.
- The speed sampled by the PLC will be multiplied with the output proportion $\mathrm{D}+0$, then the PLC will generate the actual output speed. The PLC will take the integer of the calculated value, i.e. if the calculated result is smaller than 1 Hz , the PLC will output with 0 Hz . For example, input speed: 10 Hz , output proportion: $\mathrm{K} 5(5 \%)$, then the calculation result will be $10 \times 0.05=0.5 \mathrm{~Hz}$. Pulse output will be 0 Hz ; if output proportion is modified as $\mathrm{K} 15(15 \%)$, then the calculation result will be $10 \times 0.15=1.5 \mathrm{~Hz}$. Pulse output will be 1 Hz .


## 2, Program Example:

- If DO is set as K2, D10 is set as K100:

When the sampled speed on (X0, X1) is $+10 \mathrm{~Hz}(\mathrm{DI}=\mathrm{K} 10),(\mathrm{YO}, \mathrm{Y} 1)$ will output pulses with $+10 \mathrm{~Hz}(\mathrm{D} 12, \mathrm{D} 11$ $=\mathrm{K} 10)$; When the sampled speed is $-10 \mathrm{~Hz}(\mathrm{Dl}=\mathrm{K}-10),(\mathrm{YO}, \mathrm{Y})$ will output pulses with $-10 \mathrm{~Hz}(\mathrm{D} 12, \mathrm{D} 11=\mathrm{K}-$ 10).

- If DO is set as K2, D10 is set as K1000:

When the sampled speed on $(\mathrm{XO}, \mathrm{X1})$ is $+10 \mathrm{~Hz}(\mathrm{Dl}=\mathrm{K} 10),(\mathrm{YO}, \mathrm{Y} 1)$ will output pulses with $+100 \mathrm{~Hz}(\mathrm{D} 12, \mathrm{D} 11$ $=\mathrm{K} 100$ ); When the sampled speed is $-100 \mathrm{~Hz}(\mathrm{DI}=\mathrm{K}-100)$, (YO, Y1) will output pulses with -100 Hz (D12, D11 $=K-100)$.

- If DO is set as K10, D10 is set as K10:

When the sampled speed on $(\mathrm{XO}, \mathrm{X})$ is $+10 \mathrm{~Hz}(\mathrm{D} 1=\mathrm{K} 10),(\mathrm{YO}, \mathrm{Y} 1)$ will output pulses with $+1 \mathrm{~Hz}(\mathrm{D} 12, \mathrm{D} 11=$ K1); When the sampled speed is $-10 \mathrm{~Hz}(\mathrm{Dl}=\mathrm{K}-10)$, $(\mathrm{YO}, \mathrm{Y} 1)$ will output pulses with $-1 \mathrm{~Hz}(\mathrm{D} 12, \mathrm{D} 11=\mathrm{K}-1)$.

| M0 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| H | CSFO | X 1 | D0 | D10 |


| $\begin{gathered} \text { ZL } \\ 215 \sim 217 \end{gathered}$ |  | LD\# |  |  |  |  |  |  | $S 1$ S2 |  |  |  |  |  |  | Contact Logical Operation LD\# |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | LD\#: 5 steps DLD\#: 9 step | 16-bit <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY |  |  | KnM | KnS | T | C | D |  |  | E | F |
| S 1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |

1, Explanations:

- S1: Data source device $1 \quad$ S2: Data source device 2
- This instruction compares the content in $S 1$ and $S 2$. If the result is not " 0 ", the continuity of the instruction is enabled. If the result is " 0 ", the continuity of the instruction is disabled.
- LD\# instruction is used for direct connection with BUS

| ZL No. | 16-bit instruction | 32-bit instruction | Continuity <br> condition | No-continuity <br> condition |
| :--- | :--- | :--- | :--- | :--- |
| 215 | LD\& | DLD\& | S1 \& S2 $\neq 0$ | S1 \& S2 $=0$ |
| 216 | LD $\mid$ | DLD $\mid$ | S1 $\mid S 2 \neq 0$ | S1 $\mid S 2=0$ |
| 217 | LD^ | DLD $\wedge$ | S2 $\neq 0$ | S $1 \wedge S 2=0$ |

- \&: Logical "AND" operation.
- |: Logical "OR" operation.
- $\wedge$ : Logical "XOR" operation
- When 32-bit counters (C200 ~ C255) are used in this instruction for comparison, make sure to adopt 32-bit instruction (DLD\#). If 16-bit instructions (LD\#) is adopted, a "program error" will occur and the ERROR indicator on the MPU panel will flash.

2, Program Example:

- When the result of logical AND operation of $C 0$ and $C 10 \neq 0, Y 10=O n$.
- When the result of logical OR operation of D200 and D300 $\neq 0$ and $X 1=O n, Y 11=O n$ will be retained.
- When the result of logical XOR operation of C201 and C200 $\neq 0$ or M3 $=O n, M 50=O n$.


| $\begin{gathered} \text { ZL } \\ 218 \sim 220 \end{gathered}$ |  | AND\# |  |  |  |  | $S 1$ S2 |  |  |  |  |  |  |  |  | Contact Logical Operation AND\# |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | AND\#: 5 steps DAND\#: 9 steps | 16-bit <br> 32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S 1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |

1. Explanations:

- S1: Data source device $1 \quad$ S2: Data source device 2
- This instruction compares the content in $S 1$ and $S 2$. If the result is not " 0 ", the continuity of the instruction is enabled. If the result is " 0 ", the continuity of the instruction is disabled.
- AND\# is an operation instruction used on series contacts.

| ZL No. | 16-bit instruction | 32-bit instruction | Continuity <br> condition | No-continuity <br> condition |
| :--- | :--- | :--- | :--- | :--- |
| 218 | AND\& | DAND\& | $S 1 \& S 2 \neq 0$ | $S 1 \& S 2=0$ |
| 219 | AND $\mid$ | DAND 1 | $S 1 \mid S 2 \neq 0$ | $S 1 \mid S 2=0$ |
| 220 | AND^ | DAND^ | $S 1 \wedge S 2 \neq 0$ | $S 1 \wedge S 2=0$ |

- \&:Logical "AND" operation
- |:Logical "OR" operation
- $\wedge$ : Logical "XOR" operation.
- When 32-bit counters (C200 ~ C255) are used in this instruction for comparison, make sure to adopt 32-bit instruction (DAND\#). If 16-bit instructions (AND\#) is adopted, a "program error" will occur and the ERROR indicator on the MPU panel will flash.

2, Program Example:

- When $\mathrm{XO}=\mathrm{On}$ and the result of logical AND operation of CO and $\mathrm{C} 10 \neq 0, Y 10=O$.
- When $\mathrm{XI}=\mathrm{Off}$ and the result of logical OR operation of D 10 and $\mathrm{DO} \neq 0$ and $\mathrm{XI}=\mathrm{On}, \mathrm{Y} 11=$ On will be
retained.
- When X2 = On and the result of logical XOR operation of 32-bit register D200 (D201) and 32-bit register D100 (D101) $\neq 0$ or M3 $=O n, M 50=O n$.



1, Explanations:

- S1: Data source device $1 \quad$ S2: Data source device 2
- This instruction compares the content in $S 1$ and $S 2$. If the result is not " 0 ", the continuity of the instruction is enabled. If the result is " 0 ", the continuity of the instruction is disabled.
- OR\# is an operation instruction used on parallel contacts.

| ZL No. | 16-bit instruction | 32-bit instruction | Continuity <br> condition | No-continuity <br> condition |
| :--- | :--- | :--- | :--- | :--- |
| 221 | OR\& | DOR\& | $S 1 \& S 2 \neq 0$ | $S 1 \& S 2=0$ |
| 222 | OR $\mid$ | DOR $\mid$ | S1 $\mid S 2 \neq 0$ | $S 1 \mid S 2=0$ |
| 223 | OR^ | DOR^ | $S 1 \wedge S 2 \neq 0$ | $S 1 \wedge S 2=0$ |

- \&: Logical "AND" operation.
- |: Logical "OR" operation.
- $\wedge$ : Logical "XOR" operation.
- When 32-bit counters (C200 ~ C255) are used in this instruction for comparison, make sure to adopt 32-bit instruction (DOR\#). If 16 -bit instructions (OR\#) is adopted, a "program error" will occur and the ERROR indicator on the MPU panel will flash.

2, Program Example:

- When $\mathrm{X1}=$ On and the result of logical AND operation of CO and $\mathrm{ClO} \neq 0, Y 10=O n$.
- M60 will be On when X2 $=$ On and M30 $=$ On, or the result of logical OR operation of 32-bit register D10 (D11) and 32-bit register D20 (D21) $\neq 0$, or the result of logical XOR operation of 32-bit register D200 (D201) and 32- bit counter C235 $\neq 0$.



## 10.3 (ZL 224-246) Contact type comparison command

| $\begin{array}{\|c\|} \hline \text { ZL } \\ 224 \sim 230 \end{array}$ | D | LD※ |  |  |  |  | S 1 S2 |  |  |  |  |  |  |  |  | Load Compare |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | LD※: 5 steps <br> DLD※: 9 steps | $\begin{aligned} & \text { 16-bit } \\ & \text { 32-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |

1. Explanations:

- S1: Data source device $1 \quad$ S2: Data source device 2.
- This instruction compares the content in S 1 and S 2 . Take API 224 ( $\mathrm{LD}=$ ) for example, if the result is " $=$ " , the continuity of the instruction is enabled. If the result is " $\neq$ ", the continuity of the instruction is disabled.
- LD※ instruction is used for direct connection with BUS.

| ZL No. | 16-bit instruction | 32-bit instruction | Continuity condition | No-continuity condition |
| :---: | :---: | :---: | :---: | :---: |
| 224 | LD= | DLD= | S1 = S2 | S1 $=$ S2 |
| 225 | LD> | DLD> | S1 > S2 | S1 $\leqq$ S2 |
| 226 | LD< | DLD< | S1 < S2 | S1 $\geqq$ S2 |
| 228 | LD<> | DLD<> | S1 $\ddagger$ S2 | S1 = S2 |
| 229 | LD<= | DLD<= | S1 $\leqq$ S2 | S1 > S2 |
| 230 | LD>= | DLD>= | S1 $\geqq$ S2 | S1 < S2 |

- When 32-bit counters (C200 ~ C255) are used in this instruction for comparison, make sure to adopt 32-bit instruction (DLD※). If 16-bit instructions (LD※) is adopted, a "program error" will occur and the ERROR indicator on the MPU panel will flash.

2, Program Example:

- When the content in $\mathrm{C} 10=\mathrm{K} 200, \mathrm{Y} 10=\mathrm{On}$.
- When the content in D200 > K-30 and X1 $=$ On, Y11 = On will be retained.
- When the content in $\mathrm{C} 200<\mathrm{K} 678,493$ or M3 $=$ On, M50 $=$ On.


| $\begin{array}{\|c\|} \hline \text { ZL } \\ 232 ~ 238 \end{array}$ |  | AND |  |  |  |  | S 1 |  |  | S2 |  |  |  |  |  | AND Compare |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | AND ※: 5 steps <br> DAND※: 9 steps |
|  | X | Y | M | S | K | H | KnX | KnY | KnM |  |  | KnS | T | C | D |  | E | F |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |

1. Explanations:

- S1: Data source device 1 S2: Data source device 2
- This instruction compares the content in S1 and S2. Take API232 (AND=) for example, if the result is "=" , the continuity of the instruction is enabled. If the result is " $\neq$ ", the continuity of the instruction is disabled.
- AND $※$ is a comparison instruction is used on series contacts.

| ZL No. | 16-bit instruction | 32-bit instruction | Continuity condition | No-continuity condition |
| :---: | :---: | :---: | :---: | :---: |
| 232 | AND= | DAND= | S $1=$ S2 | S1 $\neq$ S2 |
| 233 | AND> | DAND> | S1 > S2 | S1 S 2 |
| 234 | AND< | DAND< | S1 < S2 | $S 1 \geqq$ S2 |
| 236 | AND<> | DAND<> | S1 $\neq$ S2 | S1 = S2 |
| 237 | AND<= | DAND<= | S1 S 2 | $\mathrm{S} 1>\mathrm{S} 2$ |
| 238 | AND>= | DAND>= | S1 $\geqq$ S2 | S1 < S2 |

- When 32-bit counters (C200 ~ C255) are used in this instruction for comparison, make sure to adopt 32 -bit instruction (DAND※). If 16 -bit instructions (AND※) is adopted, a "program error" will occur and the ERROR indicator on the MPU panel will flash.

2. Program Example:

- When $\mathrm{XO}=$ On and the content in $\mathrm{ClO}=\mathrm{K} 200, \mathrm{Y} 10=\mathrm{On}$.
- When $\mathrm{XI}=$ Off and the content in $\mathrm{DO} \neq \mathrm{K}-10, \mathrm{Y} 11=$ On will be retained.
- When $\mathrm{X} 2=$ On and the content in 32-bit register $D 0(D 11)<678,493$ or $M 3=O n, M 50=$ On.



1, Explanations:

- S1: Data source device 1 S2: Data source device 2
- This instruction compares the content in S1 and S2. Take API240 (OR=) for example, if the result is "=" , the continuity of the instruction is enabled. If the result is " $\neq$ ", the continuity of the instruction is disabled.
- OR※ is an comparison instruction used on parallel contacts.

| ZL No. | 16-bit instruction | 32-bit instruction | Continuity <br> condition | No-continuity <br> condition |
| :--- | :--- | :--- | :--- | :--- |
| 240 | OR= | DOR= | $\mathrm{S} 1=\mathrm{S} 2$ | $\mathrm{~S} 1 \neq \mathrm{S} 2$ |
| 241 | $\mathrm{OR}>$ | DOR> | $\mathrm{S} 1>\mathrm{S} 2$ | $\mathrm{~S} 1 \leqq \mathrm{~S} 2$ |
| 242 | $\mathrm{OR}<$ | DOR< | $\mathrm{S} 1<\mathrm{S} 2$ | $\mathrm{~S} 1 \geqq \mathrm{~S} 2$ |
| 244 | $\mathrm{OR}<>$ | DOR<> | $\mathrm{S} 1 \neq \mathrm{S} 2$ | $\mathrm{~S} 1=\mathrm{S} 2$ |
| 245 | OR<= | DOR<= | $\mathrm{S} 1 \leqq \mathrm{~S} 2$ | $\mathrm{~S} 1>\mathrm{S} 2$ |
| 246 | OR>= | DOR>= | $\mathrm{S} 1 \geqq \mathrm{~S} 2$ | $\mathrm{~S} 1<\mathrm{S} 2$ |

- When 32-bit counters (C200~C255) are used in this instruction for comparison, make sure to adopt 32-bit instruction (DOR ※). If 16-bit instructions (OR ) is adopted, a "program error" will occur and the ERROR indicator on the MPU panel will flash.

2, Program Example:

- When $\mathrm{XI}=$ On and the present value of $\mathrm{C} 10=\mathrm{K} 200, \mathrm{YO}=\mathrm{On}$.
- M60 will be On when X2 = On, M30 = On and the content in 32-bit register D100 (D101) $\geqslant$ K100,000.



## 10.4 (ZL 266-274) Character device bit command



1, Explanations:

- D: Destination output device n: Device specifying the output bit..
- BOUT instruction performs bit output on the output device according to the value specified by operand n. Status of Coils and Associated Contacts:

| Evaluation result | BOUT instruction |  | Associated Contacts |  |
| :--- | :--- | :--- | :--- | :---: |
|  | Coil | NO contact (normally <br> open) | NC contact (normally <br> closed) |  |
|  |  | Current blocked | Current flows |  |
| TRUE |  | Current flows | Current blocked |  |

2 Program Example:

| $\mathrm{H}_{\mathrm{Ho}}^{\mathrm{x} 1}$ |  |  |  | Instruction: |  | Operation: <br> Load NC contact X0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | BOUT | K4Y0 | D0 | LDI | X0 |  |
|  |  |  |  | AND | X1 | Connect NO contact |
|  |  |  |  |  |  | X1 in series. |
|  |  |  |  | BOUT | K4Y0 | D0 When $\mathrm{D} 0=\mathrm{k} 1$, executes output on Y 1 |
|  |  |  |  |  |  | When D0 = k2, executes output on Y2 |



1. Explanations:

- D: Destination device to be Set ON n: Device specifying the bit to be Set ON
- When BSET instruction executes, the output device specified by operand $n$ will be On and latched. To reset the On state of the device, BRST instruction is required.

2, Program Example:


| Instruction: | Operation: |  |
| :--- | :--- | :--- |
| LDI | X0 | Load NC contact X0 |
| AND | X1 |  |
|  |  | Connect NO contact |
|  | X1 in series. |  |
| BSET | K4Y0 $\quad$ D0 | When D0 $=k 1$, |
|  |  | Y1 is ON and latched |
|  |  | When D0 $=k 2$, |
|  |  | Y2 $=\mathrm{ON}$ and latched |


| $\begin{gathered} \text { ZL } \\ 268 \end{gathered}$ |  | BRST |  |  |  |  |  |  | D $n$ |  | n |  |  |  |  | Reset Specified Bit of a Word |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | BRST: 5 steps DBRST: 9 step | $\begin{aligned} & \text { 16-bit } \\ & \text { 32-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY |  |  | KnM | KnS | T | C | D |  |  | E | F |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| n |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |

1. Explanations:

- D: Destination device to be reset n : Device specifying the bit to be reset
- When BRST instruction executes, the output device specified by operand $n$ will be reset (OFF).

2, Program Example:



1. Explanations:

- S : Reference source device n : Reference bit
- Available range for the value in operand n: K0~K15 for 16-bit instruction; K0~K31 for 32-bit instruction.
- BLD instruction is used to load NO contact whose contact state is defined by the reference bit n in reference device D, i.e. if the bit specified by n is ON , the NO contact will be ON , and vice versa.

2, Program Example:


Instruction:

| BLD | D0 | K3 | Load NO contact X0 with bit |
| :--- | :--- | :--- | :--- |
| OUT | Y0 | Status of bit 3 in D0 |  |



1. Explanations:

- S : Reference source device n : Reference bit
- BLD instruction is used to load NC contact whose contact state is defined by the reference bit n in reference device $D$, i.e. if the bit specified by $n$ is $O N$, the $N C$ contact will be $O N$, and vice versa.

2. Program Example:

| BLDI | D0 | K1 | Instruction: | Operation: <br> BLDI |
| :--- | :--- | :--- | :--- | :--- | :--- |
| OO K1 | Load NC contact with bit <br> status of bit1 in D0 |  |  |  |
| OUT Y0 | Drive coil Y0 |  |  |  |


| $\begin{gathered} \text { ZL } \\ 271 \end{gathered}$ |  | BAND |  |  |  |  |  |  | S n |  |  |  |  |  |  | Connect NO Contact in Series by Specified Bit |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | BAND: 5 steps DBAND: 9 steps | $\begin{aligned} & \text { 16-bit } \\ & \text { 32-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY |  |  | KnM | KnS | T | C | D |  |  | E | F |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| n |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |

1. Explanations:

- S : Reference source device n : Reference bit
- BAND instruction is used to connect NO contact in series. The current state of the contact which is connected in series is read, and then the logical AND operation is performed on the current state and the previous logical operation result. The final result is stored in the accumulative register.

2, Program Example:



1．Explanations：
－ S ：Reference source device n ：Reference bit
－BANI instruction is used to connect NC contact in series．The current state of the contact which is connected in series is read，and then the logical AND operation is performed on the current state and the previous logical operation result．The final result is stored in the accumulative register．

2，Program Example：


| X1 |  |  |
| :--- | :--- | :--- |
| BANI | DO | KO |

指令：
LDI X1
BANI
D0
OUT YO

操作说明：
载入 X 1 的 B 接点
串联 D0 Bit0 的 B 接点
驱动 YO 线圈

| $\begin{gathered} \text { ZL } \\ 273 \end{gathered}$ | D | BOR |  |  |  |  |  |  | n |  | n |  |  |  |  | Connect NO Contact in Parallel by Specified Bit |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | BOR: 5 steps $\quad 16$-bitDBOR: 9 steps 32 -bit |  |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| D |  |  |  |  |  |  |  | * | * | * | * | * | * |  |  |  |  |
| n |  |  |  |  | * | * | * | * | * | * | * | * | * | * | * |  |  |

1. Explanations:

- S : Reference source device n : Reference bit
- BOR instruction is used to connect NO contact in parallel. The current state of the contact which is connected in series is read, and then the logical OR operation is performed on the current state and the previous logical operation result. The final result is stored in the accumulative register.

2, Program Example:


| Instruction: | Operation: |  |  |
| :--- | :--- | :--- | :--- |
| LD | X0 | Load NO contact X0 |  |
| BOR | D0 | K0 | Connect NO contact in <br>  |
|  |  | parallel, whose state is <br> defined by bit0 of D0 |  |
| OUT | Y1 | Drive coil Y1 |  |



1. Explanations:

- S : Reference source device n : Reference bit
- BORI instruction is used to connect NC contact in parallel. The current state of the contact which is connected in series is read, and then the logical OR operation is performed on the current state and the previous logical operation result. The final result is stored in the accumulative register.

2, Program Example:


| Instruction: |  | Operation: |  |
| :--- | :--- | :--- | :--- |
| LD | X0 | Load NO contact X0 |  |
| BORI | DO | K0 | Connect NC contact in | parallel, whose state is defined by bit0 of D0

OUT Y1 Drive coil Y1

## 10.5 (ZL 275-313) Floating point contact type command

| $\begin{gathered} \text { ZL } \\ 275-280 \end{gathered}$ |  | FLD $\%$ |  |  |  |  | S 1 |  |  | S2 |  |  |  |  |  | Floating Point Contact <br> Type Comparison |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | FLD※: 9 steps32-bit |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |
| S1 |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |
| S2 |  |  |  |  |  |  |  |  |  |  | * | * | * |  |  |  |

1. Explanations:

- $s 1$ :Source device $1 \quad$ S2: Source device 2
- This instruction compares the content in S 1 and S 2 . Take API275 (FLD=) for example, if the result is " $=$ " , the continuity of the instruction is enabled. If the result is " $\neq$ ", the continuity of the instruction is disabled.
- The user can specify the floating point value directly into operands S1 and S2 (e.g. F1.2) or store the floating point value in D registers for further operation.
- FLD $(\nVdash:=,>,<,<>, \leqslant, \geqslant)$ instruction is used for direct connection with left hand bus bar.

| ZL No. | 32-bit instruction | Continuity condition | Discontinuity condition |
| :---: | :---: | :---: | :---: |
| 275 | FLD = | S1 = S2 | S1 $=$ S2 |
| 276 | FLD > | S1 > S2 | S1 § S2 |
| 277 | FLD < | S1 < S2 | S1 $\geqq$ S2 |
| 278 | FLD < > | S1 $\neq$ S2 | S1 = S2 |
| 279 | FLD $<=$ | S1 $\leqq$ S2 | S1 > S2 |
| 280 | FLD > = | S1 $\geqq$ S2 | S1 < S2 |

## 2. Program Example:

- When the content in D200 (D201) $\leqslant \mathrm{Fl} .2$ and X 1 is $\mathrm{ON}, \mathrm{Y} 21=\mathrm{ON}$ and latched.



1. Explanations:

- S1:Source device $1 \quad$ S2: Source device 2
- This instruction compares the content in S 1 and S 2 . Take API 281 (FAND =) for example, if the result is " $=$ ", the continuity of the instruction is enabled. If the result is " $\neq$ ", the continuity of the instruction is disabled.
- The user can specify the floating point value directly into operands S1 and S2 (e.g. F1.2) or store the floating point value in D registers for further operation.
- FAND $(\mathbb{F}:=,>,<,<>, \leqslant, \geqslant$ ) instruction is used for serial connection with contacts.

| ZL No. | 32-bit instruction | Continuity condition | Discontinuity condition |
| :--- | :--- | :--- | :--- |
| 281 | FAND $=$ | S1 $=$ S2 | S1 $\neq$ S2 |
| 282 | FAND $>$ | S1 $>$ S2 | S1 S2 |
| 283 | FAND $<$ | S1 $<$ S2 | S1 S2 |
| 284 | FAND $<>$ | S1 $\neq$ S2 | S1 $=$ S2 |
| 285 | FAND $<=$ | S1 | FAND $>=$ |
| 286 | S1 S2 | S1 $>$ S2 |  |

2, Program Example

- When X1 is OFF and the content in D0 (D1) does not equal to F1.2, Y21 = ON and latched.



1. Explanations:

- S1:Source device 1 S2: Source device 2
- This instruction compares the content in S1 and S2. Take API287 (FOR =) for example, if the result is "=" , the continuity of the instruction is enabled. If the result is " $\neq$ ", the continuity of the instruction is disabled.
- The user can specify the floating point value directly into operands S1 and S2 (e.g. F1.2) or store the floating point value in $D$ registers for further operation.
- OR※ ( $※:=,>,<,<>, \leqslant, \geqslant$ ) instruction is used for parallel connection with contacts.

| ZL No. | 32-bit instruction | Continuity condition | Discontinuity condition |
| :--- | :--- | :--- | :--- |
| 287 | FOR $=$ | $\mathrm{S} 1=\mathrm{S} 2$ | $\mathrm{~S} 1 \neq \mathrm{S} 2$ |
| 288 | FOR $>$ | $\mathrm{S} 1>\mathrm{S} 2$ | $\mathrm{~S} 1 \leqq \mathrm{~S} 2$ |
| 289 | $\mathrm{FOR}<$ | $\mathrm{S} 1<\mathrm{S} 2$ | $\mathrm{~S} 1 \geqq \mathrm{~S} 2$ |
| 290 | FOR < > | $\mathrm{S} 1 \neq \mathrm{S} 2$ | $\mathrm{~S} 1=\mathrm{S} 2$ |
| 291 | FOR $=$ | $\mathrm{S} 1 \leqq \mathrm{~S} 2$ | $\mathrm{~S} 1>\mathrm{S} 2$ |
| 292 | FOR $>=$ | $\mathrm{S} 1 \geqq \mathrm{~S} 2$ | $\mathrm{~S} 1<\mathrm{S} 2$ |

2, Program Example:

- When both X 2 and M30 are OFF and the content in D100 (D101) $\geqslant \mathrm{FI} .234, \mathrm{M} 60=\mathrm{ON}$.


| $\begin{array}{\|c\|} \hline \text { ZL } \\ 296 \sim 301 \end{array}$ | D | LDZ $\ldots$ |  |  |  |  |  | S 1 |  |  | S3 |  |  |  |  | Comparing contact type absolute values LDZ※ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bit Devices |  |  |  | Word Devices |  |  |  |  |  |  |  |  |  |  | $\begin{aligned} & \text { LDZ }: 7 \text { steps } \\ & \text { DLDZ } \#: 13 \text { steps } \end{aligned}$ | $\begin{aligned} & \text { 16-bit } \\ & \text { 32-bit } \end{aligned}$ |
|  | X | Y | M | S | K | H | KnX | KnY | KnM | KnS | T | C | D | E | F |  |  |
| S1 |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |  |  |
| S2 |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |  |  |
| A3 |  |  |  |  | * | * | * | * | * | * | * | * | * |  |  |  |  |

1, Explanations:

- S1:Source device $1 \quad$ S2:Source device 2 S3:Source device 3
- The absolute value of the difference between $S 1$ and $S 2$ is compared with the absolute value of $S 3$. Take LDZ> for example. If the comparison result is that the absolute value of the difference between S1 and S2 is greater than the absolute value of $\$ 3$, the condition of the instruction is met. If the comparison result is that the absolute value of the difference between S 1 and S 2 is less than or equal to the absolute value of S3, the condition of the instruction is not met.
- The instruction can be connected to a busbar:

| ZL No. | 16-bit instruction | 32-bit instruction | On | Off |
| :---: | :---: | :---: | :---: | :---: |
| 296 | LDZ> | DLDZ> | $\mid S 1-$ S2 $\|>\|S 3\|$ | $\mid$ S1-S2\| $\leqq \mid$ S3\| |
| 297 | LDZ>= | DLDZ>= | $\|S 1-S 2\| \geqq\|S 3\|$ | $\mid S 1-$ S2 $\|<\|S 3\|$ |
| 298 | LDZ< | DLDZ< | $\mid S 1-$ S2 $\|<\|S 3\|$ | $\|S 1-S 2\| \geqq\|S 3\|$ |
| 299 | LDZ<= | DLDZ<= | $\mid S 1-$ S2 $\|\leqq\|S 3\|$ | $\mid S 1-$ S2 $\|>\|S 3\|$ |
| 300 | LDZ $=$ | DLDZ = | $\|S 1-S 2\|=\|S 3\|$ | $\|S 1-S 2\| \neq\|S 3\|$ |
| 301 | LDZ<> | DLDZ<> | $\|\mathrm{S} 1-\mathrm{S} 2\| \neq\|\mathrm{S} 3\|$ | $\|S 1-\mathrm{S} 2\|=\|S 3\|$ |

- A 32-bit counter (C200~C255) must be used with the 32-bit instruction DLDZ . If it is used with the 16 -bit instruction LDZ※, a program error will occur, and the ERROR LED indicator on the PLC will blink.


## 2, Program Example:

- If the absolute value of the difference between D10 and D11 is greater than K200, YO will be On. If the absolute value of the difference between D10 and D11 is less than or equal to K200, YO will be Off.



1, Explanations:

- S1:Source device $1 \quad$ S2:Source device 2 S3:Source device 3
- The absolute value of the difference between S 1 and S 2 is compared with the absolute value of S 3 . Take ANDZ> for example. If the comparison result is that the absolute value of the difference between Sl and $\$ 2$ is greater than the absolute value of $\$ 3$, the condition of the instruction is met. If the comparison result is that the absolute value of the difference between S 1 and $\$ 2$ is less than or equal to the absolute value of S 3 , the condition of the instruction is not met.
- The instruction ANDZ※ is connected to a contact in series.

| ZL No. | 16-bit instruction | 32-bit instruction | On | Off |
| :---: | :---: | :---: | :---: | :---: |
| 302 | ANDZ> | DANDZ> | $\mid S 1-$ S2 ${ }^{\text {P }}$ \|S3| | $\mid S 1-$ S2\| $\leqq \mid$ S3\| |
| 303 | ANDZ>= | DANDZ>= | $\|S 1-S 2\| \geqq\|S 3\|$ | $\mid S 1-$ S2 $\|<\|S 3\|$ |
| 304 | ANDZ< | DANDZ< | $\mid S 1-$ S2 $\|<\|S 3\|$ | $\mid S 1-$ S2 $\|\geqq\|S 3\|$ |
| 305 | ANDZ<= | DANDZ<= | $\|S 1-S 2\| \leqq\|S 3\|$ | $\|S 1-S 2\|>\|S 3\|$ |
| 306 | ANDZ = | DANDZ = | $\|\mathrm{S} 1-\mathrm{S} 2\|=\|S 3\|$ | $\mid S 1-$ S2 $\|\neq\|S 3\|$ |
| 307 | ANDZ<> | DANDZ<> | $\|\mathrm{S} 1-\mathrm{S} 2\| \neq \mid \mathrm{S3} 3$ | $\mid S 1-$ S2\| $=\mid$ \|S3| |

- A 32-bit counter (C200~C255) must be used with the 32-bit instruction DANDZ . If it is used with the 16 -bit instruction ANDZ※, a program error will occur, and the ERROR LED indicator on the PLC will blink.


## 2, Program Example:

- If MO is On, and the absolute value of the difference between D10 and D11 is greater than K200, YO will be On. If the absolute value of the difference between D10 and D11 is less than or equal to K200, YO will be Off.

| MO |  |  |  |  |
| :---: | :--- | :--- | :--- | :--- |
| -11 | ANDZ $>$ | D10 | D11 | K200 |



- S1:Source device 1 S2: Source device 2 S3: Source device 3
- The absolute value of the difference between S1 and S2 is compared with the absolute value of S3. Take ORZ> for example. If the comparison result is that the absolute value of the difference between S1 and S2 is greater than the absolute value of $\$ 3$, the condition of the instruction is met. If the comparison result is that the absolute value of the difference between $S 1$ and $S 2$ is less than or equal to the absolute value of S3, the condition of the instruction is not met.
- The instruction ORZ $※$ is connected to a contact in parallel.

| ZL No. | 16-bit instruction | 32-bit instruction | On | Off |
| :---: | :---: | :---: | :---: | :---: |
| 308 | ORZ> | DORZ> | $\mid S 1-$ S2 ${ }^{\text {P }}$ \|S3| | $\|S 1-S 2\| \leqq\|S 3\|$ |
| 309 | ORZ>= | DORZ>= | $\mid S 1-$ S2\| $\geqq\|S 3\|$ | $\mid S 1-$ S2\| $<\|S 3\|$ |
| 310 | ORZ< | DORZ< | $\mid S 1-$ S2\| $<\|S 3\|$ | $\|S 1-S 2\| \geqq\|S 3\|$ |
| 311 | ORZ<= | DORZ<= | $\mid S 1-$ S2\| $\leqq\|S 3\|$ | $\mid S 1-$ S2\| $>\|S 3\|$ |
| 312 | ORZ $=$ | DORZ $=$ | $\|S 1-S 2\|=\|S 3\|$ | $\mid S 1-$ S2\| $=\mid$ S3\| |
| 313 | ORZ<> | DORZ<> | $\mid S 1-$ S2\| $\neq\|S 3\|$ | $\mid S 1-$ S2\| $=\|S 3\|$ |

- A 32-bit counter (C200~C255) must be used with the 32-bit instruction DORZ . If it is used with the 16 -bit instruction ORZ※, a program error will occur, and the ERROR LED indicator on the PLC will blink.

2, Program Example:

- If M0 is On, or the absolute value of the difference between D10 and D11 is greater than K200, Y0 will be On.


